



Article An Advanced Crow Search Algorithm for Solving Global Optimization Problem

Donwoo Lee ¹, Jeonghyun Kim ², Sudeok Shon ¹, and Seungjae Lee ^{1,*}

- ¹ School of Industrial Design & Architectural Engineering, Korea University of Technology & Education, 1600 Chungjeol-ro, Byeongcheon-myeon, Cheonan 31253, Republic of Korea; lov1004ely@koreatech.ac.kr (D.L.); sdshon@koreatech.ac.kr (S.S.)
- ² Faculty of Civil Engineering, Wrocław University of Science and Technology, Wybrzeze Wyspianskiego 27, 50-370 Wrocław, Poland; jeonghyun.kim@pwr.edu.pl
- * Correspondence: leeseung@koreatech.ac.kr

Abstract: The conventional crow search (CS) algorithm is a swarm-based metaheuristic algorithm that has fewer parameters, is easy to apply to problems, and is utilized in various fields. However, it has a disadvantage, as it is easy for it to fall into local minima by relying mainly on exploitation to find approximations. Therefore, in this paper, we propose the advanced crow search (ACS) algorithm, which improves the conventional CS algorithm and solves the global optimization problem. The ACS algorithm has three differences from the conventional CS algorithm. First, we propose using dynamic *AP*(awareness probability) to perform exploration of the global region for the selection of the initial population. Second, we improved the exploitation performance by introducing a formula that probabilistically selects the best crows instead of randomly selecting them. Third, we improved the exploration phase by adding an equation for local search. The ACS algorithm proposed in this paper has improved exploitation and exploration performance over other metaheuristic algorithms in both unimodal and multimodal benchmark functions, and it found the most optimal solutions in five engineering problems.

Keywords: advanced crow search algorithm; metaheuristic; convergence performance; engineering problem; benchmark function

1. Introduction

The optimization of engineering problems is of great interest to many researchers, and various strategies for incorporating optimization into the engineering field are being studied [1]. As an example, metaheuristic algorithms that are easy to apply to engineering problems are being developed for optimization. These algorithms are applied to various fields in order to optimize engineering problems by minimizing costs, shortening paths, and maximizing performance.

Metaheuristic algorithms originated in 1965 with the development of the evolution strategy (ES) algorithm [2], and algorithms using various natural phenomena have been proposed. Figure 1 classifies the metaheuristic algorithms based on the natural phenomena that they emulate. Metaheuristic algorithms can be classified into four main categories: evolutionary, swarm, physic, and human behavior [3–7]. Evolution-based algorithms are based on the genetic characteristics and evolutionary methods of nature, and representative algorithms include ES, evolutionary programming (EP), genetic algorithms are based on the behavior of organisms such as birds or ants in clusters, and representative algorithms include ant colony optimization (ACO), particle swarm optimization (PSO), artificial bee colony (ABC), cuckoo search, and crow search (CS). Physical-based algorithms are based on physical phenomena, and representative algorithms include simulated annealing (SA), harmony search (HS), gravitational search (GS), black hole (BH), and sine cosine (SC).



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Copyright: © 2023 by the authors. Licensee MDPI, Basel, Switzerland. This article is an open access article distributed under the terms and conditions of the Creative Commons Attribution (CC BY) license (https:// creativecommons.org/licenses/by/ 4.0/). Finally, the human behavior-based algorithms are based on human intelligent behavior, and representative algorithms include human-inspired (HI), social emotional optimization (SEO), brain storm optimization (BSO), teaching learning-based optimization (TLBO), and social-based (SB) [8]. All metaheuristic algorithms perform optimization using exploitation and exploration. If the metaheuristic algorithm mainly uses exploration, then it can easily find the global minima but has a difficult time finding the exact solution. Conversely, metaheuristic algorithms which mainly use exploitation can find accurate solutions but are prone to falling into local minima [9–11]. Therefore, the convergence performance of the algorithm varies greatly depending on the method of using exploitation and exploration, and exploration should be used in harmony [12].



Figure 1. Classification of metaheuristic algorithms.

Swarm-based algorithms are efficient in searching for global optima and are easy to apply to a variety of optimization problems. They also lend themselves well to parallelization, making them a popular choice for many researchers [13]. With these advantages, swarm-based algorithms are applied to various engineering fields, and many researchers are working to improve the convergence performance of algorithms. The conventional CS algorithm, originally proposed by Askarzadeh in 2016 and ranked among the swarm-based algorithms optimization by mimicking the high intelligence of crows [14]. Crow brains are intelligent enough to recognize food or humans because they are large compared to their body size. As crows are intelligent, they can remember the location of food hidden by other birds and steal this hidden food when the other birds are not around. The conventional CS algorithm proposes to perform optimization using these characteristics of crows and has the following four principles:

- Crows live in groups.
- Crows remember the location of their hidden prey.
- Crows steal food from other birds.
- Crows are protected by probability.

The conventional CS algorithm utilizes a small number of parameters and demonstrates excellent convergence performance. Due to its easy application in problems and excellent performance, it is widely applied in civil and architectural engineering, electrical engineering, mechanical engineering, and image processing [15]. The conventional CS algorithm is more likely to fall into local minima because it mainly performs optimization using exploitation rather than exploration. However, given that real optimization problems are often characterized by multimodal functions, optimization algorithms should mainly use exploration rather than exploitation to find accurate solutions [16]. In order to address this issue, Mohammadi et al. proposed a modified crow search (MCS) algorithm in 2018 that performs optimization through the adoption of a new method for selecting a target crow as well as variation of fl (flight length) based on distance depending on the distance of the crow [17]. In the same year, Díaz et al. proposed the improved crow search (ICS) algorithm, which is improved by random adoption methods using AP (awareness probability) and Lévy flight varying by fitness in the *t* generation [18]. AP is one of the important parameters used in the conventional CS algorithms, and depending on the size of the AP, the conventional CS algorithms perform exploitation or exploration. In 2019, Zamani et al. proposed the conscious neighborhood-based crow search (CCS) algorithm, which utilizes three strategies: neighborhood-based local search (NLS), non-neighborhood-based global search (NGS), and wandering around-based search (WAS) [19]. In the same year, Javidi et al. proposed the enhanced crow search (ECS) algorithm [20], which used three additional mechanisms. In addition, the convergence performance of the ECS algorithm was evaluated compared to the conventional CS algorithm to which three mechanisms were applied. In 2020, Wu et al. proposed the Lévy flight crow search (LFCS) algorithm combining Lévy flight and the conventional CS algorithm, and it utilizes AP, which linearly decreases with the number of generations, and *fl*, which is randomly selected by the parity probability density function [22].

In this paper, the advanced crow search (ACS) algorithm was proposed as a means to solve the global optimization problem. The ACS algorithm uses *dynamicAP*—which varies nonlinearly with changes in the number of generations—and suggests that we follow the best results of previous generations with a probability-based approach, rather than randomly chasing the prey selected by crows. In addition, instead of randomly selecting from the entire problem range, the algorithm proposes reducing the randomly selected space as the number of generations increases. Section 2 briefly reviews conventional CS algorithms and papers that improve on these conventional CS algorithms, and Section 3 compares the explanation of the ACS algorithm with the convergence performance according to parameter changes. In Section 4, we solve the numerical optimization problem and compare the results with those of other algorithms. Section 5 presents the conclusions drawn from this study.

2. Related Work

In this section, we explain the process of optimizing the conventional CS algorithm and briefly outline research projects that have improved upon the conventional CS algorithm.

2.1. Conventional CS Algorithm

The conventional CS algorithm proposed by Askarzadeh describes the intelligent behavior of crows and performs optimization by repeating the following five steps [14]:

Step 1. Define the problem and set the parameters

The problem undergoing optimization is defined, and the initial value of the parameters used in the conventional CS algorithm are set. The parameters used in the conventional CS algorithms are AP (awareness probability), fl (flight length), N (flock size), pd (dimension of problem), and t_{max} (maximum number of generations).

Step 2. Initialize the memory of crows and evaluate

The size of the crow group, determined by pd and the size of N, is expressed as Equation (1), and the initial position of each crow is randomly assigned within the range between lb (lower boundary) and ub (upper boundary). In this context, i is 1, 2, ..., N, t are $1, 2, ..., t_{max}$, and d is pd. The initial position of the randomly placed crow is remembered as Equation (2), and the initial position of the crow is evaluated by object function.

$$Crows_{i,t} = \begin{bmatrix} x_i^1 & \cdots & x_i^d \\ \vdots & \vdots & \vdots \\ x_N^1 & \cdots & x_N^d \end{bmatrix}$$
(1)

$$CrowsMemory_{i,t} = \begin{bmatrix} m_i^1 & \cdots & m_i^d \\ \vdots & \vdots & \vdots \\ m_N^1 & \cdots & m_N^d \end{bmatrix}$$
(2)

Step 3. Generate and evaluate the new positions for crows

Step 3 is the most important step that the conventional CS algorithm uses to perform optimization. Crow $i(x_{i,t})$ follows crow $j(m_{j,t})$, and two cases are proposed depending on whether crow j is aware of crow i's following. The first case is that crow $j(m_{j,t})$ does not recognize crow $i(x_{i,t})$'s following. The position of crow $i(x_{i,t})$ is adjusted by Equation (3), where r_i is a random number between 0 and 1. In addition, a local (fl < 1) or global (fl > 1) area search is performed depending on the size of fl, and it is known to have the best convergence performance when using fl = 2.0. Figure 2 is a diagram that expresses this characteristic.



(a)

Figure 2. Comparison results of benchmark function: (a) fl < 1. (b) fl > 1.

$$x_{i,t+1} = x_{i,t} + r_i \times fl \times (m_{j,t} - x_{i,t})$$
(3)

$$x_{i,t+1} = a \text{ random position} \tag{4}$$

The second case is that crow $j(m_{j,t})$ recognizes crow $i(x_{i,t})$'s following. In this instance, the position of crow i is adjusted by Equation (4), moving to a random position in the range between lb and ub. Two cases are selected from each generation by the AP, and the AP mainly uses 0.1. Relative to the size of AP, the conventional CS algorithms perform exploitation and exploration in order to find the optimal solutions. The positions of the newly moved crows are again evaluated by the objective function.

Step 4. Update the memory

The results are compared by evaluating the crow position change using Equations (3) and (4) with the evaluation of crows stored in memory. Comparing the evaluation results, the better crow position is updated in the crow's memory.

Step 5. Termination of repetition

The process of Steps 2–4 is repeated continuously, and when *t* reaches t_{max} , the performance of the conventional CS algorithm is terminated in order to derive optimization results. The pseudo code of the above-mentioned process is provided in Algorithm 1.

Algorithm 1 Pseudo code of the conventional CS algorithm

Initialize the parameters(*AP*, *fl*, *N*, *pd*, *t*_{max}) Initialize the position of crows in the search space and memorize Evaluate the position of crows while $t < t_{max}$ do Randomly choose the position of crows for i = 1 : N do if $r_i \geq AP$ then $x_{i,t+1} = x_{i,t} + r_i \times fl \times (m_{i,t} - x_{i,t})$ else $x_{i,t+1} = a$ random position end if end for Evaluate the new position of crows Update the memory of crows end while Show the results

2.2. Modified CS Algorithm

The modified CS (MCS) algorithm was proposed by Mohammadi et al. in 2018 [17]. The MCS algorithm has a similar structure compared to the conventional CS algorithm, but two new equations have been proposed.

First, MCS algorithm uses *K* parameters, which use random variables between '0' and '1' to select the target crow (Crow *j*), unlike the conventional CS algorithm. *K* is defined as Equation (5) and consists of K_{max} and K_{min} . *K* has values that decrease with the number of generations by K_{max} and K_{min} .

$$K_t = round \left(K_{max} - \frac{K_{max} - K_{min}}{t_{max}} \times t \right)$$
(5)

If *K* has a large value, then the probability that a crow in a bad position will be selected increases; if *K* has a small value, then the probability that the crow in the best position will be selected increases. Therefore, exploration is primarily performed in the initial generations, and exploitation is primarily performed in the latter generations.

Second, the MCS algorithm uses a value of fl differently depending on the distance between crow *i* and crow *j*, where fl is defined as Equation (6). Here, fl_{thr} and D_{thr} are initially set parameters, and $D_{i,j}$ is the distance vector of crow *i* and crow *j*.

$$fl_{i,t} = \begin{cases} 2 & if \quad D_{i,j} > D_{thr} \\ fl_{thr} & if \quad D_{i,j} \le D_{thr} \end{cases}$$
(6)

Askarzadeh noted that when fl = 2, the conventional CS algorithm has the best convergence performance [14]. However, the MCS algorithm uses $fl_{i,t}$ with a value greater than 2 when the crow's distance $(D_{i,j})$ is closer than D_{thr} .

2.3. Dynamic CS Algorithm

The dynamic CS (DCS) algorithm was proposed by Necira et al. in 2022 [22], and it proposed dynamic *AP* and *fl* that change with the number of generations.

First, dynamic *AP*, which varies dynamically with the number of generations, is defined as Equation (7). dynamic *AP* decreases linearly within the range of AP_{max} and AP_{min} as the number of generations increases. This change causes the initial number of generations to perform the exploration in the global search space.

$$AP = AP_{max} + \frac{AP_{max} - AP_{min}}{t_{max}} \times t \tag{7}$$

Second, fl_c was used instead of fl used by the conventional CS algorithm, which is defined as Equation (8). The conventional CS algorithm initially determines fl and performs a local search or global search based on the determined value. However, DCS algorithm mainly performs a global search when it is less than a certain number of generations, and a local search when it exceeds a certain number of generations. These changes are determined by τ and are mainly used at 0.9.

$$fl_{c} = \begin{cases} fl \times \left[F\left(\frac{y_{max}}{10}\right) - F(y_{min}) \times r\right] & if \quad t \le \tau \times t_{max} \\ fl \times \left[F(y_{max}) - F(y_{max}) - F(0.6 \times y_{max})\right] & else \end{cases}$$
(8)

3. Proposed Method

3.1. Advanced CS Algorithm

The conventional CS algorithm, which repeats the above process to perform optimization, performs exploration and exploitation according to the size of AP, mainly using 0.1 for AP. That is, the conventional CS algorithm mainly performs the exploitation rather than the exploration. Figure 3 is a diagram showing the exploitation and exploration that occurs in the process of optimizing the Sphere function for 1000 generations of the conventional CS algorithms with a N of 50. It can be seen that exploitation mainly occurs in all generations. Optimization algorithms that mainly use excitation in optimization performance are likely to fall into local minima [10], and the performance of the conventional CS algorithms is largely dependent on the initial population. In this paper, to address this problem, we improve the performance of the initial population using dynamic AP that varies dynamically with the number of generations, and the performance of exploitation and exploration using two proposed equations.



Figure 3. Exploitation and exploration of the conventional CS algorithm.

Similar to the conventional CS algorithm, the ACS algorithm consists of a total of five steps.

Step 1. Define the problem and set the parameters

Like the conventional CS algorithm, the problem for performing optimization is defined in Step 1, and the parameters used in the ACS algorithm are set. The parameters added in the ACS algorithm are AP_{max} , AP_{min} , and FAR (Flight Awareness Ratio). Here, AP_{max} and AP_{min} are used for dynamic AP.

Step 2. Initialize the memory of crows and evaluate

The size of the crew group used in the ACS algorithm is expressed as Equation (1) as in the conventional CS algorithm, and the initial position is remembered as Equation (2). The initial position of the remembered crow is evaluated by the objective function.

Step 3. Generate and evaluate the new positions for crows

The ACS algorithm displays the biggest difference from the conventional CS algorithm in Step 3. First, The ACS algorithm uses dynamic *AP*, which changes dynamically with the number of generations. dynamic *AP* uses Equation (9) for dynamic changes, and *AP_{max}* and *AP_{min}* have a value between 0 and 1. Figure 4 shows an *AP* that changes dynamically according to the number of generations when t_{max} is 2000. Using dynamic *AP*, as shown in Figure 5, increases the probability of exploration at the beginning of the generation, which can increase the performance of the initial population. Compared to Figure 3, the number of explorations increases at lower numbers of generations. Thus, the larger the *AP*, the higher the probability of the initial population performing exploration, and the smaller the *AP*, the higher the probability of performing exploitation. In addition, a dynamic *AP* of an appropriate size is required for harmony between exploitation and exploration.



$$AP_t = AP_{min} + \frac{AP_{max} - AP_{min}}{ln(t) + 1}$$
⁽⁹⁾

Figure 4. *dynamic* AP of ACS algorithm($AP_{max} = 0.4$, $AP_{min} = 0.01$).



Figure 5. Exploitation and exploration of ACS algorithm($AP_{max} = 1.0, AP_{min} = 0.1$).

Second, unlike the conventional Equation (3) in which crow *i* follows randomly selected crow *j* ($m_{j,t}$), in the ACS algorithm, it follows the best crow *j* ($gb_{j,t}$) by *FAR*. This can be expressed as Equation (10). Here, $r_{i,t}^2$, $r_{i,t}^3$ is a random number between 0 and 1, and *FAR* is an initial set value between 0 and 1. The change in this equation improves the exploitation performance compared to the conventional CS algorithm. If *FAR* approaches 0,

it follows the best solution stored in the crow's memory. Conversely, when *FAR* approaches 1, it follows a randomly selected crow, just like the conventional CS algorithm. Therefore, using the appropriate *FAR*, it is possible to improve the convergence performance of the optimization algorithm by harmonizing the exploitation and exploration.

$$x_{i,t+1} = \begin{cases} x_{i,t} + r_{i,t}^2 \times fl \times (m_{j,t} - x_{j,t}) & r_{i,t}^3 \le FAR \\ x_{i,t} + r_{i,t}^2 \times fl \times (gb_{j,t} - x_{j,t}) & else \end{cases}$$
(10)

Third, using this algorithm, the exploration phase of the conventional CS algorithm was improved. The conventional CS algorithms are randomly adopted in the *lb* and *ub* ranges if the random number is less than the *AP*. That is, global search is mainly performed. The global search can contribute to the convergence performance of the algorithm because it searches a large area at the beginning of the generation. However, it does not contribute significantly to the convergence performance of the algorithm as the generation progresses. Therefore, the process of reducing the range that can be selected toward the end of the generation was added as Equation (11), which allows the ACS algorithm to perform a local search. Here, $r_{i,t}^4$ and $r_{i,t}^5$ are random numbers between 0 and 1. Figure 6 illustrates this method.

$$x_{i,t+1} = \begin{cases} 2x_{i,t} + (lb + r_{i,t}^5 \times (lb - ub))/t & r_{i,t}^4 < 0.5\\ a \text{ random position} & else \end{cases}$$
(11)



Figure 6. Random position of ACS algorithm.

Step 4. Update the memory

The results are compared through evaluation of the crow position change by Equation (3), Equation (4) with the evaluation of crows stored in memory. Comparing the evaluation results, the better crow position is updated in the crow's memory.

Step 5. Termination of repetition

The ACS algorithm performs optimization by repeating the process of Steps 2–4. When the current number of generations (t) reaches the maximum number of generations (t_{max}), the execution of the ACS algorithm ends, and the optimization result of the problem is derived. Pseudo code of the above-mentioned process is provided in Algorithm 2.

Algorithm 2 Pseudo code of the ACS algorithm

Initialize the parameters(*AP_{max}*, *AP_{min}*, *FAR*, *fl*, *N*, *pd*, *t_{max}*) Initialize the position of crows in the search space and memorize Evaluate the position of crows while $t < t_{max}$ do Randomly choose the position of crows **for** i = 1 : N **do** if $r_{i,t}^1 \ge AP_t$ then if $r_{i,t}^3 \leq FAR$ then $x_{i,t+1} = x_{i,t} + r_{i,t}^2 \times fl \times (m_{j,t} - x_{i,t})$ else $x_{i,t+1} = x_{i,t} + r_{i,t}^2 \times fl \times (gb_{j,t} - x_{i,t})$ end if else if $r_{i,t}^4 \leq 0.5$ then $x_{i,t+1} = 2x_{i,t} + (lb + r_{i,t}^5 \times (lb - ub))/t$ else $x_{i,t+1} = a$ random position end if end if end for Evaluate the new position of crows Update the memory of crows end while Show the results

3.2. Characteristic of the ACS Algorithm

Unlike the conventional CS algorithm, the ACS algorithm adds the parameters of *dynamic AP* and *FAR*. Therefore, this section compares the convergence performance according to the change in the newly added parameters and seeks the value with the best convergence performance. The benchmark function was used to compare convergence performance, and it was summarized in Table 1. Here, *d* was set to 10 in order to identify the characteristics of the ACS algorithm.

Table 1. Benchmark function for comparison.

Fun	Equation	В	Min
f1	$f(x) = \sum_{i=1}^n x_i^2$	$[-100\ 100]^{d}$	0
f2	$f(x) = \sum_{i=1}^{n} x_i + \prod_{i=1}^{n} x_i $	$[-10\ 10]^{d}$	0
f3	$f(x) = \sum_{i=1}^{n} (\sum_{j=1}^{i} x_j)^2$	$[-100\ 100]^{d}$	0
f4	$f(x) = max\{ x_i , 1 \le i \le n\}$	$[-100\ 100]^{d}$	0
f5	$f(x) = \sum_{i=1}^{n-1} \left[100 (x_{i=1} - x_i^2)^2 + (x_i - 1)^2 \right]$	$[-30\ 30]^{d}$	0
<i>f</i> 6	$f(x) = \sum_{i=1}^{n} ([x_i + 0.5])^2$	$[-100\ 100]^{d}$	0
f7	$f(x) = \sum_{i=1}^{n} ix_i^4 + rand(0, 1)$	$[-1.28\ 1.28]^{d}$	0
f8	$f(x) = \sum_{i=1}^{n} -x \sin \sqrt{ x_i }$	$[-500\ 500]^{d}$	$-418.9829 \times d$
f9	$f(x) = \sum_{i=1}^{n} \left[x_i^2 - 10 \cos(2\pi x_i) \right] + 10$	$[-5.12\ 5.12]^{d}$	0
<i>f</i> 10	$f(x) = -20 \exp\left(-0.2 \sqrt{\frac{1}{n} \sum_{i=1}^{n} x_i^2}\right) - \exp\left(\frac{1}{n} \sum_{i=1}^{n} \cos(2\pi x_i)\right) + 20 + e$	$[-32\ 32]^{d}$	0
f11	$f(x) = \frac{1}{4000} \sum_{i=1}^{n} x_i^2 - \prod_{i=1}^{n} \cos\left(\frac{x_i}{\sqrt{i}}\right)$	$[-600\ 600]^{d}$	0
<i>f</i> 12	$f(x) = \frac{\pi}{n} \left\{ 10\sin^2(\pi y_1) + \sum_{i=1}^{n-1} (y_i - 1)^2 \left[1 + 10\sin^2(\pi y_{i+1}) \right] + (y_n - 1)^2 \right\} + \sum_{i=1}^n u(x_i, a, k, m)$	$[-50\ 50]^{d}$	0
f13	$f(x) = 0.1\left\{\sin^2(3\pi x_1) + \sum_{i=1}^{n-1} (x_i - 1)^2 \left[1 + \sin^2(3\pi x_{i+1})\right] + (x_n - 1)^2 \left[1 + \sin^2(2\pi x_n)\right]\right\} + \sum_{i=1}^n u(x_i, a, k, m)$	[-50 50] ^d	0

A total of 13 functions were used to compare the convergence performance according to the value of the added parameter. In Table 1, f1-f7 is a unimodal benchmark function that can test the exploitation performance of each algorithm. Additionally, f8-f13 is a multimodal benchmark function that can test the exploration performance of each algorithm. The multimodal benchmark function has many local minima, making it difficult to find an exact solution.

3.2.1. Dynamic AP

The ACS algorithm uses *dynamic* AP, which varies with the number of generations, to increase the performance of the exploration initially. *dynamic* AP is calculated by Equation (9) and has a different value depending on the size of the AP_{max} . Figure 7 is a graph that changes according to the size of the AP_{max} . The larger the AP_{max} , the higher the probability of randomly selecting the entire boundary initially and the better the initial population selection. Therefore, this section compares results that change according to the value of the AP_{max} .



Figure 7. Dynamic AP according to AP_{max}.

When *AP* becomes 0, only exploitation occurs in all generations. Therefore, the *AP*_{min} was set to a minimum value of (=0.01). *AP*_{max} was changed to 0.01, 0.1, 0.2, 0.4, 0.6, 0.8, and 1.0, and *N*, *fl*, and *FAR* were set to 20, 2.0, and 1.0. t_{max} was set to 2000, and each analysis was repeated a total of 50 times.

Table 2 presents the analysis result of each benchmark function according to the change of AP_{max} , and the last row indicates the average ranking of the BF (best fitness) or MF (mean fitness) according to the AP_{max} . If two or more values were ranked the same, then the average ranking was derived. The average ranking of BF was best at 1.88 when $AP_{max} = 0.4$, and the average ranking of MF was best at 2.31 when $AP_{max} = 0.8$. Conversely, when $AP_{max} = 0.01$, both BF and MF performance deteriorated. In other words, using *dynamic AP* as an appropriate value yields better convergence performance than the conventional CS algorithms, and the convergence performance of the ACS algorithm is the best when the *dynamic AP* has a range of 0.4–0.6.

Fun

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on results according to <i>AP_{max}</i> .							
	AP _{max}						
.2	0.4	0.6	0.8	1.0			
10 ⁻²⁶	2.593×10^{-24}	1.061×10^{-21}	1.529×10^{-18}	2.833×10^{-17}			
10^{-22}	2.924×10^{-21}	1.009×10^{-18}	$2.040 imes10^{-16}$	1.780×10^{-14}			
10^{-21}	6.870×10^{-21}	$2.496 imes 10^{-18}$	$3.732 imes 10^{-16}$	$3.336 imes10^{-14}$			
$\times 10^{-9}$	$1.453 imes 10^{-9}$	$2.523 imes 10^{-8}$	$5.769 imes10^{-8}$	$4.234 imes10^{-7}$			
$< 10^{-2}$	5.967×10^{-3}	3.480×10^{-3}	1.651×10^{-3}	1.669×10^{-2}			
$< 10^{-1}$	2.146×10^{-2}	1.474×10^{-2}	6.987×10^{-3}	5.230×10^{-2}			
$< 10^{-15}$	1.478×10^{-15}	$6.072 imes 10^{-14}$	3.762×10^{-13}	$9.214 imes10^{-11}$			
$< 10^{-10}$	2.695×10^{-12}	$1.810 imes10^{-11}$	$1.778 imes 10^{-10}$	$2.906 imes10^{-9}$			
0	10	4.4	10	0			

Table 2. Benchmark functi

0.01 0.1 0 2.262×10^{-25} BF 1.539×10^{-16} 2.565 × 1.210×10^{-11} 3.174×10^{-20} f1MF 4.473 × Std 3.856×10^{-11} 1.148×10^{-19} 2.366 × BF 9.000×10^{-4} 2.982×10^{-7} 3.825 > 4.804×10^{-1} $3.235\times~10^{-1}$ f2MF 5.726 > Std 8.730×10^{-1} $6.932\times~10^{-1}$ 1.713 > 8.264×10^{-7} $3.387 imes 10^{-12}$ BF 3.318 > 7.755×10^{-4} 3.325×10^{-8} fЗ MF 2.260 × 1.426×10^{-3} $8.402 imes 10^{-8}$ Std 4.797×10^{-12} 3.574×10^{-11} 3.233×10^{-10} 1.294×10^{-9} 7.114×10^{-9} 1.288×10^{-6} 7.312×10^{-8} 5.775×10^{-8} 3.196×10^{-8} 4.060×10^{-7} 1.179×10^{-6} BF 6.825×10^{-4} 1.968×10^{-5} 2.281×10^{-1} $1.375\times~10^{-3}$ $4.209 imes 10^{-5}$ $3.244 imes 10^{-6}$ 3.038×10^{-6} 3.329×10^{-5} f4MF 5.381×10^{-1} 5.247×10^{-3} 9.829×10^{-6} 4.470×10^{-6} 5.200×10^{-5} 5.856×10^{-5} Std 1.190×10^{-4} 3.408×10^{-1} 8.382×10^{-2} 1.713×10^{0} $4.013\times~10^{-1}$ 5.917×10^{-1} 1.636×10^{-1} 2.719×10^{-1} BF f5MF 4.907×10^{1} 3.163×10^{1} 1.235×10^{1} 5.890×10^{0} 1.942×10^{1} 4.255×10^{0} 6.337×10^{0} 1.146×10^2 Std 7.602×10^{1} 3.155×10^{1} 1.722×10^{1} 6.666×10^{1} 1.329×10^{1} 1.842×10^{1} 1.153×10^{-15} 4.338×10^{-25} 1.302×10^{-26} 8.952×10^{-24} $1.329 \times \ 10^{-21}$ 3.406×10^{-18} 4.081×10^{-16} BF 1.516×10^{-11} $1.065\times~10^{-20}$ $5.884\times~10^{-23}$ 3.664×10^{-21} 5.147×10^{-19} $2.082 imes 10^{-14}$ *f*6 MF 1.780×10^{-16} 8.877×10^{-11} 3.800×10^{-20} $1.452\times~10^{-22}$ 1.099×10^{-20} 6.811×10^{-19} 3.103×10^{-16} 2.444×10^{-14} Std $1.402\times~10^{-3}$ 3.638×10^{-3} $2.790 imes 10^{-4}$ 2.662×10^{-4} $4.150 imes 10^{-4}$ 5.504×10^{-4} $5.838 imes 10^{-4}$ BF 1.861×10^{-2} 3.338×10^{-3} f7 MF 7.457×10^{-3} 5.353×10^{-3} 3.847×10^{-3} 3.086×10^{-3} 3.419×10^{-3} 1.243×10^{-2} $2.699\times~10^{-3}$ $4.381\times~10^{-3}$ $3.128\times~10^{-3}$ $2.168\times~10^{-3}$ 1.838×10^{-3} 2.241×10^{-3} Std BF -3.517×10^{3} -3.475×10^{3} -3.616×10^{3} -3.835×10^{3} -3.736×10^{3} -3.617×10^{3} -3.476×10^{3} f8MF -2.618×10^{3} -2.790×10^{3} -2.784×10^{3} -2.797×10^{3} -2.799×10^{3} -2.911×10^{3} -2.816×10^{3} 4.082×10^2 3.450×10^{2} 3.783×10^{2} 3.988×10^{2} 3.013×10^2 Std 3.896×10^{2} 3.289×10^{2} BF 5.970×10^{0} 4.975×10^{0} 5.970×10^{0} 4.975×10^{0} 3.980×10^{0} 5.970×10^{0} 5.970×10^{0} f9 MF 2.507×10^{1} 2.454×10^{1} 2.306×10^{1} 2.255×10^{1} 1.988×10^{1} 2.366×10^{1} 2.312×10^{1} Std 1.058×10^{1} 1.282×10^{1} 1.374×10^{1} 1.164×10^{1} 1.095×10^{1} 1.166×10^{1} 1.061×10^{1} 2.013×10^{0} $1.150 imes 10^{-5}$ 1.131×10^{-12} 2.077×10^{-11} 1.267×10^{-10} $2.340 imes 10^{-9}$ $1.762 imes 10^{-8}$ BF 2.436×10^{0} 2.361×10^0 f10 3.965×10^{0} 2.984×10^{0} 2.625×10^{0} 2.100×10^{0} 2.096×10^{0} MF 1.073×10^{0} $1.097 imes 10^{0}$ $8.876\times~10^{-1}$ 9.702×10^{-1} $8.554\times~10^{-1}$ 9.932×10^{-1} 1.000×10^{0} Std $7.378\times~10^{-2}$ BF 7.874×10^{-2} 5.899×10^{-2} 6.637×10^{-2} 9.106×10^{-2} 8.115×10^{-2} 7.132×10^{-2} $7.107\times~10^{-1}$ f11 $7.328\times~10^{-1}$ $5.611\times~10^{-1}$ 3.788×10^{-1} 4.330×10^{-1} $3.077\times~10^{-1}$ $2.621\times~10^{-1}$ MF 4.147×10^{-1} $3.670\times~10^{-1}$ 3.402×10^{-1} 2.323×10^{-1} 2.936×10^{-1} 1.644×10^{-1} 1.784×10^{-1} Std $2.089\times~10^{-3}$ 1.272×10^{-5} 2.475×10^{-15} 1.017×10^{-16} 2.176×10^{-15} 3.198×10^{-14} 2.542×10^{-12} BF f12 6.269×10^{0} MF 1.001×10^{1} 6.160×10^{0} 2.527×10^{0} 1.988×10^{0} 1.685×10^{0} 2.078×10^{0} 7.024×10^{0} 5.187×10^{0} $6.034 imes 10^{0}$ 3.390×10^{0} 3.928×10^{0} 3.189×10^{0} 3.076×10^{0} Std 1.144×10^{-8} 3.114×10^{-14} 2.572×10^{-18} 4.356×10^{-19} 9.371×10^{-16} 7.594×10^{-15} 5.353×10^{-13} BF 1.994×10^{-2} $8.973\times~10^{-3}$ $6.078\times~10^{-3}$ 6.096×10^{-3} 6.679×10^{-3} $4.364\times~10^{-3}$ 6.731×10^{-3} f13 MF 3.138×10^{-2} $1.057\times~10^{-2}$ 6.584×10^{-3} 7.094×10^{-3} Std 8.789×10^{-3} 7.714×10^{-3} 6.168×10^{-3} BF 4.58 2.85 4.08 5.23 6.77 2.62 1.88 Ranking MF 6.92 5.54 4.08 2.85 2.92 2.31 3.38

3.2.2. FAR

The ACS algorithm follows a randomly selected crow $(m_{i,t})$ by FAR or a crow $(gb_{i,t})$ with favorite prey. The closer FAR = 1.0 is, the more likely it is to follow a randomly selected crow $(m_{i,t})$ like the conventional CS algorithm, and the closer it is to FAR = 0.0 the more likely is to follow a crow $(gb_{j,t})$ with favorite prey. In this section, *FAR* was changed to 0.0, 0.2, 0.4, 0.6, 0.8, and 1.0 in order to compare convergence performance with changes in FAR, and N, fl, AP_{max}, and AP_{min} were set to 20, 2.0, 0.4, and 0.01, respectively. t_{max} was set to 2000, and each analysis was repeated a total of 100 times.

Table 3 presents the analysis results according to a change in FAR. The mean ranking of BF was the best at 2.04 when FAR = 0.2, and the mean ranking of MF was the best at 2.92 when FAR = 0.6. Conversely, the closer FAR = 0.0 or 1.0, the worse the average ranking. Furthermore, the closer the local minima were to FAR = 0.0 in F1, F4, and F6, the better the convergence performance. In other words, using the appropriate value of FAR yielded better convergence performance than the conventional CS algorithm, and when FAR had a range of 0.2–0.4, it had the best convergence performance.

_		FAR						
Fun	Index -	0.0	0.2	0.4	0.6	0.8	1.0	
	BF	1.787×10^{-20}	1.754×10^{-20}	3.271×10^{-19}	1.532×10^{-17}	$6.406 imes 10^{-16}$	$1.360 imes 10^{-12}$	
f1	MF	3.076×10^{-18}	2.509×10^{-18}	3.189×10^{-17}	$1.136 imes 10^{-15}$	$4.515 imes10^{-14}$	$2.813 imes10^{-11}$	
	Std	5.849×10^{-18}	3.748×10^{-18}	$6.464 imes 10^{-17}$	2.002×10^{-15}	$6.196 imes10^{-14}$	$4.601 imes 10^{-11}$	
	BF	$7.143 imes10^{-4}$	$6.113 imes 10^{-8}$	5.538×10^{-8}	$8.745 imes10^{-8}$	$1.645 imes10^{-7}$	$2.121 imes 10^{-6}$	
f2	MF	$7.625 imes10^{-1}$	4.168×10^{-2}	2.789×10^{-3}	1.890×10^{-2}	3.333×10^{-3}	2.349×10^{-2}	
	Std	$9.166 imes 10^{-1}$	$2.084 imes 10^{-1}$	1.214×10^{-2}	$1.300 imes10^{-1}$	1.657×10^{-2}	8.600×10^{-2}	
	BF	$1.318 imes10^{-10}$	$1.385 imes10^{-13}$	$1.723 imes 10^{-13}$	2.560×10^{-12}	$2.984 imes10^{-11}$	$1.954 imes10^{-8}$	
f3	MF	$9.056 imes10^{-9}$	$2.084 imes10^{-11}$	$8.017 imes10^{-11}$	$9.212 imes 10^{-10}$	$1.486 imes10^{-8}$	$5.695 imes10^{-6}$	
	Std	$1.382 imes10^{-8}$	$2.847 imes 10^{-11}$	$1.780 imes 10^{-10}$	$1.767 imes10^{-9}$	$2.950 imes10^{-8}$	$9.067 imes10^{-6}$	
	BF	$1.954 imes10^{-8}$	$6.332 imes 10^{-8}$	$1.195 imes 10^{-7}$	$2.276 imes10^{-7}$	$1.533 imes10^{-6}$	$8.818 imes10^{-6}$	
f4	MF	$5.695 imes10^{-6}$	$8.908 imes10^{-6}$	$8.421 imes 10^{-6}$	$1.605 imes10^{-5}$	$2.570 imes 10^{-5}$	$1.203 imes10^{-4}$	
	Std	$9.067 imes10^{-6}$	$1.448 imes 10^{-5}$	1.102×10^{-5}	$2.305 imes10^{-5}$	$3.901 imes 10^{-5}$	$1.628 imes 10^{-4}$	
	BF	$6.351 imes10^{-1}$	$2.218 imes10^{-1}$	$1.365 imes10^{-1}$	$1.857 imes10^{-1}$	$5.801 imes10^{-1}$	$4.578 imes10^{-1}$	
f5	MF	$3.904 imes10^1$	$1.206 imes 10^1$	$5.356 imes 10^0$	$8.779 imes 10^0$	$1.262 imes 10^1$	$1.656 imes 10^1$	
	Std	$9.006 imes10^1$	$3.350 imes10^1$	$5.356 imes 10^0$	$2.897 imes10^1$	$3.539 imes10^1$	$4.085 imes10^1$	
	BF	2.275×10^{-20}	1.137×10^{-20}	$9.394 imes 10^{-20}$	$9.199 imes 10^{-19}$	$2.792 imes 10^{-16}$	$1.856 imes 10^{-12}$	
<i>f</i> 6	MF	2.655×10^{-18}	4.212×10^{-18}	3.363×10^{-17}	$7.367 imes 10^{-16}$	$4.578 imes 10^{-14}$	$2.756 imes 10^{-11}$	
	Std	4.478×10^{-18}	$7.407 imes 10^{-18}$	7.556×10^{-17}	$1.182 imes 10^{-15}$	$8.264 imes 10^{-14}$	$3.879 imes 10^{-11}$	
	BF	$6.747 imes10^{-4}$	$2.623 imes 10^{-4}$	3.859×10^{-4}	$4.041 imes10^{-4}$	$3.774 imes10^{-4}$	$6.165 imes10^{-4}$	
f7	MF	4.244×10^{-3}	3.991×10^{-3}	4.146×10^{-3}	3.277×10^{-3}	3.619×10^{-3}	$4.088 imes 10^{-3}$	
	Std	2.879×10^{-3}	2.630×10^{-3}	2.624×10^{-3}	2.069×10^{-3}	2.272×10^{-3}	$2.417 imes 10^{-3}$	
	BF	-3.953×10^{3}	-3.595×10^{3}	-3.953×10^{3}	-3.716×10^{3}	-3.953×10^{3}	-4.071×10^{3}	
f8	MF	-2.808×10^{3}	-2.792×10^{3}	-2.804×10^{3}	-2.830×10^{3}	-2.887×10^{3}	-2.923×10^{3}	
	Std	4.065×10^{2}	3.371×10^{2}	3.740×10^{2}	3.620×10^{2}	3.549×10^2	3.878×10^{2}	
	BF	$4.975 imes 10^0$	$4.975 imes 10^0$	3.980×10^{0}	$3.980 imes 10^0$	$3.980 imes 10^0$	3.980×10^{0}	
<i>f</i> 9	MF	$2.983 imes 10^1$	$2.320 imes 10^1$	$1.996 imes 10^1$	$2.016 imes 10^1$	$1.617 imes 10^1$	$1.135 imes 10^1$	
	Std	$1.272 imes 10^1$	$1.105 imes 10^1$	8.363×10^{0}	$9.589 imes 10^0$	$7.980 imes 10^0$	5.342×10^0	
	BF	$1.033 imes 10^{-9}$	$6.115 imes 10^{-10}$	2.005×10^{-9}	$2.950 imes 10^{-9}$	$3.231 imes 10^{-8}$	5.811×10^{-7}	
f10	MF	$2.446 imes 10^0$	2.473×10^{0}	2.325×10^{0}	2.057×10^{0}	$2.063 imes 10^{0}$	1.666×10^{0}	
	Std	$8.111 imes 10^{-1}$	7.402×10^{-1}	7.449×10^{-1}	$8.948 imes 10^{-1}$	$9.754 imes 10^{-1}$	9.920×10^{-1}	
	BF	6.149×10^{-2}	2.219×10^{-2}	2.709×10^{-2}	6.886×10^{-2}	3.937×10^{-2}	3.446×10^{-2}	
f11	MF	$4.724 imes 10^{-1}$	$4.276 imes 10^{-1}$	$4.516 imes 10^{-1}$	$3.836 imes 10^{-1}$	$2.900 imes 10^{-1}$	1.898×10^{-1}	
	Std	$2.797 imes 10^{-1}$	2.809×10^{-1}	$2.783 imes 10^{-1}$	2.295×10^{-1}	$1.468 imes 10^{-1}$	9.766×10^{-2}	
	BF	1.667×10^{-13}	3.925×10^{-16}	$1.685 imes 10^{-14}$	9.559×10^{-14}	2.608×10^{-11}	$7.787 imes 10^{-11}$	
f12	MF	$4.175 imes 10^{0}$	3.322×10^{0}	2.212×10^{0}	3.049×10^{0}	2.289×10^{0}	$2.446 imes 10^{0}$	
	Std	3.913×10^{0}	5.352×10^{0}	3.187×10^{0}	3.988×10^{0}	3.595×10^{0}	3.381×10^{0}	
	BF	1.171×10^{-14}	4.129×10^{-16}	6.199×10^{-15}	3.664×10^{-14}	2.733×10^{-13}	5.752×10^{-11}	
f13	MF	6.429×10^{-3}	6.192×10^{-3}	6.377×10^{-3}	4.708×10^{-3}	5.325×10^{-3}	5.285×10^{-3}	
	Std	8.209×10^{-3}	7.961×10^{-3}	8.446×10^{-3}	6.981×10^{-3}	6.743×10^{-3}	7.286×10^{-3}	
Ranking	BF	3.85	2.04	2.31	3.73	4.27	4.81	
	MF	4.54	3.69	3.08	2.92	3.23	3.54	

Table 3. Benchmark function results according to FAR.

4. Numerical Examples

In this section, the ACS algorithm was applied to benchmark function and engineering problems and compared with the results of other algorithms. The benchmark function used 23 functions shown in Tables 1 and 4 [23]. Five engineering problems were performed: a pressure vessel design problem (PVD), a welded beam design problem, a weight of a tension/compression string problem, a three-bar truss optimization problem, and a stepped cantilever beam design problem.

Fun	Equation	В	Min
<i>f</i> 14	$f(\mathbf{x}) = \left(\frac{1}{500} + \sum_{j=1}^{25} \frac{1}{j + \sum_{i=1}^{2} (x_i - a_{ij})^6}\right)^{-1}$	$[-65\ 65]^2$	1
f15	$f(\mathbf{x}) = \sum_{i=1}^{11} \left[a_i - \frac{x_1(b_i^2 + b_i x_2)}{b_i^2 + b_i x_3 + x_4} \right]^2$	$[-5\ 5]^4$	0.0003
f16	$f(\mathbf{x}) = 4x_1^2 - 2.1x_1^4 + \frac{1}{3}x_1^6 + x_1x_2 - \frac{1}{4}x_2^2 + 4x_2^4$	$[-5\ 5]^2$	-1.0316
f17	$f(\mathbf{x}) = \left(x_2 - \frac{5.1}{4\pi^2}x_1^2 + \frac{5}{\pi}x_1 - 6\right)^2 + 10\left(1 - \frac{1}{8\pi}\right)\cos x_1 + 10$	$[-5 5]^2$	0.398
f18	$f(\mathbf{x}) = \begin{bmatrix} 1 + (x_1 + x_2 + 1)^2 (19 - 14x_1 + 3x_1^2 - 14x_2 + 6x_1x_2 + 3x_2^2) \end{bmatrix}$	$[-2 2]^2$	3
	$\times \left[30 + (2x_1 - 3x_2)^2 \left(18 - 32x_1 + 12x_1^2 + 48x_2 - 36x_1x_2 + 27x_2^2 \right) \right]$		
f19	$f(\mathbf{x}) = -\sum_{i=1}^{4} c_{i} \exp\left(-\sum_{j=1}^{3} a_{ij} (x_{j} - p_{ij})^{2}\right)$	[1 3] ³	-3.86
f20	$f(\mathbf{x}) = -\sum_{i=1}^{4} c_{i} \exp\left(-\sum_{j=1}^{6} a_{ij} (x_{j} - p_{ij})^{2}\right)$	$[0\ 1]^6$	-3.32
f21	$f(\mathbf{x}) = -\sum_{i=1}^{5} \left[(X - a_i) (X - a_i)^T + c_i \right]^{-1}$	$[0\ 10]^4$	-10.1532
f22	$f(\mathbf{x}) = -\sum_{i=1}^{7} \left[(X - a_i)(X - a_i)^T + c_i \right]^{-1}$	$[0\ 10]^4$	-10.4028
f23	$f(\mathbf{x}) = -\sum_{i=1}^{10} \left[(X - a_i)(X - a_i)^T + c_i \right]^{-1}$	$[0\ 10]^4$	-10.5363

Table 4. Fixed-dimension multimodal benchmark function for comparison.

4.1. Benchmark Function Problems

The algorithms used to compare the convergence performance of the ACS algorithm were the conventional CS algorithm, HS, DE, the grasshopper optimization (GO) algorithm, the salp swarm (SS) algorithm, and GA. Table 5 presents the parameters used in each algorithm. t_{max} , N, and Dim used 2000, 30, and 30, respectively, and each analysis was repeated a total of 30 times.

Table 5. Parameters for benchmark function problems.

Algorithm	Parameters	
ACS	$fl = 2.0, AP_{max} = 0.4, AP_{min} = 0.01, FAR = 0.4$	
Conventional CS	fl = 2.0, AP = 0.1	
HS	HMCR = 0.9, PAR = 0.1, bw = 0.03	
DE	$PC_r = 0.5, F = 0.2$	
GO	$C_{max} = 1.0, C_{min} = 0.00001$	
SS	Non – parameters	
GA	$P_m = 0.005, P_c = 0.9$	

Figure 8 is a graph representing the convergence of each algorithm, and the red line is the result of the ACS algorithm. In all of the benchmark functions except for five (f14, f20, f21, f22, and f23), it can be seen that the ACS algorithm finds the value closest to the **Min** the fastest. Table 6 presents the analysis results of each algorithm, and the last row shows the ranking using the BF of each algorithm. The ACS algorithm has the best convergence performance on unimodal (f1-f7) and multimodal (f8-f13) functions. In the fixed-dimension multimodal function (f14-f23), f15-f19 confirmed the best convergence performance, but f14 and f20-f23 did not. However, the ACS algorithm showed better convergence performance than the conventional CS algorithm. As a result of using the rankings of BF and MF, the ACS algorithm was derived as 1.65 and 1.78, confirming that it was the best. Therefore, it can be seen that the ACS algorithm has improved exploitation and exploration capabilities compared to the conventional CS algorithm.



Fitness

 -2×10

-1×

-3×10

-4×10

Fitness





Figure 8. Comparison results of the benchmark function: (a) *f*1; (b) *f*2; (c) *f*3; (d) *f*4; (e) *f*5; (f) *f*6; (g) *f*7; (h) *f*8; (i) *f*9; (j) *f*10; (k) *f*11; (l) *f*12; (m) *f*13; (n) *f*14; (o) *f*15; (p) *f*16; (q) *f*17; (r) *f*18; (s) *f*19; (t) *f*20; (u) *f*21; (v) *f*22; (w) *f*23.

Table 6. Comparison results with other algorithms using the benchmark function.

					Algorithm			
Fun	Index	ACS	Conventional CS	HS	DE	GO	SS	GA
	BF	1.648×10^{-13}	$3.416 imes 10^{-5}$	2.280×10^{-3}	$5.594 imes10^{0}$	1.361×10^{-3}	$4.905 imes 10^{-9}$	$4.447 imes 10^{-6}$
f1	MF	$9.166 imes 10^{-12}$	$1.054 imes10^{-4}$	$1.067 imes10^{-1}$	$9.928 imes10^1$	$1.799 imes10^{-1}$	$7.454 imes10^{-9}$	$1.038 imes10^{-4}$
	Std	$1.281 imes10^{-11}$	$6.157 imes10^{-5}$	$1.134 imes10^{-1}$	$9.165 imes 10^1$	$3.703 imes10^{-1}$	$1.465 imes10^{-9}$	$1.437 imes10^{-4}$
	BF	$6.511 imes 10^{-7}$	$5.118 imes10^{-1}$	1.274×10^{-2}	$3.360 imes10^{-5}$	$3.545 imes10^{-1}$	2.477×10^{-3}	$2.827 imes10^{-5}$
f2	MF	$6.843 imes10^{-5}$	$1.493 imes10^{0}$	1.764×10^{-2}	$2.358 imes10^{-1}$	$5.134 imes10^{0}$	$1.053 imes 10^0$	$1.537 imes10^{-4}$
	Std	$7.103 imes 10^{-5}$	$6.184 imes10^{-1}$	2.859×10^{-3}	$2.938 imes10^{-1}$	$7.628 imes 10^0$	$1.256 imes 10^0$	$1.744 imes10^{-4}$
	BF	$7.314 imes10^{-9}$	$4.392 imes 10^{0}$	1.081×10^3	3.612×10^2	5.409×10^2	$7.628 imes10^{-1}$	6.919×10^2
f3	MF	$2.980 imes 10^{-6}$	$9.869 imes 10^0$	3.426×10^3	1.376×10^{3}	1.340×10^{3}	$4.734 imes10^{0}$	2.276×10^{3}
	Std	$4.618 imes10^{-6}$	$4.004 imes 10^{0}$	1.137×10^{3}	8.058×10^2	9.813×10^2	$4.891 imes 10^0$	9.368×10^{2}
	BF	$3.277 imes 10^{-7}$	$7.622 imes 10^{-1}$	$2.729 imes 10^0$	$1.391 imes 10^1$	$3.172 imes 10^0$	$1.229 imes 10^0$	$1.608 imes 10^1$
f4	MF	7.729×10^{-5}	$2.682 imes 10^0$	$3.733 imes 10^0$	$2.569 imes 10^1$	$7.949 imes 10^0$	$4.523 imes 10^0$	$2.468 imes 10^1$
	Std	$1.092 imes 10^{-4}$	$1.112 imes 10^0$	$6.870 imes10^{-1}$	$6.344 imes 10^0$	$3.152 imes 10^0$	$2.623 imes 10^0$	$5.934 imes10^{0}$
	BF	$6.970 imes10^{-8}$	$2.489 imes 10^1$	$2.604 imes 10^1$	2.021×10^3	$2.778 imes 10^1$	2.231×10^1	$3.549 imes10^{0}$
f5	MF	$8.623 imes 10^0$	$6.917 imes10^1$	1.254×10^2	$5.880 imes 10^4$	2.734×10^2	1.700×10^2	3.619×10^2
	Std	$1.240 imes 10^1$	$6.094 imes 10^1$	$6.013 imes 10^1$	$8.939 imes 10^4$	4.733×10^2	3.320×10^2	6.670×10^2

Table 6. Cont.

					Algorithm			
Fun	Index	ACS	Conventional CS	HS	DE	GO	SS	GA
	BF	1.945×10^{-12}	$2.620 imes 10^{-5}$	$7.714 imes 10^{-4}$	$3.110 imes 10^{-1}$	1.644×10^{-3}	$4.494 imes 10^{-9}$	$6.886 imes 10^{-6}$
<i>f</i> 6	MF	$5.110 imes10^{-11}$	$1.157 imes10^{-4}$	$1.028 imes 10^{-1}$	1.757×10^2	$2.023 imes10^{-1}$	$7.604 imes10^{-9}$	$1.080 imes10^{-4}$
	Std	$4.041 imes10^{-11}$	$6.755 imes 10^{-5}$	$1.153 imes10^{-1}$	2.134×10^2	$3.159 imes10^{-1}$	$1.469 imes10^{-9}$	$1.668 imes10^{-4}$
	BF	$2.558 imes10^{-5}$	6.412×10^{-3}	1.142×10^{-2}	7.737×10^{-3}	4.156×10^{-3}	1.382×10^{-2}	2.693×10^{-2}
<i>f</i> 7	MF	$5.264 imes10^{-4}$	2.555×10^{-2}	3.257×10^{-2}	5.536×10^{-2}	9.534×10^{-3}	$4.395 imes 10^{-2}$	9.956×10^{-2}
	Std	$3.602 imes 10^{-4}$	1.058×10^{-2}	1.276×10^{-2}	7.461×10^{-2}	4.320×10^{-3}	1.765×10^{-2}	4.806×10^{-2}
	BF	$-1.134 imes 10^4$	-1.012×10^{4}	-1.049×10^{4}	-1.005×10^{4}	-8.594×10^{3}	-8.758×10^{3}	-1.872×10^{3}
f8	MF	-6.895×10^{3}	-7.660×10^{3}	-9.614×10^{3}	-9.172×10^{3}	-7.288×10^{3}	-7.467×10^{3}	-1.803×10^{3}
	Std	1.848×10^{3}	1.311×10^{3}	3.569×10^2	4.755×10^{2}	6.386×10^2	7.267×10^2	3.793×10^{1}
4-	BF	6.928×10^{-14}	1.393×10^{1}	3.545×10^{-3}	$2.327 \times 10^{\circ}$	4.287×10^{1}	2.786×10^{1}	1.095×10^{1}
<i>f</i> 9	MF	8.291×10^{-1}	2.852×10^{1}	5.622×10^{-2}	9.072×10^{0}	8.816×10^{1}	6.106×10^{1}	1.862×10^{1}
	Std	4.541×10^{0}	1.209×10^{1}	1.904×10^{-1}	$3.387 \times 10^{\circ}$	3.199×10^{1}	1.822×10^{1}	$5.276 \times 10^{\circ}$
61 0	BF	8.848×10^{-8}	$2.661 \times 10^{\circ}$	8.655×10^{-3}	6.096×10^{-1}	$1.905 \times 10^{\circ}$	2.066×10^{-5}	4.441×10^{0}
<i>f</i> 10	MF	4.896×10^{-7}	3.946×10^{0}	1.314×10^{-1}	$2.219 \times 10^{\circ}$	$3.729 \times 10^{\circ}$	2.011×10^{0}	1.948×10^{1}
	Std	3.530×10^{-7}	7.918×10^{-1}	1.848×10^{-1}	$1.328 \times 10^{\circ}$	$1.031 \times 10^{\circ}$	9.248×10^{-1}	$2.841 \times 10^{\circ}$
(11	BF	1.682×10^{-11}	1.314×10^{-3}	9.000×10^{-1}	1.131×10^{-1}	1.943×10^{-1}	1.349×10^{-8}	1.147×10^{-7}
<i>f</i> 11	MF	2.466×10^{-4}	1.826×10^{-2}	$1.022 \times 10^{\circ}$	$1.653 \times 10^{\circ}$	4.813×10^{-1}	6.895×10^{-3}	2.937×10^{-2}
	Std	1.350×10^{-3}	1.558×10^{-2}	3.036×10^{-2}	$1.292 \times 10^{\circ}$	1.694×10^{-1}	8.483×10^{-3}	2.875×10^{-2}
(10	BF	1.718×10^{-8}	$5.955 \times 10^{\circ}$	2.056×10^{-4}	5.004×10^{-1}	$1.786 \times 10^{\circ}$	5.058×10^{-1}	9.760×10^{-7}
<i>f</i> 12	MF	8.806×10^{-1}	1.012×10^{1}	5.314×10^{-3}	1.895×10^4	$5.242 \times 10^{\circ}$	3.970×10^{0}	3.877×10^{-2}
	Std	1.538×10^{6}	3.799×10^{6}	1.873×10^{-2}	4.376×10^{4}	$2.275 \times 10^{\circ}$	3.220×10^{6}	7.596×10^{-2}
(10	BF	1.050×10^{-10}	5.023×10^{-4}	9.494×10^{-3}	4.811×10^{6}	8.914×10^{-1}	3.471×10^{-10}	1.272×10^{-3}
<i>f</i> 13	MF	1.099×10^{-3}	1.440×10^{0}	7.514×10^{-2}	5.142×10^{4}	1.427×10^{1}	5.096×10^{-3}	8.399×10^{-3}
	Sta	3.353×10^{-9}	7.727×10^{-1}	4.988×10^{-1}	8.230×10^{-1}	1.421×10^{-1}	9.396×10^{-1}	$1.2/1 \times 10^{-1}$
£14	BF ME	9.980×10^{-1}	9.980×10^{-1}	9.980×10^{-1}	9.980×10^{-1}	9.980×10^{-1}	9.980×10^{-1}	9.980×10^{-1}
J 14	MF Ctd	$1.593 \times 10^{\circ}$	$1.741 \times 10^{\circ}$	$1.525 \times 10^{\circ}$ 1 504 × 10 ⁰	$3.259 \times 10^{\circ}$	9.980×10^{-16}	9.980×10^{-16}	9.980×10^{-16}
	BE	0.009×10^{-4}	0.320×10^{-4}	1.304×10^{-4}	5.227×10^{-4}	5.557×10^{-4}	1.912×10^{-10}	0.403×10^{-4}
<i>f</i> 15		3.073×10^{-3}	5.075×10^{-4}	7.034×10^{-2}	4.929×10^{-2}	4.401×10^{-2}	5.430×10^{-3}	5.437×10^{-3}
<i>J</i> 15	Std	2.712×10^{-3}	0.970×10^{-4}	1.042×10^{-3}	1.401×10 2.585 $\times 10^{-2}$	2.342×10 3.703×10^{-2}	1.467×10 3.573×10^{-3}	4.471×10^{-3} 7 153 × 10 ⁻³
	BE	1.032×10^{0}	4.103×10^{-1}	9.107×10 1.032×10^{0}	$2.363 \times 10^{-1.032} \times 10^{-0.000}$	1.032×10^{0}	1.032×10^{0}	1.032×10^{0}
£16	MF	-1.032×10^{-1}	-1.032×10^{-1}	-1.032×10^{-1}	$-1.032 \times 10^{-1.032}$	$-1.032 \times 10^{-1.032}$	$-1.032 \times 10^{-1.032}$	-1.032×10^{-1}
<i>j</i> 10	Std	6.649×10^{-16}	6.775×10^{-16}	5.616×10^{-6}	6.674×10^{-16}	1.002×10^{-15} 1.422×10^{-15}	3.227×10^{-15}	1.002×10^{-11} 1.763 × 10 ⁻¹¹
	BF	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}
f17	MF	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}	3.979×10^{-1}
) =:	Std	0.000×10^{0}	0.000×10^{0}	5.302×10^{-6}	2.739×10^{-10}	2.699×10^{-15}	1.174×10^{-15}	2.203×10^{-10}
	BF	3.000×10^{0}	3.000×10^{0}	3.000×10^{0}	3.000×10^{0}	3.000×10^{0}	3.000×10^{0}	3.000×10^{0}
f18	MF	3.000×10^{0}	3.000×10^{0}	8.194×10^{0}	3.030×10^{0}	1.110×10^{1}	3.000×10^{0}	3.900×10^{0}
,	Std	$6.171 imes 10^{-16}$	1.588×10^{-15}	1.021×10^1	1.654×10^{-1}	$2.472 imes 10^1$	2.944×10^{-14}	4.930×10^{0}
	BF	-3.863×10^{0}	-3.863×10^{0}	-3.863×10^{0}	$-3.854 imes10^{0}$	$-3.846 imes10^{0}$	-3.863×10^{0}	-3.863×10^{0}
f19	MF	-3.862×10^{0}	-3.863×10^{0}	$-3.811 imes10^{0}$	-3.172×10^{0}	-2.962×10^{0}	$-3.863 imes10^{0}$	$-3.418 imes10^{0}$
-	Std	$4.462 imes 10^{-3}$	$1.766 imes10^{-6}$	$1.961 imes 10^{-1}$	$7.050 imes10^{-1}$	$7.169 imes10^{-1}$	$1.685 imes10^{-9}$	$7.426 imes 10^{-1}$
	BF	-3.322×10^{0}	-3.322×10^{0}	-3.322×10^{0}	$-3.142 imes 10^0$	$-3.094 imes10^{0}$	$-3.322 imes 10^0$	$-3.317 imes10^{0}$
f20	MF	-3.273×10^{0}	-3.275×10^{0}	$-3.286 imes 10^{0}$	$-2.146 imes 10^0$	$-1.844 imes 10^0$	$-3.200 imes 10^0$	$-2.774 imes10^{0}$
	Std	8.411×10^{-2}	6.329×10^{-2}	5.541×10^{-2}	$7.468 imes10^{-1}$	$8.682 imes10^{-1}$	2.442×10^{-2}	$6.166 imes10^{-1}$
	BF	-9.999×10^{0}	$-9.938 imes10^{0}$	$-1.015 imes10^1$	-9.006×10^{0}	$-1.015 imes10^1$	$-1.015 imes10^1$	$-1.015 imes10^1$
f21	MF	-9.528×10^{0}	$-8.965 imes 10^{0}$	-5.059×10^{0}	$-2.017 imes10^{0}$	$-4.425 imes 10^0$	$-6.713 imes 10^{0}$	-4.262×10^{0}
	Std	$7.279 imes10^{-1}$	$9.988 imes10^{-1}$	$3.442 imes 10^0$	$1.845 imes 10^{0}$	$3.578 imes 10^{0}$	$3.194 imes 10^{0}$	$2.889 imes 10^{0}$
	BF	$-1.000 imes 10^1$	-9.671×10^{0}	$-1.040 imes 10^{1}$	$-8.953 imes 10^{0}$	$-1.040 imes10^1$	$-1.040 imes10^1$	-1.015×10^{1}
f22	MF	-9.108×10^{0}	-8.100×10^{0}	-5.100×10^{0}	-2.922×10^{0}	-4.303×10^{0}	-8.646×10^{0}	-4.803×10^{0}
	Std	$1.946 imes 10^0$	$1.258 imes 10^0$	$3.054 imes 10^0$	1.835×10^{0}	2.878×10^{0}	3.042×10^{0}	$3.071 imes 10^0$
-	BF	$-9.999 imes 10^{0}$	-9.236×10^{0}	-1.054×10^{1}	-1.053×10^{1}	-1.054×10^{1}	-1.054×10^{1}	-1.015×10^{1}
f23	MF	-9.850×10^{0}	-6.796×10^{0}	-6.058×10^{0}	-5.512×10^{0}	-4.306×10^{0}	-8.093×10^{0}	-5.809×10^{0}
	Std	$2.426 imes 10^{-1}$	1.460×10^{0}	3.753×10^{0}	3.357×10^{0}	2.964×10^{0}	3.121×10^{0}	3.636×10^{0}
	BF	1.65	3.70	3.83	4.74	4.43	2.70	3.30
Kanking	MF	1.78	3.39	3.52	5.57	5.22	2.96	3.96

4.2. Engineering Problems

Table 7 is a parameter of the ACS algorithm used to solve the numerical problem. The engineering problem was repeatedly interpreted 20 times. The fitness of the engineering problem was calculated as shown in Equation (12). Here, f(x), P(x), and x indicate a result

value, a penalty value, and a design variable defined in each problem, respectively. P(x) can be defined as in Equation (13). Here, np, p_i represent the number of constraints and a value assigned by the constraint, respectively. If the constraint is met, then p_i is 0, and if the constraint is not met, then a penalty of 10^4 is imposed.

Table 7. Parameters for engineering problems.

Algorithm	Parameters
ACS Conventional CS	$ t_{max} = 200, N = 50, AP_{max} = 0.4, AP_{min} = 0.01, fl = 2.0, FAR = 0.4 t_{max} = 200, N = 50, AP = 0.1, fl = 2.0 $

$$F(x) = f(x) \times P(x)$$
(12)

$$P(x) = (1 + 10 \times \sum_{i=1}^{np} p_i)^2$$
(13)

4.2.1. Pressure Vessel Design (PVD) Problem

This problem posed here is to design a cylindrical container with both ends blocked by a hemispherical head as shown in Figure 9 in such a way as to minimize material, forming, and welding costs. The design variables are T_s (shell thickness; x_1), T_h (head thickness; x_2), R (inner radius; x_3), and L (container length; x_4), and the range that the design variables can have is given by Equation (14). The cost minimization problem for cylindrical containers is expressed as an equation in Equation (15). In addition, each design variable has the constraint presented by Equation (16).



Figure 9. PVD problem.

$$\begin{array}{l} 0.0 \le x_1 \ or \ x_2 \le 99.0 \\ 10.0 \le x_3 \ or \ x_4 \le 200.0 \end{array} \tag{14}$$

$$\min f(x) = 0.6224x_1x_3x_4 + 1.7781x_2x_3^2 + 3.1661x_1^2x_4 + 19.84x_1^2x_3 \tag{15}$$

Subject to
$$:g_1(x) = -x_1 + 0.0193x_3 \le 0$$

 $g_2(x) = -x_2 + 0.00954x_3 \le 0$
 $g_3(x) = -\pi x_3^2 x_4 - \frac{4}{3}\pi x_3^3 + 1,296,000 \le 0$
 $g_4(x) = -x_4 + 240 \le 0$
(16)

Table 8 compares the results of the ACS algorithm with those of previous studies [24–27]. The ACS algorithm derived the smallest cost of 5885.333 (design variables were 0.7782, 0.3846, 40.3196, 200.0), and all of the constraints were satisfied. The ACS algorithm reduced the cost by about 0.08% compared to the conventional CS algorithm and by 6.85% compared to Coello's results.

Design				Sandgren	This Paper	
Variables	Coeno [24]	Deb [25]	Kramer [26]	[27]	Conventional CS	ACS
<i>x</i> ₁	0.8125	0.9375	1.125	1.125	0.7783	0.7782
x_2	0.4375	0.5000	0.625	0.625	0.3860	0.3846
<i>x</i> ₃	40.3239	48.3290	58.291	47.700	40.3211	40.3196
x_4	2000.0000	112.6790	43.690	117.701	200.0000	200.0000
$g_1(x)$	-0.0343	-0.0048	0.0000	-0.2044	$-7.2388 imes 10^{-5}$	$-3.8296 imes 10^{-9}$
$g_2(x)$	-0.0528	-0.0389	-0.0689	-0.1699	-0.014	$-6.4738 imes 10^{-9}$
$g_3(x)$	-27.1058	-3652.8768	-21.2201	54.2260	-1.0402×10^{2}	$-3.4597 imes 10^{-4}$
$g_4(x)$	-40.0000	-127.3210	-196.3100	-122.2990	-40.0000	-40.0000
F(x)	6288.7445	6410.3811	7198.0428	8,129.1036	5890.288	5885.333

Table 8. Results of PVD problem.

4.2.2. Welded Beam Design Problem

This problem posed here is to minimize the costs of welding and materials for the welding of two beams, as shown in Figure 10. *h* (welding height; x_1), *l* (welding length; x_2), *t* (thickness of beam 2; x_3), and *b* (width of beam 2; x_4) are design variables, and the range that the design variables can have is given by Equation (17). The welding cost minimization problem is expressed as an equation in Equation (18). Here, the load (*P*) applied to Beam 2 is 6000 lb, the length (*L*) of Beam 2 is 14.0 inches, the modulus of elasticity (*E*) is 30×10^6 psi, the modulus of shear elasticity (*G*) is 12×10^6 psi, the maximum shear stress (τ_{max}) is 13,600 psi, maximum stress (σ_{max}) is 30,000 psi, and the maximum displacement (δ_{max}) is 0.25 inches. In addition, each design variable has the constraints provided by Equation (19).



Figure 10. Welded beam design problem.

$$0.1 \le x_1 \text{ or } x_4 \le 2.0$$

$$0.1 \le x_2 \text{ or } x_3 \le 10.0$$
(17)

$$\min f(x) = 1.10471x_1^2x_2 + 0.04811x_3x_4(14.0 + x_2)$$
(18)

Subject to
$$:g_1(x) = \tau(x) - \tau_{max} \le 0$$

 $g_2(x) = \sigma(x) - \sigma_{max} \le 0$
 $g_3(x) = x_1 - x_4 \le 0$
 $g_4(x) = 0.10471x_1^2 + 0.04811x_3x_4(14.0 + x_2) - 5.0 \le 0$ (19)
 $g_5(x) = 0.125 - x_1 \le 0$
 $g_6(x) = \delta(x) - \delta_{max} \le 0$
 $g_7(x) = P - P_c(x) \le 0$

Table 9 compares the results of the ACS algorithm with those of previous studies [24,28–30]. The ACS algorithm derived the smallest cost of 1.7254 (design variables were 0.2057, 3.4747, 9.0365, and 0.2057), and all of the constraints were satisfied. The ACS algorithm reduced the cost by about 0.23% compared to the conventional CS algorithm and by 1.33% compared to a study by Coello [24].

Design	Caella [24]	Dah [29]	6:44-11 [20]	Ragsdell &	This	Paper
Variables		Deb [28]	Siddaii [29]	Phillips [30]	Conventional CS	ACS
<i>x</i> ₁	0.2088	0.2489	0.2444	0.2455	0.2060	0.2057
x_2	3.4205	6.1730	6.2189	6.1960	3.4724	3.4747
<i>x</i> ₃	8.9975	8.1789	8.2915	8.2730	9.0174	9.0365
x_4	0.2100	-0.2533	0.2444	0.2455	0.2067	0.2057
$g_1(x)$	-0.3378	-5758.6038	-5743.5020	-5743.8265	$-1.2036 imes10^1$	$-3.871 imes 10^{-1}$
$g_2(x)$	-353.9026	-255.5769	-4.0152	-4.7151	$-1.05166 imes 10^{1}$	$-6.9459 imes 10^{-9}$
$g_3(x)$	-0.0012	-0.0044	0.0000	0.0000	$-6.8759 imes 10^{-4}$	$-2.1207 imes 10^{-5}$
$g_4(x)$	-3.4119	-2.9829	-3.0226	-3.0203	-3.4289	-3.4326
$g_5(x)$	-0.0838	-0.1239	-0.1194	-0.1205	-0.0810	-0.0807
$g_6(x)$	-0.2356	-0.2342	-0.2342	-0.2342	-0.2355	-0.2355
$g_7(x)$	-363.2324	-4465.2709	-3490.4694	-3604.2750	-75.0733	-0.4201
F(x)	1.7483	2.4331	2.3815	2.3859	1.7294	1.7254

Table 9. Results of Welded beam design problem.

4.2.3. Weight of a Tension/Compression Spring Problem

This problem presented here is to minimize the weight of a spring that satisfies the constraints when a load is applied to said spring, as shown in Figure 11. The design variables are d (spring thickness; x_1), D (spring diameter; x_2), and N (spring coil count; x_3), and the range that the design variables can have is given by Equation (20). The spring weight minimization problem is expressed as an equation by Equation (21). In addition, each design variable has the constraints provided by Equation (22).

0 0**-** /



Figure 11. Weight of a tension/compression spring problem.

$$\begin{array}{l}
0.05 \le x_1 \le 2.00 \\
0.25 \le x_2 \le 1.30 \\
2.00 \le x_3 \le 15.0
\end{array}$$
(20)

$$min f(x) = (N+2)Dd^2$$
 (21)

Subject to
$$:g_1(x) = 1 - \frac{D^3 N}{71,785d^4} \le 0$$

 $g_2(x) = \frac{4D^2 - dD}{12,566(Dd^3 - d^4)} + \frac{1}{5108d^2} - 1 \le 0$
 $g_3(x) = 1 - \frac{140.45d}{D^2 N} \le 0$
 $g_4(x) = \frac{D+d}{1.5} - 1 \le 0$
(22)

Table 10 shows the results of the ACS algorithm and those of other researchers. The ACS algorithm derived the smallest spring weight of 1.2665×10^{-2} (the design variables were 0.0517, 0.3578, and 11.2240), and all of the constraints were satisfied. The ACS algorithm reduced the weight by about 0.03% compared to the conventional CS algorithm and by 0.31% compared to study by Coello [24].

Design			Pologun du [20]	This Paper		
Variables	Afora [51]	belegundu [32]	Conventional CS	ACS		
x_1	0.0515	0.0534	0.0500	0.0520	0.0517	
<i>x</i> ₂	0.3517	0.3992	0.3159	0.3642	0.3578	
<i>x</i> ₃	11.6322	9.1854	14.2500	10.8626	11.2240	
$g_1(x)$	-0.00218	0.00002	-0.00001	$-4.5997 imes 10^{-5}$	$-20183 imes 10^{-11}$	
$g_2(x)$	-0.00011	-0.00002	-0.00378	-7.9229×10^{-5}	$-2.6946 imes 10^{-10}$	
$g_3(x)$	-4.02632	-4.12383	-3.93830	-4.0677	-4.0560	
$g_4(x)$	-4.02632	-0.69828	-0.75607	-0.7225	-0.7270	
F(x)	$1.2704 imes 10^{-2}$	$1.2730 imes 10^{-2}$	$1.2833 imes 10^{-2}$	$1.2669 imes 10^{-2}$	$1.2665 imes 10^{-2}$	

Table 10. Results of weight of a spring problem.

4.2.4. Weight of a Three-Bar Truss Problem

This problem aims to find the minimum truss weight that satisfies the constraints when a load (*P*) is applied to a truss structure of three members, such as in Figure 12. The design variables are A_1 (cross-sectional area of Member 1; $x_1 = x_3$) and A_2 (cross-sectional area of Member 2; x_2), and the range that the design variables can have is given by Equation (23). The three-bar truss weight minimization problem is expressed as an equation by Equation (24). Here, the distance (*L*) of the node is 100 cm, the load (*P*) is 2 kN/cm^2 , and the maximum stress (σ_{max}) is 2 kN/cm^2 . In addition, each design variable has the constraints provided in Equation (25). The maximum number of generations (t_{max}) was set at 20 in this problem.



Figure 12. Weight of a three-bar truss problem.

$$0.0 \le x_1 \text{ or } x_2 \le 1.0$$
 (23)

$$\min f(x) = (2\sqrt{2x_1} + x_2)l \tag{24}$$

Subject to
$$:g_1(x) = \frac{\sqrt{2}x_1 + x_2}{\sqrt{2}x_1^2 + 2x_1x_2}P - \sigma \le 0$$

 $g_2(x) = \frac{x_2}{\sqrt{2}x_1^2 + 2x_1x_2}P - \sigma \le 0$
 $g_3(x) = \frac{1}{\sqrt{2}x_2 + x_1}P - \sigma \le 0$
(25)

Table 11 shows the results of the ACS algorithm and those of a previous study [14]. Here, SoC, MB, and DSS-MDE refer to the society and civilization (SoC) algorithm, the mine blast (MB) algorithm, and the dynamic stochastic selection with multimember differential evolution (DSS-MDE) algorithm. The ACS algorithm determined the weight of the three-bar truss structure to be 263.895843 (the design variables were 0.7887 and 0.4081), and all of

the constraints were satisfied. The result of the ACS algorithm was lighter than the results of the conventional CS algorithm and Askarzadeh.

Table 11. Results o	f weight of a three-	bar truss problem.
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Design	Askarzadeh [14]			This Paper		
Variables	SoC	MB	DSS-MDE	Conventional CS	ACS	
<i>x</i> ₁	-	-	-	0.7887	0.7887	
<i>x</i> ₂	-	-	-	0.4081	0.4081	
$g_1(x)$	-	-	-	$-1.7977 imes 10^{-9}$	$-3.9746 imes 10^{-14}$	
$g_2(x)$	-	-	-	-1.4642	-1.4643	
$g_3(x)$	-	-	-	-0.5358	-0.5357	
F(x)	263.895846	263.895852	263.895849	263.895844	263.895843	

4.2.5. Stepped Cantilever Beam Design Problem

The problem posed here is to calculate the width of a stepped cantilever beam as shown in Figure 13 and minimize its weight. The λ_{1-5} (width) of the five-cantilever beam is a design variable (x_{1-5}), and the range that the design variable can have is provided by Equation (26). Equation (27) is an expression of the stepped cantilever beam design problem, and Equation (28) is a constraint of the stepped cantilever beam design problem.



Figure 13. Weight of a three-bar truss problem.

$$0.01 \le x_1 \text{ or } x_2 \text{ or } x_3 \text{ or } x_4 \text{ or } x_5 \le 100.0 \tag{26}$$

$$\min f(x) = 0.0624 \sum_{i=1}^{5} x_i \tag{27}$$

Subject to
$$:g_1(x) = \frac{61}{x_1^3} + \frac{37}{x_2^3} + \frac{19}{x_3^3} + \frac{7}{x_4^3} + \frac{1}{x_5^3} - 1 \le 0$$
 (28)

Table 12 shows the results of the ACS algorithm and those of Hijjawi et al. [33]. Here, AOACS and HHO stand for the hybrid algorithm of the arithmetical optimization algorithm and cuckoo search and for Harris hawks optimization, respectively. The ACS algorithm determined a minimum weight of the step cantilever beam of 1.3418 and satisfied the constraints. The ACS algorithm showed better results than the conventional CS algorithm.

Design	Hijjawi et al. [33]			This Paper		
Variables	AOACS	ННО	PSO	Conventional CS	ACS	
<i>x</i> ₁	6.01	5.13	6.05	7.1816	6.0064	
x_2	5.31	5.62	5.26	4.5530	5.3169	
x_3	4.49	5.10	4.51	4.5403	4.3240	
x_4	3.50	3.93	3.46	3.3118	3.6624	
x ₅	2.15	2.32	2.19	2.7603	2.1929	
$g_1(x)$	0.0019	-0.0011	0.0010	0.0000	0.0000	
F(x)	1.34	1.38	1.34	1.3944	1.3418	

Table 12. Results of the stepped cantilever beam design problem.

5. Conclusions

This paper proposed the ACS algorithm, which improves Step 3 of the conventional CS algorithm. The ACS algorithm added three methods to the conventional CS algorithm. First, unlike conventional CS algorithms that use fixed-value *AP*, we proposed the use of dynamic *AP*, which decreases nonlinearly with the number of generations. This change improved the algorithm's exploration performance. Second, we proposed an expression that follows the crow in the best position rather than following a randomly adopted crow, and this improved the algorithm's exploitation performance. Third, we proposed a local search based on the adopted value rather than a global search of the entire area at later generations. The convergence performance according to the value change of *AP_{max}* and *FAR*—parameters added to the ACS algorithm—was compared, and it was verified that the convergence performance was the best when the *AP_{max}* was in the range of 0.4–0.6 and the *FAR* was in the range of 0.2–0.4. Finally, the ACS algorithm was applied to benchmark functions and four engineering problems in order to confirm that the convergence speed was the fastest and the best convergence performance compared to the results of other metaheuristic algorithms.

In future work, if the ACS algorithm is applied to various large-scale or real-scale engineering problems, it is believed that the optimal solutions for a variety of engineering problems would be obtained.

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