

Article

Quantifying Redundant Information in Predicting a Target Random Variable

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Abstract: We consider the problem of defining a measure of redundant information that quantifies how much common information two or more random variables specify about a target random variable. We discussed desired properties of such a measure, and propose new measures with some desirable properties.

Keywords: synergy; information theory; complex systems; irreducibility; synergistic information; intersection-information

1. Introduction

Many molecular and neurological systems involve multiple interacting factors affecting an outcome synergistically and/or redundantly. Attempts to shed light on issues such as population coding in neurons, or genetic contribution to a phenotype (e.g., eye-color), have motivated various proposals to leverage principled information-theoretic measures for quantifying informational synergy and redundancy, e.g., [1–5]. In these settings, we are concerned with the statistics of how two (or more) random variables X_1, X_2 , called predictors, jointly or separately specify/predict another random variable Y , called a target random variable. This focus on a target random variable is in contrast to Shannon’s mutual information which quantifies statistical dependence between two random variables, and various notions of common information, e.g., [6–8].

The concepts of synergy and redundancy are based on several intuitive notions, e.g., positive informational synergy indicates that X_1 and X_2 act cooperatively or antagonistically to influence Y ; positive redundancy indicates there is an aspect of Y that X_1 and X_2 can each separately predict. However, it has been challenging [9–12] to come up with precise information-theoretic definitions of synergy and redundancy that are consistent with all intuitively desired properties.

2. Background: Partial Information Decomposition

Partial Information Decomposition (PID) [13] defines the concepts of synergistic, redundant and unique information in terms of *intersection information*, $I_{\cap}(\{X_1, \dots, X_n\}:Y)$, which quantifies the common information that each of the n predictors X_1, \dots, X_n conveys about a target random variable Y . An antichain lattice [14] of redundant, unique, and synergistic partial informations is built from the intersection information.

Partial information diagrams (PI-diagrams) extend Venn diagrams to represent synergy. A PI-diagram is composed of nonnegative *partial information regions* (PI-regions). Unlike the standard Venn entropy diagram in which the sum of all regions is the joint entropy $H(X_{1\dots n}, Y)$, in PI-diagrams the sum of all regions (i.e. the space of the PI-diagram) is the mutual information $I(X_{1\dots n}:Y)$. PI-diagrams show how the mutual information $I(X_{1\dots n}:Y)$ is distributed across subsets of the predictors. For example, in the PI-diagram for $n = 2$ (Figure 1): $\{1\}$ denotes the unique information about Y that only X_1 carries (likewise $\{2\}$ denotes the information only X_2 carries); $\{1, 2\}$ denotes the redundant information about Y that X_1 as well as X_2 carries, while $\{12\}$ denotes the information about Y that is specified only by X_1 and X_2 synergistically or jointly.

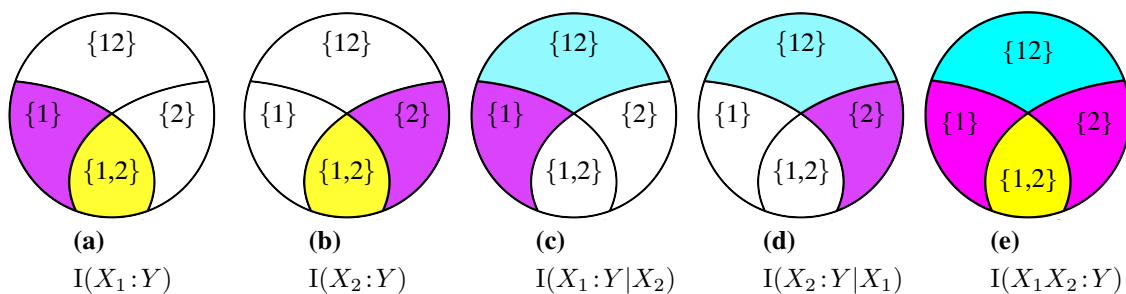


Figure 1. PI-diagrams for $n = 2$ predictors, showing the amount of redundant (yellow/bottom), unique (magenta/left and right) and synergistic (cyan/top) information with respect to the target Y .

Each PI-region is either redundant, unique, or synergistic, but any combination of positive PI-regions may be possible. Per [13], for two predictors, the four partial informations are defined as follows: the redundant information as $I_{\cap}(\{X_1, X_2\}:Y)$, the unique informations as

$$\begin{aligned} I_{\partial}(\{X_1\}:Y) &= I(X_1:Y) - I_{\cap}(\{X_1, X_2\}:Y) \\ I_{\partial}(\{X_2\}:Y) &= I(X_2:Y) - I_{\cap}(\{X_1, X_2\}:Y), \end{aligned} \quad (1)$$

and the synergistic information as

$$\begin{aligned} I_{\partial}(\{X_1, X_2\}:Y) &= I(X_1, X_2:Y) - I_{\partial}(\{X_1\}:Y) - I_{\partial}(\{X_2\}:Y) - I_{\cap}(\{X_1, X_2\}:Y) \\ &= I(X_1, X_2:Y) - I(X_1:Y) - I(X_2:Y) + I_{\cap}(\{X_1, X_2\}:Y). \end{aligned} \quad (2)$$

3. Desired I_\cap properties and canonical examples

There are a number of intuitive properties, proposed in [5,9–13], that are considered desirable for the intersection information measure I_\cap to satisfy:

- (S₀) Weak Symmetry: $I_\cap(\{X_1, \dots, X_n\}:Y)$ is invariant under reordering of X_1, \dots, X_n .
- (M₀) Weak Monotonicity: $I_\cap(\{X_1, \dots, X_n, Z\}:Y) \leq I_\cap(\{X_1, \dots, X_n\}:Y)$ with equality if there exists $X_i \in \{X_1, \dots, X_n\}$ such that $H(Z, X_i) = H(Z)$.

Weak Monotonicity is a natural generalization of the monotonicity property from [13]. Weak monotonicity is inspired by the property of mutual information that if $H(X|Z) = 0$, then $I(X:Y) \leq I(Z:Y)$.

- (SR) Self-Redundancy: $I_\cap(\{X_1\}:Y) = I(X_1:Y)$. The intersection information a single predictor X_1 conveys about the target Y is equal to the mutual information between the X_1 and the target Y .
- (M₁) Strong Monotonicity: $I_\cap(\{X_1, \dots, X_n, Z\}:Y) \leq I_\cap(\{X_1, \dots, X_n\}:Y)$ with equality if there exists $X_i \in \{X_1, \dots, X_n\}$ such that $I(Z, X_i:Y) = I(Z:Y)$.

Strong Monotonicity captures more precisely what is meant by “redundant information”, it says explicitly that it *information about* Y that is redundant, not just any redundancy among the predictors (weak monotonicity).

- (LP) Local Positivity: For all n , the derived “partial informations” defined in [13] are nonnegative. This is equivalent to requiring that I_\cap satisfy *total monotonicity*, a stronger form of supermodularity. For $n = 2$ this can be concretized as, $I_\cap(\{X_1, X_2\}:Y) \geq I(X_1:X_2) - I(X_1:X_2|Y)$.

- (TM) Target Monotonicity: If $H(Y|Z) = 0$, then $I_\cap(\{X_1, \dots, X_n\}:Y) \leq I_\cap(\{X_1, \dots, X_n\}:Z)$.

There are also a number of canonical examples for which one or more of the partial informations have intuitive values, which are considered desirable for the intersection information measure I_\cap to attain.

Example UNQ, shown in Figure 2, is a canonical case of unique information, in which each predictor carries independent information about the target. Y has four equiprobable states: ab, aB, Ab, and AB. X_1 uniquely specifies bit a/A, and X_2 uniquely specifies bit b/B. Note that the states are named so as to highlight the two bits of unique information; it is equivalent to choose any four unique names for the four states.

Example RdnXor, shown in Figure 3, is a canonical example of redundancy and synergy coexisting. The r/R bit is redundant, while the 0/1 bit of Y is synergistically specified as the XOR of the corresponding bits in X_1 and X_2 .

Example And, shown in Figure 4, is an example where the relationship between X_1, X_2 and Y is nonlinear, making the desired partial information values less intuitively obvious. Nevertheless, it is desired that the partial information values should be nonnegative.

Example ImperfectRdn, shown in Figure 5, is an example of “imperfect” or “lossy” correlation between the predictors, where it is intuitively desirable that the derived redundancy should be positive. Given (LP), we can determine the desired decomposition analytically. First, $I(X_1, X_2:Y) =$

$I(X_1:Y) = 1$ bit; therefore, $I(X_2:Y|X_1) = I(X_1, X_2:Y) - I(X_1:Y) = 0$ bits. This determines two of the partial informations—the synergistic information $I_\partial(\{X_1, X_2\}:Y)$ and the unique information $I_\partial(\{X_2\}:Y)$ are both zero. Then, the redundant information $I_\partial(\{X_1, X_2\}:Y) = I(X_2:Y) - I_\partial(\{X_2\}:Y) = I(X_2:Y) = 0.99$ bits. Having determined three of the partial informations, we compute the final unique information $I_\partial(\{X_1\}:Y) = I(X_1:Y) - 0.99 = 0.01$ bits.

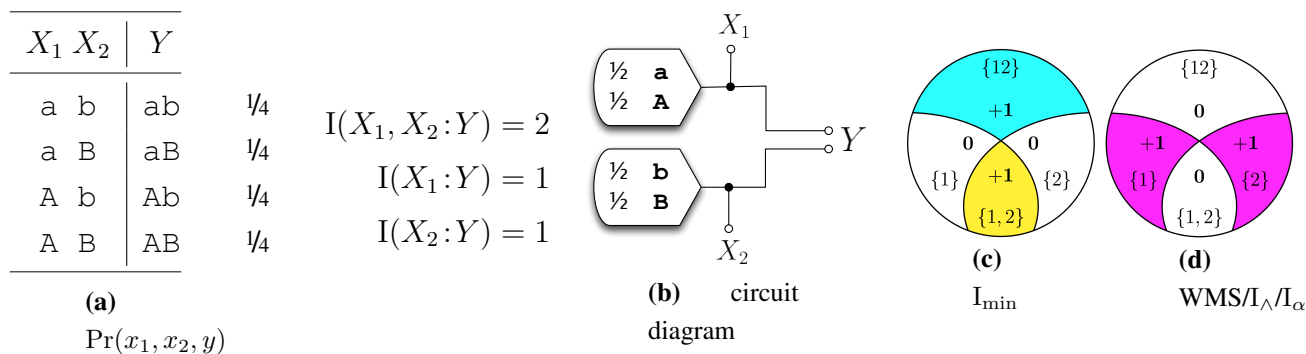


Figure 2. Example UNQ. X_1 and X_2 each uniquely carry one bit of information about Y . $I(X_1 X_2:Y) = H(Y) = 2$ bits.

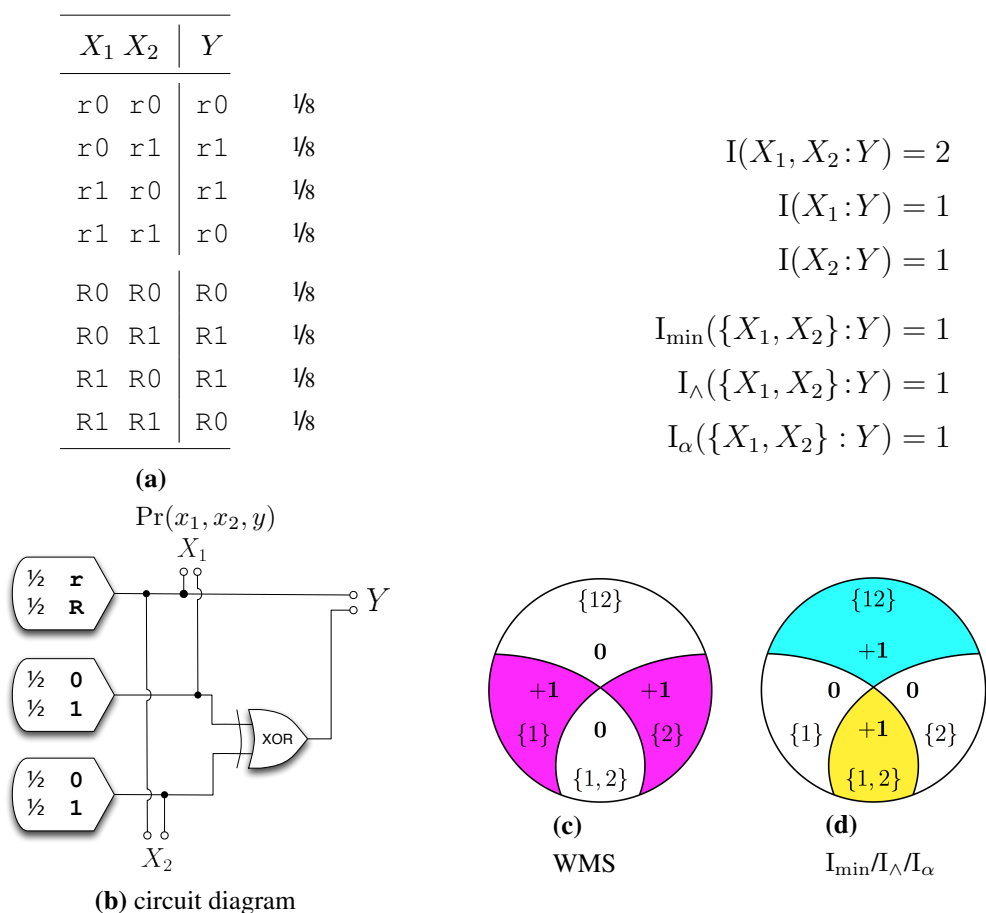


Figure 3. Example RDNXOR. This is the canonical example of redundancy and synergy coexisting. I_{\min} and I_\wedge each reach the desired decomposition of one bit of redundancy and one bit of synergy. This example demonstrates I_\wedge correctly extracting the embedded redundant bit within X_1 and X_2 .

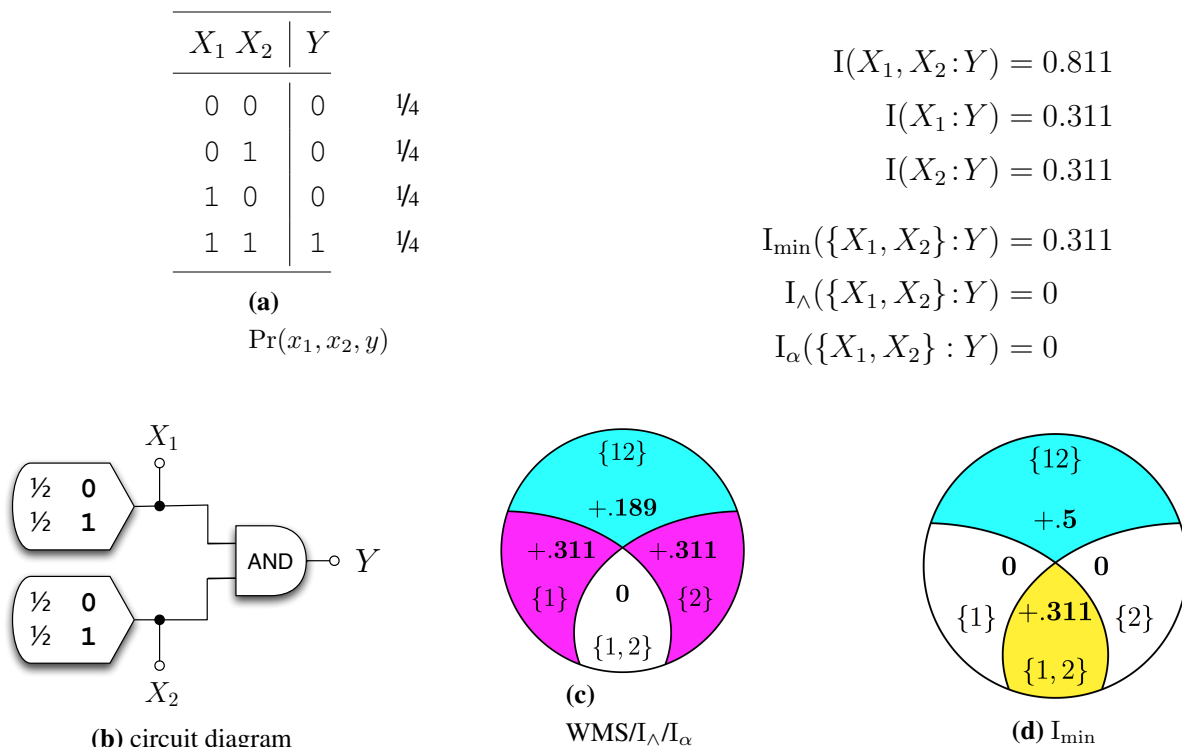


Figure 4. Example AND. It is universally agreed that the redundant information is between $[0, 0.311]$ bits. The most compelling argument is from [15] arguing for 0.311 bits of redundant information.

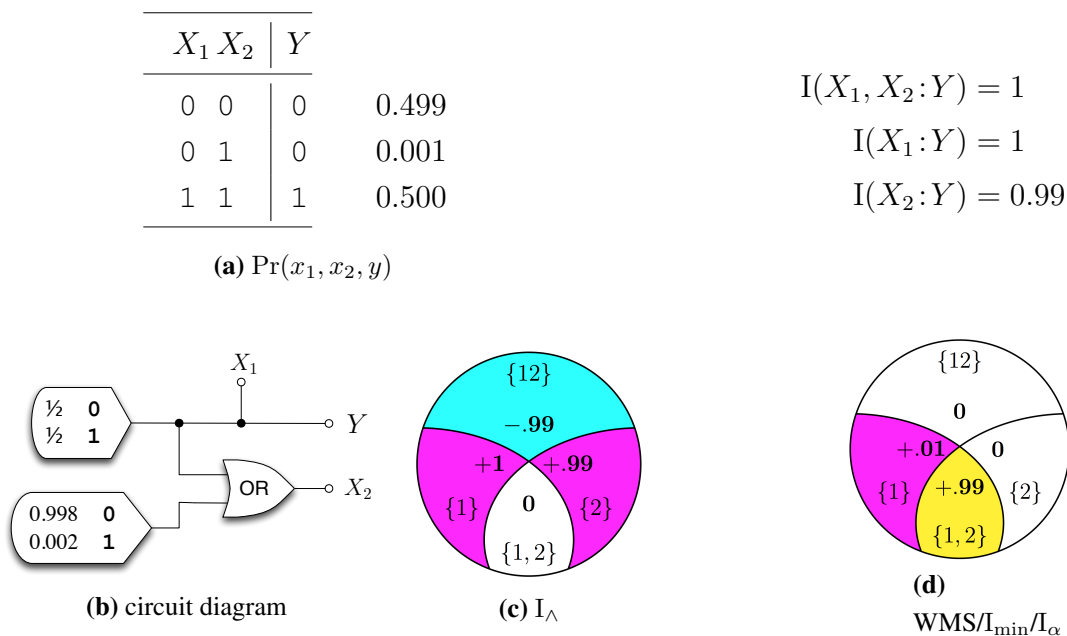


Figure 5. Example IMPERFECTRDN. I_{\wedge} is blind to the noisy correlation between X_1 and X_2 and calculates zero redundant information. An ideal I_{\cap} measure would detect that all of the information X_2 specifies about Y is also specified by X_1 to calculate $I_{\cap}(\{X_1, X_2\} : Y) = 0.99$ bits.

4. Previous candidate measures

In [13], the authors propose to use the following quantity, I_{\min} , as the intersection information measure:

$$\begin{aligned} I_{\min}(X_1, \dots, X_n : Y) &\equiv \sum_{y \in Y} \Pr(y) \min_{i \in \{1, \dots, n\}} I(X_i : Y = y) \\ &= \sum_{y \in Y} \Pr(y) \min_{i \in \{1, \dots, n\}} D_{\text{KL}}[\Pr(X_i | y) \| \Pr(X_i)] , \end{aligned} \quad (3)$$

where D_{KL} is the Kullback-Leibler divergence.

Though I_{\min} is an intuitive and plausible choice for the intersection information, [9] showed that I_{\min} has counterintuitive properties. In particular, I_{\min} calculates one bit of redundant information for example UNQ (Figure 2). It does this because each input shares one bit of information with the output. However, it is quite clear that the shared informations are, in fact, different: X_1 provides the low bit, while X_2 provides the high bit. This led to the conclusion that I_{\min} *overestimates* the ideal intersection information measure by focusing only on *how much* information the inputs provide to the output. Another way to understand why I_{\min} overestimates redundancy in example UNQ is to imagine a hypothetical example where there are exactly two bits of unique information for every state $y \in Y$ and no synergy or redundancy. I_{\min} would calculate the redundancy as the minimum over both predictors which would be $\min[1, 1] = 1$ bit. Therefore I_{\min} would calculate 1 bit of redundancy even though by definition there was no redundancy but merely two bits of unique information.

Another candidate measure of synergy, WholeMinusSum (WMS) [9,16], calculates zero synergy and redundancy for Example RDNXOR, as opposed to the intuitive value of one bit of redundancy and one bit of synergy.

5. New candidate measures

5.1. The I_{\wedge} measure

Based on [17], we can consider a candidate intersection information as the maximum mutual information $I(Q : Y)$ that some random variable Q conveys about Y , subject to Q being a function of each predictor X_1, \dots, X_n . After some algebra, this leads to,

$$\begin{aligned} I_{\wedge}(\{X_1, \dots, X_n\} : Y) &\equiv \max_{\Pr(Q|Y)} I(Q : Y) \\ &\text{subject to } \forall i \in \{1, \dots, n\} : H(Q | X_i) = 0 \end{aligned} , \quad (4)$$

which reduces to a simple expression in [12].

Example IMPERFECTRDN highlights the foremost shortcoming of I_{\wedge} ; I_{\wedge} does not detect “imperfect” or “lossy” correlations between X_1 and X_2 . Instead, I_{\wedge} calculates zero redundant information, that $I_{\wedge}(\{X_1, X_2\} : Y) = 0$ bits. This arises from $\Pr(X_1 = 1, X_2 = 0) > 0$. If this were zero, IMPERFECTRDN reverts to being determined by the properties (SR) and the (M_0) equality condition. Due to the nature of the common random variable, I_{\wedge} only sees the “deterministic” correlations between X_1 and X_2 —add even an iota of noise between X_1 and X_2 and I_{\wedge} plummets to zero. This highlights a

related issue with I_\wedge ; it is not continuous—an arbitrarily small change in the probability distribution can result in a discontinuous jump in the value of I_\wedge .

Despite this, I_\wedge is a useful stepping-stone, it captures what is inarguably redundant information (the common random variable). In addition, unlike earlier measures, I_\wedge satisfies (TM).

5.2. The I_α measure Intuitively, we expect that if Q only specifies redundant information, that conditioning on any predictor X_i would vanquish all of the information Q conveys about Y . We take this intuition to its final conclusion and find it yields a tighter lowerbound on I_\cap than I_\wedge . Moreover, I_α pleasantly reduces to a I_\wedge but loosens the constraint in Equation (4) from $H(Q|X_i) = 0$ to $H(Q|X_i) = H(Q_i|X_i, Y)$:

$$I_\alpha(\{X_1, \dots, X_n\} : Y) \equiv \max_{\Pr(Q|Y)} I(Q : Y) \quad (5)$$

$$\text{subject to } \forall i \in \{1, \dots, n\} : I(Q, X_i : Y) = I(X_i : Y)$$

$$= \max_{\Pr(Q|Y)} I(Q : Y) \quad (6)$$

$$\text{subject to } \forall i \in \{1, \dots, n\} : H(Q|X_i) = H(Q_i|X_i, Y)$$

This measure obtains the desired values for the canonical examples in Section 3. However, its implicit definition makes it more difficult to verify whether or not it satisfies the desired properties in Section 3. Pleasingly, I_α also satisfies (TM). We can also show (See Lemmas 1 and 2 in Appendix A) that

$$0 \leq I_\wedge(\{X_1, \dots, X_n\} : Y) \leq I_\alpha(\{X_1, \dots, X_n\} : Y) \leq I_{\min}(\{X_1, \dots, X_n\} : Y) \quad (7)$$

While I_α satisfies previously defined canonical examples, we have found another example, shown in Figure 6, for which I_\wedge and I_α both calculate negative synergy. This example further complicates Example AND by making the predictors mutually dependent.

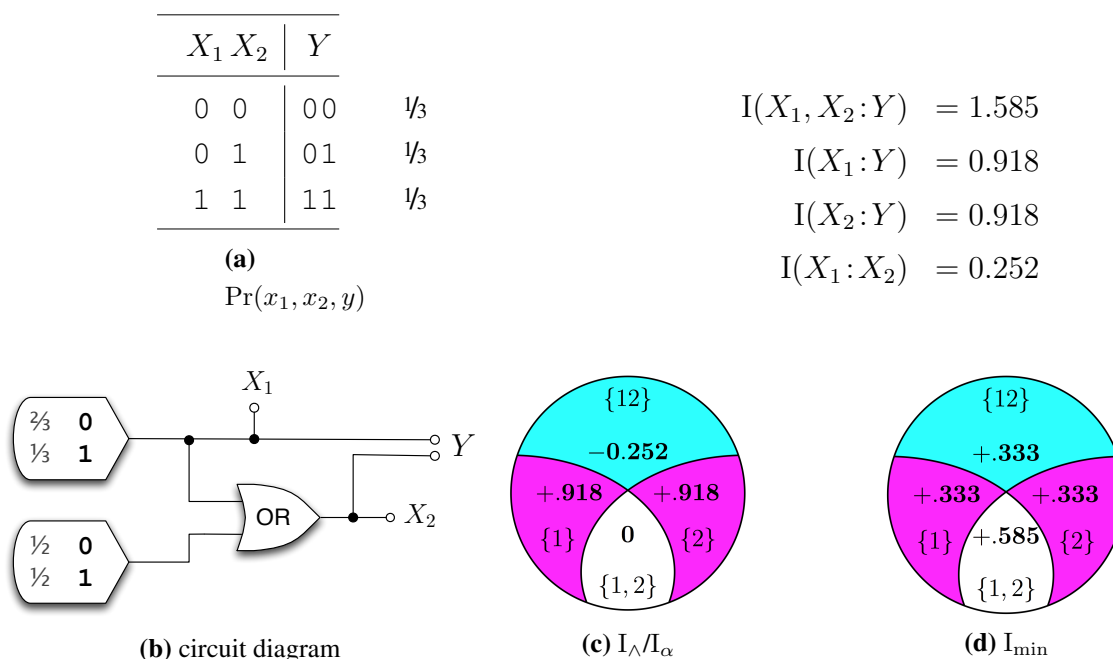


Figure 6. Example SUBTLE.

6. Conclusion

The important part of this paper is exchanging (M_0) with (M_1) thus further constraining the space of acceptable I_\cap measures. The complexity community aspires to eventually find a unique I_\cap measure that satisfies a large portion of the desired properties, and *any noncontroversial tightening of the space of possible I_\cap measures, even (or especially?) if obvious in hindsight, is immensely welcome.*

As discussed in [12], I_\cap measures fail (LP) if and only if they are *too strict* a measure of redundant information. Loosening the constraints on I_\cap yields I_α and achieves a nonnegative decomposition on example IMPERFECTRDN. A natural next step is to loosen the constraints on I_α until achieving a nonnegative decomposition for example SUBTLE. Alternatively, a very plausible measure of the “unique information” [9,15,18] that satisfies (LP) for $n = 2$ yet does not satisfy (TM). It seems that (LP) and (TM) will be incompatible, and it would be nice to prove this.

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Author Contributions

Both authors shared in this research equally. Both authors have read and approved the final manuscript.

Conflicts of Interest

The authors declare no conflict of interest.

A. Appendix

Proof I_α does not satisfy (LP). Proof by counter-example SUBTLE (Figure 6).

For $I(Q:Y|X_1) = 0$, then Q must not distinguish between states of $Y = 00$ and $Y = 01$ (because X_1 does not distinguish between these two states). This entails that $\Pr(Q|Y = 00) = \Pr(Q|Y = 01)$. By symmetry, likewise for $I(Q:Y|X_2) = 0$, Q must be distinguish between states $Y = 01$ and $Y = 11$. Altogether, this entails that $\Pr(Q|Y = 00) = \Pr(Q|Y = 01) = \Pr(Q|Y = 11)$, which then entails, $\Pr(q|y_i) = \Pr(q|y_j) \quad \forall q \in Q, y_i \in Y, y_j \in Y$, which is only achievable when $\Pr(q) = \Pr(q|y) \quad \forall q \in Q, y \in Y$. This makes $I(Q:Y) = 0$, therefore for example SUBTLE, $I_\alpha(\{X_1, X_2\} : Y) = 0$.

Lemma 1. We have $I_\cap(\{X_1, \dots, X_n\} : Y) \leq I_\alpha(\{X_1, \dots, X_n\} : Y)$.

Proof. We define a random variable $Q' = X_1 \wedge \dots \wedge X_n$. We then plugin Q' for Q in the definition of I_α . This newly plugged-in Q satisfies the constraint $\forall i \in \{1, \dots, n\}$ that $I(Q:Y|X_i) = 0$. Therefore, Q' is always a possible choice for Q , and the maximization of $I(Q:Y)$ in I_α must be at least as large as $I(Q':Y) = I_\cap(\{X_1, \dots, X_n\} : Y)$. \square

Lemma 2. We have $I_\alpha(\{X_1, \dots, X_n\} : Y) \leq I_{\min}(X_1, \dots, X_n : Y)$

Proof. For a given state $y \in Y$ and two arbitrary random variables Q and X , given $I(Q:y|X) = D_{\text{KL}}[\text{Pr}(QX|y) \parallel \text{Pr}(Q|X) \text{Pr}(X|y)] = 0$, we show that, $I(Q:y) \leq I(X:y)$,

$$\begin{aligned} I(X:y) - I(Q:y) &= \sum_{x \in X} \text{Pr}(x|y) \log \frac{\text{Pr}(x|y)}{\text{Pr}(x)} - \sum_{q \in Q} \text{Pr}(q|y) \log \frac{\text{Pr}(q|y)}{\text{Pr}(q)} \\ &\geq 0. \end{aligned}$$

Generalizing to n predictors X_1, \dots, X_n , the above shows that that the maximum $I(Q:y)$ under constraint $I(Q:y|X_i)$ will always be less than $\min_{i \in \{1, \dots, n\}} I(X_i:y)$, which completes the proof. \square

Lemma 3. Measure I_{\min} satisfies desired property Strong Monotonicity, (M_1) .

Proof. Given $H(Y|Z) = 0$, then the specific-surprise $I(Z:y)$ yields,

$$\begin{aligned} I(Z:y) &\equiv D_{\text{KL}}[\text{Pr}(Z|y) \parallel \text{Pr}(Z)] \\ &= \sum_{z \in Z} \text{Pr}(z|y) \log \frac{\text{Pr}(z|y)}{\text{Pr}(z)} \\ &= \sum_{z \in Z} \text{Pr}(z|y) \log \frac{1}{\text{Pr}(y)} \\ &= \log \frac{1}{\text{Pr}(y)}. \end{aligned}$$

Given that for an arbitrary random variable X_i , $I(X_i:y) \leq \log \frac{1}{\text{Pr}(y)}$. As I_{\min} takes only uses the $\min_i I(X_i:y)$, the minimum is invariant under adding any predictor Z such that $H(Y|Z) = 0$. Therefore, measure I_{\min} satisfies property (M_1) . \square

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