

Article

A Novel Geo-Broadcast Algorithm for V2V Communications over WSN

José J. Anaya, Edgar Talavera, Felipe Jiménez, Nuria Gómez and José E. Naranjo *

University Institute for Automobile Research (INSIA), Universidad Politécnica de Madrid, Campus Sur UPM, Carretera de Valencia, km. 7, 28031 Madrid, Spain; E-Mails: jjanaya.c@gmail.com (J.J.A.); c_freeman@hotmail.com (E.T.); felipe.jimenez@upm.es (F.J.); nuria.gomez@upm.es (N.G.)

* Author to whom correspondence should be addressed; E-Mail: joseeugenio.naranjo@upm.es; Tel.: +34-91-336-5287; Fax: +34-91-366-7522.

Received: 31 March 2014; in revised form: 8 July 2014 / Accepted: 16 July 2014 /

Published: 28 August 2014

Abstract: The key for enabling the next generation of advanced driver assistance systems (ADAS), the cooperative systems, is the availability of vehicular communication technologies, whose mandatory installation in cars is foreseen in the next few years. The definition of the communications is in the final step of development, with great efforts on standardization and some field operational tests of network devices and applications. However, some inter-vehicular communications issues are not sufficiently developed and are the target of research. One of these challenges is the construction of stable networks based on the position of the nodes of the vehicular network, as well as the broadcast of information destined to nodes concentrated in a specific geographic area without collapsing the network. In this paper, a novel algorithm for geo-broadcast communications is presented, based on the evolution of previous results in vehicular mesh networks using wireless sensor networks with IEEE 802.15.4 technology. This algorithm has been designed and compared with the IEEE 802.11p algorithms, implemented and validated in controlled conditions and tested on real vehicles. The results suggest that the characteristics of the designed broadcast algorithm can improve any vehicular communications architecture to complement a geo-networking functionality that supports a variety of ADAS.

Keywords: intelligent transportation systems; vehicular communications; V2V; cooperative systems; geo-networking

1. Introduction

A great effort is being carried out to increase safety in the traffic environment through inter-vehicle communication. It is expected that vehicle safety applications will rely on communication capabilities for information gathering beyond their immediate environment and line-of-sight vision and for sharing the necessary safety information [1]. In this context, the position of vehicles rather than their identity is relevant, as vehicles share their actual positions and anticipated trajectories, their cooperative merging maneuvers among sideways neighbors [2], instantly notify the vehicle traveling behind of a braking action [3] or warn oncoming traffic [4] of an icy patch [5].

The key to supporting this communication technique is the geolocation and tracking of the vehicles (nodes) that are circulating in the proximity of each node of the network. The core of these systems is based on the combination of two key elements to manage the effective information flow among an apparently random distribution of network nodes: routing algorithms and GPS positioning. The routing algorithms manage the logical structure of the wireless vehicular network, selecting the best route to transmit messages from one or more receivers and dealing with the addition and removal of network nodes. There are two main routing structures considered in a vehicular network, also known as geo-routing. The first is the structure that supports unicast (GeoUnicast) communication; that is, the transmission of data from an emitter to a specific receiver, directly or indirectly, using multiple nodes as intermediaries. The second structure supports broadcast (geo-broadcast) communication; that is, the transmission of data from an emitter to every receiver that is located in a specific geographical region.

However, in order to guarantee a massive impact of this technology on the vehicular environment, the necessary standards need to be developed to guarantee the right interoperability and performance, regardless of the manufacturer. Therefore, a number of working groups within the European Telecommunications Standards Institute (ETSI) initiated the geo-networking specification and family of standards [6], which are currently under deployment to support geo-networking in vehicular environments [7].

Nevertheless, the implementation of this philosophy is presently in an initial stage, where these available developments are based on the IEEE 802.11p (dedicated short range communications (DSRC)) [8] standard, radio modem, WiFi or cellular. Several research projects are being conducted towards the adaptation of IPv6 protocols to support vehicular communications, including geographical routing [9,10], but efficient routing algorithms that guarantee the network performance and stability are still an open research area. An alternative is the application of communication systems oriented to wireless sensor networks (WSN) based on IEEE 802.15.4, where some experience in geo-networking has been reported, but mainly focused towards a static physical infrastructure of the network, with minimum mobility [11–14].

In fact, due to the limited availability of devices that support geo-networking protocols and the inherent difficulties in developing field trials, the usual testing of protocols and hardware has been carried out through simulations [15].

Several research projects are being conducted towards the adaptation of IPv6 protocols to support vehicular communications, including geographical routing [9,10], but efficient routing algorithms that guarantee the network performance and stability are still an open research area. In fact, some of

the proposed solutions have been designed to run offline, with a complex installation in online devices [16,17].

Many broadcast protocols [18–21] have been proposed to alleviate the broadcast network saturation at a high packet delivery rate, reducing unnecessary rebroadcasts. In [18], an enhancement of multihop vehicular broadcast is proposed, but focusing on improvements in processing and packet retransmissions, where the nodes forward only as a function of the relative position of the sender. The work of [19] proposes a directional antennas and distance-based broadcast protocol, which is not efficient in real traffic situations. In [20], a vehicular broadcast protocol is presented, whose main novelty is the selection of the forwarder nodes using an opportunistic selection in order to reduce the network saturation. The authors of [21] solve some limitations of the existing broadcast protocols, but using information of the transmission range of the nodes, which depends on the status of the interferences and other factors.

However, this is still an open research field, with specific parameters from the transportation field that may influence the network behavior, like vehicle density, transmission range, length of target region and lifetime of a message [22]. Having detected a possible weakness in the algorithm of ETSI and that in the multiple papers [23–25] in which the ETSI algorithm is used, we opted to develop a new algorithm to solve the problems identified.

In this paper, we present some results of the iVANET project (Mesh Communications for Intelligent Vehicles and Infrastructures) that is under development at the Technical University of Madrid, where WSN technology is used as a basis to implement VANETs and test them in real traffic circulation situations in order to validate the feasibility of this technology to be used as a testbed for vehicular communication systems. In the first step of this project, the technology was tested in controlled and urban traffic situations [26]. In the second step, these systems were validated on public highways [15]. In the third step, the necessary standards and novel routing algorithms were applied to the developed vehicular communication systems to support GeoUnicast, testing them in real situations with real vehicles, the preliminary results being shown in [27]. Finally, in this paper, geo-broadcast algorithms to support vehicular communications are presented, including field tests and simulations to verify the right functionality of the developed systems, as well as the differences with the default algorithms defined in the standard.

2. Experimental Equipment

In order to test the geo-networking algorithms for vehicular communications, the IEEE 802.15.4 communication protocol and Maxfor Inc. (Bundang, Korea) MTM-CM3100 devices were used, based on the TelosB platform. These devices are wireless sensor network-supporting hardware, extended to support GPS positioning and geo-broadcast algorithms. This allows connectivity with any computer platform and acts as a gateway to access the vehicular network. This device uses TinyOS (TinyOS 2.0 Core Working Group: Stanford, CA, USA) as the operating system, a reduced version of Linux for embedded systems, like TelosB platforms. TinyOS is an open code operative system and allows personalizing its features and adapting applications in accordance with the necessities of the developers. In order to access the wireless network at 2.4 GHz, it uses the IEEE 802.15.4 standard at the physical and link levels, which guarantees the desired functionality of the VANET at transmission

rates of up to 250 Kbps and in a range of 100 m. Important features of IEEE 802.15.4 include real-time suitability by reservation of guaranteed time slots, collision avoidance through CSMA/CA and integrated support for secure communications.

The novel geo-broadcast protocol is a network protocol that resides jointly in the network and transport layer [6] and is executed in each gateway. This protocol is the manager to allow or not the dissemination of the broadcast packets that each node receives in accordance with the GPS position of the network nodes.

With this technology, it is possible to develop the necessary algorithms to optimize message routing in V2V communications. This algorithm can be used as a solution to be applied to other specifically designed vehicular communication technologies when available.

The MTM-CM3100 is configured to act as a broadcasting node by modifying the internal algorithms and the default protocol stack, connected to a computer, where the data received is saved. Each computer is also connected to a GPS receiver and to the CAN bus of each testbed vehicle in order to retrieve the vehicle data in real time.

Three communication nodes were used in the field experiments, which included one GPS receiver each and a CAN bus connection. It should be noted that it is possible to acquire speed using GPS receivers. However, we obtained speed from the vehicle CAN bus in order to be able to perform the experiments from standard equipment.

3. Geo-Broadcast Algorithm

The International Organization for Standardization's (ISO) standard Open System Interconnection reference model stack [28] has been modified in order to include the geo-networking functionality by adapting the transport and network levels, as shown in Figure 1.

Application Level

Geo-Networking + Basic
Transport

IEEE 802.2
Logical Link Control

IEEE 802.15.4
Medium Access Control

Physical Layers

Figure 1. Modification of the ISO model stack for geo-networking.

The core of the geo-networking is the geo-broadcast algorithm that runs at the network level of the ISO model stack, in charge of the logical addressing and the route determination. The main difficulty when defining a broadcast algorithm is the uncontrolled repetition of data packets that can generate loops, duplicate packets and reduce the efficiency of the network. The application in

the vehicular environment also enables the need to guarantee that the communications are transparent to the configuration of the network. That is, the data flow must be assured to every node in a determined geographical range regardless of the structure of the network, the routes followed by the packets and the positions of the vehicles.

The devices have a range of 100 m near emission. This range may be appropriate or not according to traffic conditions. For example, in urban traffic, it is a sufficient retransmission distance; however, at an average speed of 60 km/h (16 m/s), other vehicles from outside of the range of the 100 m of each vehicle must be able to receive safety data. The solution is to propagate the messages in a multihop behavior in order to cover the necessary safety range, and the geo-broadcast algorithm assures efficient data flow in this range.

Therefore, the philosophy of a generic geo-broadcast algorithm is that each node receives every message emitted in the reception range, and then, it decides whether to attend or ignore the message, in accordance with a series of parameters.

In the novel algorithm presented in this paper, we have designed a new method to avoid the saturation of the network by avoiding duplicated and uncontrolled message forwarding in accordance with the circulation speed of the receiving node (vehicle).

To avoid the saturation of the network, we have adopted two different and complementary strategies instead of limiting the messages to a certain number of hops, which is the common geographic and temporal strategy, as well as to avoid resending broadcast algorithms. Both strategies have been chosen to limit the forwarding of duplicated messages.

The first strategy is to avoid duplicate message forwarding. In each node, there is a register of forwarded messages with the identifier of the sender and the timestamp of the newer message of this sender, named the neighbor table (NT). When a new message is received, the identification of the origin node and timestamp of emission is analyzed. If this timestamp is earlier than the last record or the message is discontinued, that message is discarded.

The second strategy is to avoid uncontrolled packet forwarding by limiting the geo-broadcast retransmission area in accordance with a parameter. This retransmission area, or area of influence, is defined in accordance with the circulation speed of each node in order to adapt the network features according to the traffic flow and to improve the efficiency of the data distribution. Then, a variable named "range of rebroadcasting" of a receiving node *i* (RRB*i*) was defined (Equation (1)) in order to determine the forwarding area of each node, where S represents the speed of the vehicle, RT, the reaction time of the driver (1 s) and SD represents the distance required to stop the vehicle in an emergency stop.

$$RRB_i = S \cdot RT + SD \tag{1}$$

The selection of this metric assures that, in the event of an emergency message that requires an emergency stop, the vehicle has enough space to brake safely and avoid an accident.

If a sender node is at a lower distance from node i than RRB_i, its data packets will be forwarded by the receiving node i. In any other case, node i discards the incoming message from this node.

With this strategy, we extend the range of message dissemination in an area that depends on the speed of the nodes. The maximum range for each node i is defined by Equation (2).

$$R_{i} = d_{i,n} + R_{n} \mid d_{i,n} < RRB_{n} \tag{2}$$

where $d_{i,n}$ represents the distance between the origin node (i) and the farthest node (n) and Rn represents the range of transmission of this farthest node.

Therefore, the RRB limits the uncontrolled dissemination of messages across the network. This concept is shown in Figure 2, where the top schema represents the transmission range of each node and the second schema represents the RRB range of the first two receiving nodes from the origin. Figure 2 at the top shows that the distance between Vehicle 0 (V0) and Vehicle 3 (V3) is greater than the range of Vehicle 0 (R0), and the speed of Vehicle 1 (V1) is higher than the speed of Vehicle 2 (V2). However, there is a need for data transmission between both extreme vehicles. The proposed algorithm manages the network data multihop broadcast in accordance with the RRB parameter. Then, the message from Vehicle 0 to Vehicle 3 is directly received by Vehicles 1 and 2, but only Vehicle 1 resends the message to 3, because the distance between Vehicles 0 and 1 is less than the RRB₁. On the contrary, Vehicle 2 does not resend the message due to its RRB₂ being less than the distance between Vehicles 0 and 2.

Figure 2. Graphical description of the "range of rebroadcasting" (RRB) concept.

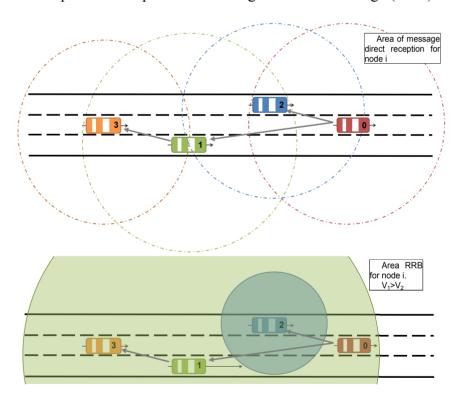
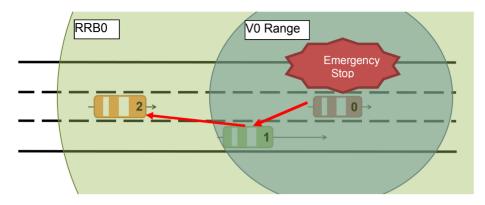


Figure 3 illustrates an example of the range of rebroadcasting concept, in an emergency stop situation. In this case, Vehicle 0 (V0) circulates at a speed of 100 km/h (3.6 m/s) and starts an emergency stop, with a deceleration of -10 m/s^2 . In this case, the value of RRB₀ is 63.45 m. On the contrary, the range of this vehicle is 30 m. This means that some hops are necessary in order for the emergency messages emitted from V0 to reach any vehicle that circulates in the RRB₀ area. In the case of Figure 3, one vehicle (V1) circulates in the V0 range and re-broadcasts the emergency stop message to the other vehicles in its range (V2).

Figure 3. Example of geo-broadcast message flow in an emergency situation. V0, Vehicle 0.



3.1. Data Packets

In order to support the data transmission and the geo-broadcast network organization, only one data packet has been defined, shown in Table 1. This data packet is based on the IEEE 802.15.4 standard WSN frame, where the PHY (physical layer of the OSI model) header includes a preamble, a start packet delimiter and the packet payload length. On the other hand, the MAC (Medium Access Control Layer of the OSI model) header includes a frame control, the frame sequence number, the destination WSN ID, the destination MAC address, the source WSN ID and the source MAC address.

Table 1. Geo-broadcast data packet.

Field	PHY Header	MAC Header	Geo-broadcast Header	Payload	MAC Footer
Length	6 bytes	10 bytes	27 bytes	0-205 bytes	2 bytes

The PHY and MAC headers and the MAC footer are standard, and the payload is defined in accordance with the destination application. The geo-broadcast header is defined in Table 2 and contains the necessary information to support the algorithm performance.

Table 2. Geo-broadcast header. Fr, header start and finish byte; PT, protocol type; Sq, sequence number; Dst, destination node; Org, origin node; Ln, length;

Byte	0	1	2	3	4	5	6	7	8	9	10	11–25	26	27	28
Name	Fr	PT	Sq	PF	D	st		rg	Ln	Gr	AM	GeoData	CI		Fr

Fr is the header start and finish byte (7Eh). PT is the protocol type and indicates if the message source is the local node or has been received from the radio. Sq is the sequence number of the packet. Dst and Org represent the identification of the destination node and the origin node, respectively, at the application level, defined by the user of the front end applications. They represent the identifiers of the network nodes at the application level. The standard size for the IEEE 802.15.4 WSN IDs is two bytes in order to reduce the size of the data packets. Ln indicates the length of the data field, Gr is the Group of the origin, and AM (Active Message) represents the type of message (geo-broadcast). Finally, the GeoData field (Table 3) contains the necessary geospatial and navigation information to support the algorithm.

Table 3. GeoData field. UTM_N, north UTM coordinates; UTM_E, east UTM coordinates; BZ, UTM band and zone; Ang, azimuth angle; TST, timestamp.

Byte	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
Name	Type	II)	J	JTM_1	1	U	TM_I	Ξ	ΒZ	Speed	Ang	TS	ST	Sum

Type indicates if the emitter node is the origin of the packet or only the forwarder. ID source is an identification of the node source of the packet. UTM_N and UTM_E represent the north and east UTM coordinates of the origin of the packet. BZ is the UTM band and zone. Speed represents the circulation speed of the node. Ang is the azimuth angle of the trajectory of the node. TST is a timestamp (mm:ss,d), and Sum is a data control field.

3.2. Algorithm Description

Figure 4 contains the flowchart of the geo-broadcast algorithm, described in detail next. When a new data packet is generated by a node *i*, it is analyzed; the geoinformation is extracted and saved. With this data, the RBBi is calculated. Once the calculation has been made, the geo-routing headers are added to the data packet, which is then sent to the vehicular network through the radio.

When data is received from the vehicular network, the geodata headers are extracted. With this information, the node checks whether the package has already been forwarded or is discontinued. If the data packet fulfills the requirements (new data packet), the payload is sent to the computer and the neighbor table (NT) is updated. To update the NT, the node checks if the node exists in the NT. In the affirmative case, the data is updated, and in the negative case, the node is included in the NT. If the NT is full, the oldest neighbor is replaced by the new node. Finally, the distance between the nodes is calculated, and if the distance is less than RRBi, the data packet is forwarded through the vehicular network.

Ш ОК NO Packet GeoData Duplicated Radio IN Serial OUT Validation Extraction or old Neighbor Table update and distance calculus YES YES NO ERROR END of Distance Radio OUT < RRB Algorithm VII RBB calculus Save Packet OK GeoData Serial IN Position Neighbor Validation Extraction Table update **ERROR** END of Radio OUT Algorithm

Figure 4. Flowchart of the broadcast algorithm.

3.3. Step by Step

Described below is the performance of the algorithm in pseudo code, following the schema of Figure 1.

- i. Reception of a packet via radio, checking its format.
 - If true: the data is passed to II.
 - If not: the packet is discarded and data packet analysis is ended.
- ii. The geoinformation is extracted: position and timestamp.
- iii. Check timestamp to verify that the packet has not already been forwarded or it is updated.
 - If true, two actions are executed: the packet is sent through the serial port to notify the PC and the packet is transmitted to step IV.
 - If not: the packet is discarded and data packet analysis is ended.
- iv. The distance is calculated between transmitter and receiver node *i*. Update the neighbor table.
- v. Check if the distance between nodes is less than RRB_i .
 - If true: the packet is forwarded by radio.
 - If not: the packet is discarded and data packet analysis is ended.
- vi. Reception of a packet via serial port, checking its format.
 - If true: the data is passed to vii and is sent via radio.
 - If not: the packet is discarded and data packet analysis is ended.
- vii. The geoinformation is extracted: position and timestamp.
- viii. The RRB_i is calculated with the speed of the vehicle, i. The position of the node is read to allow the calculation of the distances, and the NT is checked to eliminate discontinued nodes.

4. Comparison with Other Algorithms

IEEE 802.11p defines the physical and link layer protocols for the communication devices of vehicular communications, the basis for the dedicated short range communications (DSRC). The ETSI TS 102 626-4-1 protocol family defines the geo-networking algorithms to be applied in vehicular communications DSRC, operating in the network and transport layer. These algorithms include GeoUnicast and geo-broadcast.

The main difference between the IEEE 802.11p standard with ETSI TS 102 636-4-1 and the presented algorithm is the philosophy of the entity in charge of deciding if a message is forwarded.

In ETSI TS 102 636-4-1, each node forwards the received packets with the only condition that the node is in the broadcast area. This reception area is defined in accordance with a central point and a radius, but without considering the circulation speed or the traffic status. The validity of the message is calculated according to the number of hops and the value of the "lifetime" timestamp parameter. The calculus of the optimum number of hops is difficult with the information managed with the algorithm, since no traffic status considerations are taken into account, raising the possibility of network saturation.

In the proposed algorithm, the reception area is defined in accordance with the speed and some safety parameters. This means that the broadcast coverage can be adapted depending on the traffic condition and the number of vehicles in the surroundings of each node. Due the selection of this

parameter, the number of hops transparent to the algorithm is not necessary in order to calculate the validity of the message, and only the timestamp is used for this task, thereby avoiding the possibility of network saturation. This philosophy also improves some of what the broadcast algorithms lack [18–21], described previously.

Take an example for better understanding of the advantages of the proposed geo-broadcast, using the same technology of communication for both algorithms, a 128-kbs bandwidth and a range of 600 m. In the algorithm ETSI TS 102 636-4-1, a fixed-broadcast radius of 40 m is used.

Figures 5a,b and 6 show the results of a simulation that compares the performance of these algorithms. This simulation compares their behavior as a function of the circulation speed of the vehicles and the number of vehicles the maximum communication range in a V2V highway situation (six lanes). With these data, it is possible to estimate the maximum available size of the packet (MP) to avoid network saturation (Equation (3)) and the number of packets received (PR) by each vehicle, considering an emission frequency of 1 Hz, using Equation (4).

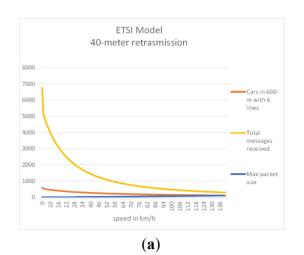
$$MP = \frac{Bandwidth}{PR} \tag{3}$$

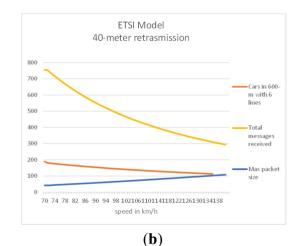
$$PR = (NV \cdot 600 \,\mathrm{m}) \cdot (NV \cdot BCR) + (NV \cdot BCR) \cdot (NR \cdot BCR)$$
(4)

where *Bandwidth* indicates the bandwidth of the network (128 Kbps), *NV* represents the number of vehicles in a range of 600 m, *BCR* represents the preset range of rebroadcast and NR represents the number of rebroadcasted messages per vehicle. The first field of Equation (4) indicates the number of the packets generated by the vehicles in the range, including the rebroadcasts; the second field indicates the packets of the vehicles from outside the range, but rebroadcast by other vehicles inside their range.

As a result of this simulation, the stability and better performance of the proposed algorithm are clear, since its packet size remains constant around 51 bytes, while the proposed ETSI TS 102 636-4-1 algorithm varies from five bytes to 109 bytes (Figure 5a), which allows one to increase the number of received messages, even in situations of high road occupancy (Figure 5b). In Figure 6 is also shown the RRB of the proposed algorithm.

Figure 5. (a) European Telecommunications Standards Institute (ETSI) model; (b) details of the ETSI model.





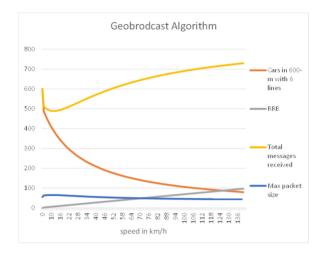


Figure 6. Proposed algorithm.

As result of these simulations, it is clear that with the proposed algorithm, the maximum size of the data packet remains constant, as well as the total number of messages received, improving the results of the ETSI algorithm.

5. Data Analysis

To verify the algorithm, we proceeded to test the system in simulation and field trials. The simulation was used to check the right performance of the RRB in accordance with its parameters. Next, a set of field tests was carried out in order to verify the right performance of the algorithms with the selected parameters.

Two kinds of field tests were developed: static and dynamic. The selected circulation speed for every test was 60 km/h, which generates an RBBi of 40 m, in order to emulate the situation of vehicle circulation in a dense traffic situation.

5.1. Static

The objective of the static test is to check that the geo-broadcast algorithm runs in the situation where the network is established and no nodes are merging or leaving the network. Therefore, in this test, the system ran for 15 min (913 s) with three nodes broadcasting in static positions, where the distance between them was always less than 40 m (Figure 7).

D < 40 m
D < 40 m

Figure 7. Static test.

In this experiment, the stability of the algorithm was proven. During the experiment, 5,431 messages were sent, and an average of 91.4% was received, as shown in Tables 4 and 5. The cause of this message loss, about 9%, is the high level of interference in the 2.4 GHz band that causes this distance to be at the limit of the coverage, and some packets are not received in the receiving nodes.

Table 4. Messages sent and received in the static test.

Node	Total received	Total Sent
1	3376	1825
2	3185	1773
3	3368	1883
System	9929	5431

Table 5. Node percentage in the static test.

	M.J.		Received	
	Node	1	2	3
	1	X	87	92
Sent	2	94	X	95
	3	93	87	X

In order to compare the results of this study with those of other related proposals, we have replicated the field experiments with a different platform. Concretely, the static tests have been replicated using an IEEE 802.11p DSRC platform, running with the geo-broadcast algorithm defined in ETSI TS 102 636-4-1, the one defined in the DSRC standard, with a duration of 206 s.

As shown in Table 6, the results for these experiments are similar, obtaining a 100% emission/reception rate. The cause of this higher rate is the 5.9-GHz emission band of the 5.9 GHz of the DSRC platform, with a low level of interferences, which assures that the coverage is right to receive every data packet.

Table 6. Messages sent and received in the static test.

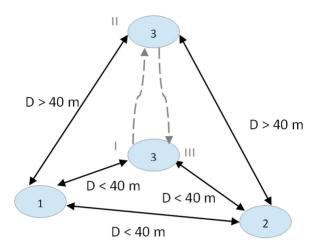
Node	Total received	Total sent
1	412	206
2	412	206
3	412	206
System	1236	618

5.2. Dynamic

In the dynamic test, the system ran for more than 6 min (370 s) with two nodes initially emitting with a distance between them of always less than 40 m (Figure 8I) and a third node at a farther RRB distance from the other two.

Therefore, the third node reduces its distance with the others, successfully merging the network (Figure 8II) and initiating the geo-broadcast message. Finally, this last node separates from the other two at a farther distance than the maximum RRB range, leaving the network (Figure 8III) and stopping the data broadcasting. This test tries to emulate the situation of a vehicle that overtakes another.

Figure 8. Dynamic test.



This test simulates the phases of one vehicle overtaking two vehicles on a road. During the experiment, 2198 messages were sent among the nodes, of which an average of 91.1% were received, as shown in Tables 7 and 8.

Table 7. Messages sent and received in dynamic test.

Node	Total received	Total sent
1	1298	720
2	1398	739
3	1337	739
System	4003	2198

Table 8. Node percentage in dynamic test.

	Mada		Received	
	Node -	1	2	3
	1	X	95	94
Sent	2	87	X	90
	3	88	92	X

In this case, a comparison was performed, with the same use case, with the IEEE 802.11p-based platform, running with the geo-broadcast algorithm defined in ETSI TS 102 636-4-1, with a duration of 188 s.

In this second situation, the results with this alternative platform are similar to the previous ones and shown in Table 9, with an emission/reception rate of 99%.

Table 9. Messages sent and received in dynamic test.

Node	Total Received	Total Sent
1	374	188
2	376	188
3	376	188
System	1126	564

The two tests demonstrate that the algorithm is stable under static and dynamic conditions, with only 9% of packets being lost and a measurement range of 87%–95%. This is because the tests were conducted in an environment of high radio interference and because of the nature of the prototype devices used, where their electronic characteristics may vary from one to another.

If we compare the results with the experiments performed with the alternative platform, both algorithms comply with their function of geo-broadcast perfectly, but the ETSI TS 102 636-4-1 has the limitation of saturation described in Section 4, a limitation that is not present in the novel geo-broadcast algorithm described in this paper. In other words, the network and transport layers of both platforms run similarly; they even have different geo-broadcast algorithms, and only some minor differences are caused by the different protocols at the physical level (WSN protocol IEEE 802.15.4 and WiFi protocol IEEE 802.11p).

6. Conclusions

A novel geo-broadcast algorithm that supports V2V mesh communications has been presented. This algorithm has been defined and implemented using IEEE 802.15.4 communication modules and tested in real situations, using the range of rebroadcasting (RRB) as a key parameter to define the range of each message. This parameter allows adapting the range of each message to the traffic circumstances, taking into account the circulation speed and the stopping time of the vehicles. The use of RRB reduces the saturation of vehicular communications, extending the data flow when vehicles circulate at high speeds and reducing it in low speed conditions. The level of robustness and efficiency of the algorithm is able to successfully manage vehicular communications to serve ADAS applications. The validation of the algorithm has been developed in controlled tests, with the expected results being obtained in performance, latency and the number of successful packets received, and is therefore appropriate to support VANET communications in real V2V. Geo-broadcast has also been compared with the former algorithms that are defined in IEEE 803.11p, showing some benefits from the point of view of management in situations with a high traffic flow. Having demonstrated its performance, it is expected that the proposed algorithm can be tested on an 802.11p communication platform, following the actual standards defined by IEEE and ETSI.

An interesting contribution of this work could be the lessons learned during the design, implementation and configuration of all of the necessary elements for the field test. In this case, some limitations have been found when we try to translate the wireless sensor network experience to a vehicular environment, recently known as the Internet of vehicles. The first limitations are that the algorithms and internal tools of WSN have been designed to run in static environments, like buildings, fences or similar environments, with a low rate of mobility. This means that the internal broadcast algorithms have had to be redesigned in order to run in a vehicular environment. Additionally, geographical capabilities have also been added to these algorithms, in order to comply with experimental requirements. Two important limitations have been found in the WSN technologies: on the one hand, the limited CPU processing capacity of the usual nodes, which makes them saturate in high CPU demand situations, like the case of mobility; the second is the use of the 2.4-GHz frequency band, shared with WiFi, and with a high level of interference, which causes a reduction in the range that is, in some cases, dramatic.

Acknowledgments

This work has received the support of the Spanish Ministry of Economy and Competitivity (TRA2010-0225-C03-03, TRA2010-15645, TRA2009-07505 and TRA2013-48314-C3-2-R) and the Region of Madrid Excellence Network SEGVAUTO-II (Safety for Automobile Vehicles).

Author Contributions

José E. Naranjo and Felipe Jiménez were in charge of the direction of the research activities. Nuria Gómez was in charge of the coordination of the technical activities. José J. Anaya and Edgar Talavera were in charge of the technical set up and the development of the research activities.

Conflicts of Interest

The authors declare no conflict of interest.

References

- 1. US DOT. *ITS Strategic Research Plan*, 2010–2014; US Department of Transportation, Research and Innovative Technology Administration: Washington, DC, USA, 2010.
- 2. Wolterink, W.K.; Heijenk, G.; Karagiannis, G. Dissemination protocols to support cooperative adaptive cruise control (CACC) merging. In Proceedings of the 2011 11th International Conference on ITS Telecommunications (ITST), St. Petersburg, Russia, 23–25 August 2011; pp. 15–20.
- 3. Demmel, S.; Gruyer, D.; Rakotonirainy, A. Comparing cooperative and non-cooperative crash risk-assessment. In Proceedings of the 2013 IEEE Intelligent Vehicles Symposium (IV), Gold Coast City, Australia, 23–26 June 2013; pp. 1007–1013.
- 4. Lee, J.; Park, B. Development and Evaluation of a Cooperative Vehicle Intersection Control Algorithm Under the Connected Vehicles Environment. In Proceedings of the IEEE Transactions on Intelligent Transportation Systems, Chemnitz, Germany, 20–23 March 2012; Volume 13, No. 1, pp. 81–90.
- 5. Bohm, M.; Frotscher, A. Bidirectional I2V communication applications for advanced invehicle ISA systems. In Proceedings of the 2009 IEEE Vehicular Networking Conference (VNC), 28–30 October 2009, Tokyo, Japan; Volume 1, No. 6, pp. 28–30.
- 6. ETSI TS 102 636-6-1 V1.1.1. Intelligent Transport Systems (ITS), Vehicular Communications, Geo-networking; Part 6: Internet Integration, Sub-part 1: Transmission of IPv6 Packets over Geo-networking Protocols; European Telecommunications Standards Institute: Sophia Antipolis, France, 2011.
- 7. Jiménez, F.; Naranjo, J.E.; García, F.; Armingol, J.M. Can low-cost road vehicles positioning systems fulfil accuracy specifications of new ADAS applications? *J. Navig.* **2011**, *64*, 251–264.
- 8. Matolak, D.W.; Xiong, I.; Sen, W. Channel Modeling for V2V Communications. In *Mobile and Ubiquitous Systems Workshop 2006*; Proceedings of the 3rd Annual International Conference, San Jose, CA, USA, 17–21 July 2006.

9. Choi, J.H.; Khaled, Y.; Tsukada, M.; Ernst, T. IPv6 support for VANET with geographical routing. In Proceedings of the 8th International Conference on Intelligent Transport System Telecommunications (ITST'08), Phuket, Thailand, 22–24 October 2008.

- 10. Gordillo, A.; Calderón, M.; Bernardos, C.J. Enabling IP geomulticast services for vehicular networks. In Proceedings of the ITST'09, Lille, France, 20–22 October 2009.
- 11. Van Hoesel, L.F.W.; Tüysüz Erman, A.; Dilo, A.; Havinga, P.J.M. Geo-casting of queries combined with coverage area reporting for wireless sensor networks. *Ad. Hoc. Net.* **2013**, *11*, 104–123.
- 12. Yu, Y.; Govindan, R.; Estrin, D. *Geographical and Energy Aware Routing: A Recursive Data Dissemination Protocol for Wireless Sensor Networks, Technical Report UCLA/CSD-TR-01-0023*; University of Southern California: Los Angeles, CA, USA, 2001.
- 13. Ko, Y.; Vaidya, N. Anycasting-based protocol for geocast service in mobile ad hoc networks. *Comput. Networks* **2003**, *41*, 743–760.
- 14. Coman, A.; Nascimento, M.A.; Sander, J. A framework for spatiotemporal query processing over wireless sensor networks. In Proceedings of the 1st International Workshop on Data Management for Sensor Networks: In Conjunction with VLDB, 2004, Toronto, ON, Canada, 30 August 2004.
- 15. Baur, M.; Fullerton, M.; Busch, F. Realizing an effective and flexible ITS evaluation strategy through modular and multi-scaled traffic simulation. *IEEE Intell. Trans. Syst. Mag.* **2010**, *2*, 34–42.
- 16. Karp, B.; Kung, H.T. GPSR: Greedy perimeter stateless routing for wireless networks. In Proceedings of the ACM MOBICOM, Boston, MA, USA, 6–11 August 2000; pp. 243–254.
- 17. Lochert, C.; Mauve, M.; Füßler, H.; Hartenstein, H. *Geographic Routing in City Scenarios*; ACM SIGMOBILE Mobile Computing and Communications Review: New York, NY, USA, 2005; Volume 9, pp. 69–72.
- 18. Mariyasagayam, M.N.; Osafune, T.; Lenardi, M. Enhanced Multi-Hop Vehicular Broadcast (MHVB) for Active Safety Applications. In Proceedings of the 7th IEEE International Conference on ITS Telecommunications (ITST), Sophia Antipolis, France, 6–8 June 2007.
- 19. Li, D.; Huang, H.; Li, X.; Li, M.; Tang, F. A Distance-based directional broadcast protocol for urban vehicular ad hoc network. In Proceeding of the WiCom 2007. International Conference on Wireless Communications, Networking and Mobile Computing, Shanghai, China, 21–25 September 2007.
- 20. Vegni, A.M.; Stramacci, A.; Natalizio, E. SRB: A selective reliable broadcast protocol for safety applications in VANETs. In Proceedings of the IEEE International Conference on Selected Topics in Mobile and Wireless Networking, Avignon, France, 2–4 July 2012.
- 21. Wu, X.; Song, S.; Wang, H. A novel position-based multi-hop broadcast protocol for vehicular ad hoc networks. *J. Netw.* **2012**, *6*, 112–120.
- 22. Rajendran, R. The Evaluation of Geo-networking Forwarding in Vehicular Ad-Hoc Networks. Master Thesis, Hamlstad University, Halmstad, Sweden, November 2013.
- 23. Ben Jemaa, I.; Shagdar, O.; Ernst, T. A framework for IP and non-IP multicast services for vehicular networks. In Proceeding of the 2012 Third International Conference on the Network of the Future (NOF), Tunis, Tunisia, 21–23 November 2012.

24. Taysi, Z.C.; Yavuz, A.G. ETSI compliant Geo-networking protocol layer implementation for IVC simulations. In Proceeding of the 2012 International Conference on Computer, Information and Telecommunication Systems (CITS), Amman, Jordan, 14–16 May 2012.

- 25. Toukabri, T.; Tsukada, M.; Ernst, T.; Bettaieb, L. Experimental evaluation of an open source implementation of IPv6 Geo-networking in VANETs. In Proceeding of the 2011 11th International Conference on ITS Telecommunications (ITST), St. Petersburg, Russia, 23–25 August 2011, pp. 237–245.
- 26. Naranjo, J.E.; Anaya, J.J.; Talavera, E.; Zato, J.G.; Jiménez, F.; Gómez, N. Evaluation of V2V and V2I Mesh Prototypes based on a Wireless Sensor Network. In Proceeding of the 14th International IEEE Conference on Intelligent Transportation Systems, Washington, DC, USA, 5–7 October 2011.
- 27. Anaya, J.J.; Talavera, E.; Zato, J.G.; Jiménez, F.; Gómez, N.; Naranjo, J.E. Geo-networking based V2V Mesh Communications over WSN. In Proceeding of the IEEE Intelligent Transportation Systems Conference 2013, The Hague, The Netherlands, 6–9 October 2013.
- 28. ISO/IEC 7498-1, Information technology—Open Systems Interconnection—Basic Reference Model: The Basic Model, International Standard; International Organization for Standardization (ISO): Geneva, Switzerland, 1994.
- © 2014 by the authors; licensee MDPI, Basel, Switzerland. This article is an open access article distributed under the terms and conditions of the Creative Commons Attribution license (http://creativecommons.org/licenses/by/3.0/).