

Article

The Complexity of Some Classes of Pyramid Graphs Created from a Gear Graph

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Abstract: The methods of measuring the complexity (spanning trees) in a finite graph, a problem related to various areas of mathematics and physics, have been inspected by many mathematicians and physicists. In this work, we defined some classes of pyramid graphs created by a gear graph then we developed the Kirchhoff's matrix tree theorem method to produce explicit formulas for the complexity of these graphs, using linear algebra, matrix analysis techniques, and employing knowledge of Chebyshev polynomials. Finally, we gave some numerical results for the number of spanning trees of the studied graphs.

Keywords: complexity; Chebyshev polynomials; gear graph; pyramid graphs

MSC: 05C05, 05C50

1. Introduction

The graph theory is a theory that combines computer science and mathematics, which can solve considerable problems in several fields (telecom, social network, molecules, computer network, genetics, etc.) by designing graphs and facilitating them through idealistic cases such as the spanning trees, see [1–10].

A spanning tree of a finite connected graph G is a maximal subset of the edges that contains no cycle, or equivalently a minimal subset of the edges that connects all the vertices. The history of enumerating the number of spanning trees $\tau(G)$ of a graph G dates back to 1842 when the physicist Kirchhoff [11] offered the matrix tree theorem established on the determinants of a certain matrix gained from the Laplacian matrix L defined by the difference between the degree matrix D and adjacency matrix A , where D is a diagonal matrix, $D = \text{diag}(d_1, d_2, \dots, d_n)$ corresponding to a graph G with n vertices that has the vertex degree of d_i in the i th position of a graph G and A is a matrix with rows and columns labeled by graph vertices, with a 1 or 0 in position (u_i, u_j) according to whether u_i and u_j are adjacent or not. That is

$$L_{ij} = \begin{cases} a_i & \text{if } i = j \\ -1 & \text{if } i \neq j \text{ and } i \text{ is adjacent to } j \\ 0 & \text{otherwise} \end{cases},$$

where a_i denotes the degree of the vertex i .

This method allows beneficial results for a graph comprising a small number of vertices, but is not feasible for large graphs. There is one more method for calculating $\tau(G)$. Let $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_k = 0$

denote the eigenvalues of the matrix L of a graph G with n vertices. “Kelmans” and “Chelnokov” [12] have derived that

$$\tau(G) = \frac{1}{k} \prod_{i=1}^{k-1} \lambda_k.$$

One of the favorite methods of calculating the complexity is the contraction–deletion theorem. For any graph G , the complexity $\tau(G)$ of G is equal to $\tau(G) = \tau(G - e) + \tau(G/e)$, where e is any edge of G , and where $G - e$ is the deletion of e from G , and G/e is the contraction of e in G . This gives a recursive method to calculate the complexity of a graph [13,14].

Another important method is using electrically equivalent transformations of networks. Yilun Shang [15] derived a closed-form formula for the enumeration of spanning trees the subdivided-line graph of a simple connected graph using the theory of electrical networks.

Many works have conceived techniques to derive the number of spanning trees of a graph, some of which can be found at [16–18].

Now, we give the following Lemma:

Lemma 1 [19]. $\tau(G) = \frac{1}{k^2} \det(kI - D^c + A^c)$ where A^c and D^c are the adjacency and degree matrices of G^c , the complement of G , respectively, and I is the $k \times k$ identity matrix.

The characteristic of this formula is to express $\tau(G)$ straightway as a determinant rather than in terms of cofactors as in Kirchhoff theorem or eigenvalues as in Kelmans and Chelnokov formula.

2. Chebyshev Polynomial

In this part we insert some relations regarding Chebyshev polynomials of the first and second types which we use in our calculations.

We start from their definitions, see Yuanping, et al. [20].

Let $A_n(x)$ be $n \times n$ matrix such that

$$A_n(x) = \begin{pmatrix} 2x & -1 & 0 & \cdots & 0 \\ -1 & 2x & -1 & \ddots & \vdots \\ 0 & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & -1 \\ 0 & \cdots & 0 & -1 & 2x \end{pmatrix}.$$

Furthermore, we render that the Chebyshev polynomials of the first type are defined by

$$T_n(x) = \cos(n \cos^{-1} x) \quad (1)$$

The Chebyshev polynomials of the second type are defined by

$$U_{n-1}(x) = \frac{1}{n} \frac{d}{dx} T_n(x) = \frac{\sin(n \cos^{-1} x)}{\sin(\cos^{-1} x)} \quad (2)$$

It is easily confirmed that

$$U_n(x) - 2xU_{n-1}(x) + U_{n-2}(x) = 0 \quad (3)$$

It can then be shown from this recursion that by expanding $\det A_n(x)$ one obtains

$$U_n(x) = \det(A_n(x)), n \geq 1 \quad (4)$$

Moreover, by solving the recursion (3), one gets the straightforward formula

$$U_n(x) = \frac{(x + \sqrt{x^2 - 1})^{n+1} - (x - \sqrt{x^2 - 1})^{n+1}}{2\sqrt{x^2 - 1}}, \quad n \geq 1, \quad (5)$$

where the conformity is valid for all complex x (except at $x = \pm 1$, where the function can be taken as the limit).

The definition of $U_n(x)$ easily yields its zeros and it can therefore be confirmed that

$$U_{n-1}(x) = 2^{n-1} \prod_{j=1}^{n-1} (x - \cos \frac{j\pi}{n}) \quad (6)$$

One further notes that

$$U_{n-1}(-x) = (-1)^{n-1} U_{n-1}(x) \quad (7)$$

From Equations (6) and (7), we have:

$$U_{n-1}^2(x) = 4^{n-1} \prod_{j=1}^{n-1} (x^2 - \cos^2 \frac{j\pi}{n}) \quad (8)$$

Finally, straightforward manipulation of the above formula produces the following formula (9), which is highly beneficial to us later:

$$U_{n-1}^2(\sqrt{\frac{x+2}{4}}) = \prod_{j=1}^{n-1} (x - 2 \cos \frac{2j\pi}{n}) \quad (9)$$

Moreover, one can see that

$$U_{n-1}^2(x) = \frac{1 - T_{2n}(x)}{2(1 - x^2)} = \frac{1 - T_n(2x^2 - 1)}{2(1 - x^2)} \quad (10)$$

$$T_n(x) = \frac{1}{2} [(x + \sqrt{x^2 - 1})^n + ((x - \sqrt{x^2 - 1})^n)] \quad (11)$$

Now we introduce the following important two Lemmas.

Lemma 2 [21]. Let $B_n(x)$ be $n \times n$ Circulant matrix such that

$$B_n(x) = \begin{pmatrix} x & 0 & 1 & \cdots & 1 & 0 \\ 0 & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & x \end{pmatrix}.$$

Then for $n \geq 3$, $x \geq 4$, one has

$$\det(B_n(x)) = \frac{2(x+n-3)}{x-3} [T_n(\frac{x-1}{2}) - 1].$$

Lemma 3 [22]. If $A \in F^{n \times n}$, $B \in F^{n \times m}$, $C \in F^{m \times n}$ and $D \in F^{m \times m}$. Suppose that A and D are nonsingular matrices, then:

$$\det \begin{pmatrix} A & B \\ C & D \end{pmatrix} = \det(A - BD^{-1}C)\det D = \det A \det(D - CA^{-1}B).$$

This Lemma gives a type of symmetry for some matrices which simplify our calculations of the complexity of graphs studied in this paper.

3. Main Results

Definition 1. The pyramid graph $A_n^{(m)}$ is the graph created from the gear graph G_{m+1} with vertices $\{u_0; u_1, u_2, \dots, u_m; w_1, w_2, \dots, w_m\}$ and m sets of vertices, say, $\{v_1^1, v_2^1, \dots, v_n^1\}, \{v_1^2, v_2^2, \dots, v_n^2\}, \dots, \{v_1^m, v_2^m, \dots, v_n^m\}$, such that for all $i = 1, 2, \dots, n$ the vertex v_i^j is adjacent to u_j and u_{j+1} , where $j = 1, 2, \dots, m-1$, and v_i^m is adjacent to u_1 and u_m . See Figure 1.

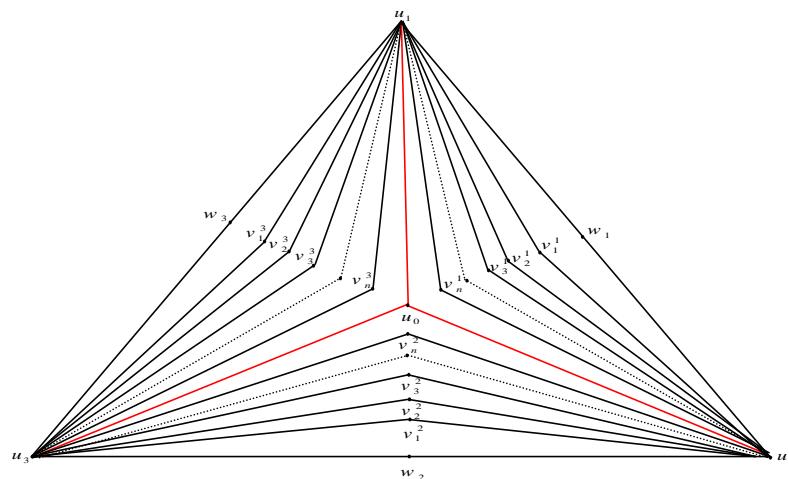


Figure 1. The pyramid graph $A_n^{(3)}$.

Theorem 1. For $n \geq 0$, $m \geq 3$, $\tau(A_n^{(m)}) = 2^{mn}[(n+2+\sqrt{2n+3})^m + (n+2-\sqrt{2n+3})^m - 2(n+1)^m]$.

Proof. Using Lemma 1, we have

$$\tau(A_n^{(m)}) = \frac{1}{(mn+2m+1)^2} \times \det((mn+2m+1)I - D^c + A^c) = \frac{1}{(mn+2m+1)^2} \times$$

Let $j = (1 \cdots 1)$ be the $1 \times n$ matrix with all one, and J_n be the $n \times n$ matrix with all one. Set $a = 2n + 4$ and $b = mn + 2m + 1$. Then we obtain:

Using Lemma 3, yields

$$\tau(A_n^{(m)}) = \frac{1}{b} \times \det \begin{pmatrix} A & B \\ C & 2I_{mn} \end{pmatrix} = \frac{1}{b} \times \det(A - B \frac{1}{2I_{mn}} C) \times 2^{mn}$$

$$= \frac{1}{b} 2^{mn} \times 2^{-2m} \times \det \begin{pmatrix} 2a & n+2 & 2(n+1) & \cdots & 2(n+1) & n+2 & -2 & 0 & \cdots & \cdots & 0 & -2 \\ n+2 & 2a & n+2 & 2(n+1) & \cdots & 2(n+1) & -2 & \ddots & \ddots & \ddots & \ddots & 0 \\ 2(n+1) & n+2 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2(n+1) & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2(n+1) & \ddots & \ddots & \ddots & \ddots & n+2 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ n+2 & 2(n+1) & \cdots & 2(n+1) & n+2 & 2a & 0 & \cdots & \cdots & 0 & -2 & -2 \\ 0 & 0 & 2 & \cdots & \cdots & 2 & 4 & 0 & \cdots & \cdots & \cdots & 0 \\ 2 & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & 0 \\ 0 & 2 & \cdots & \cdots & 2 & 0 & 0 & \cdots & \cdots & \cdots & 0 & 4 \end{pmatrix}$$

Using Lemma 3 again, yields

$$\tau(A_n^{(m)}) = \frac{2^{mn-2m}}{b} \times \det \begin{pmatrix} D & E \\ F & 4I_m \end{pmatrix} = \frac{2^{mn}}{b} \times \det(D - E \frac{1}{4I_m} F)$$

$$\tau(A_n^{(m)}) = \frac{2^{mn}}{b} \times \det \begin{pmatrix} 2a & (n+3) & 2(n+2) & \cdots & 2(n+2) & (n+3) \\ (n+3) & 2a & (n+3) & \ddots & \cdots & 2(n+2) \\ 2(n+2) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2(n+2) \\ 2(n+2) & \ddots & \ddots & \ddots & \ddots & (n+3) \\ (n+3) & 2(n+2) & \cdots & 2(n+2) & (n+3) & 2a \end{pmatrix}$$

Straightforward inducement using the properties of determinants, one can obtain

$$\begin{aligned} \tau(A_n^{(m)}) &= \frac{2^{mn}}{b} \times \frac{2b}{mn+m+2} \times \det \begin{pmatrix} (2a-n-3) & 0 & (n+1) & \cdots & (n+1) & 0 \\ 0 & (2a-n-3) & 0 & \ddots & \cdots & (n+1) \\ (n+1) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & (n+1) \\ (n+1) & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & (n+1) & \cdots & (n+1) & 0 & (2a-n-3) \end{pmatrix} \\ &= \frac{2^{mn+1}(n+1)^m}{mn+m+2} \times \det \begin{pmatrix} \frac{(2a-n-3)}{(n+1)} & 0 & 1 & \cdots & 1 & 0 \\ 0 & \frac{(2a-n-3)}{(n+1)} & 0 & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & \frac{(2a-n-3)}{(n+1)} \end{pmatrix} \end{aligned}$$

Using Lemma 2, yields

$$\begin{aligned} \tau(A_n^{(m)}) &= 2^{mn+1} \times \frac{(n+1)^m}{mn+m+2} \times \frac{2(\frac{2a-n-3}{n+1} + m - 3)}{\frac{2a-n-3}{n+1} - 3} \times [T_m(\frac{\frac{2a-n-3}{n+1} - 1}{2}) - 1] \\ &= 2^{mn+1} \times (n+1)^m \times [T_m(\frac{n+2}{n+1}) - 1]. \end{aligned}$$

Using Equation (11), yields the result. \square

Definition 2. The pyramid graph $B_n^{(m)}$ is the graph created from the gear graph G_{m+1} with vertices $\{u_0; u_1, u_2, \dots, u_m; w_1, w_2, \dots, w_m\}$ with double internal edges and m sets of vertices, say, $\{v_1^1, v_2^1, \dots, v_n^1\}, \{v_1^2, v_2^2, \dots, v_n^2\}, \dots, \{v_1^m, v_2^m, \dots, v_n^m\}$, such that for all $i = 1, 2, \dots, n$ the vertex v_i^j is adjacent to u_j and u_{j+1} , where $j = 1, 2, \dots, m-1$, and v_i^m is adjacent to u_1 and u_m . See Figure 2.

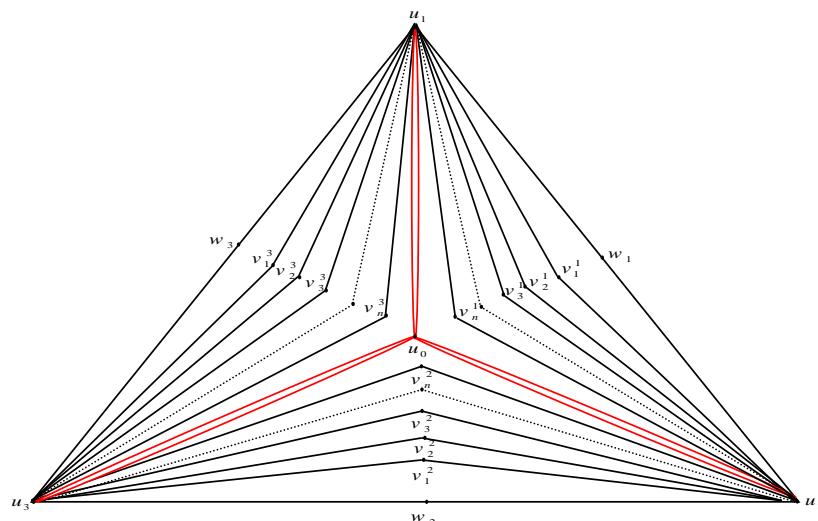


Figure 2. The pyramid graph $B_n^{(3)}$.

Theorem 2. For $n \geq 0$, $m \geq 3$, $\tau(B_n^{(m)}) = 2^{mn}[(n+3+2\sqrt{n+2})^m + (n+3-2\sqrt{n+2})^m - 2(n+1)^m]$.

Proof. Using Lemma 1, we get:

$$\tau(B_n^{(m)}) = \frac{1}{(mn+2m+1)^2} \times \det((mn+2m+1)I - D^c + A^c) = \frac{1}{(mn+2m+1)^2} \times$$

	$(2m+1)$	-1	-1	-1	1	1	1	1			
	-1	$(2n+5)$	1	1	0	1	1	0	0	...	0	1	...	1	...	1	1	...	0		
	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$:	:	:	0	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$	1	0	...	0	0	...	0	1	...	1	...	1	1	...	1			
	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$:	:	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$	1	1	...	1	0	...	0	0	...	0	1	...	1	1	...	1		
	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$:	:	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$	1	1	...	1	1	...	1	0	...	0	0	...	0	1	...	1		
	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$:	:	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$	1	1	...	1	1	...	1	0	...	0	0	...	0	1	...	1		
	-1	1	1	$(2n+5)$	1	1	0	0	1	...	1	1	...	1	...	1	0	...	0	0		
	1	0	0	1	...	1	3	1	1	1	...	1	1	...	1	...	1	...	1	1	...	1		
	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$:	:	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$	1	...	1	1	...	1	1	...	1	...	1	1	...	1				
	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$:	:	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$	1	...	1	1	...	1	1	...	1	...	1	1	...	1				
	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$:	:	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$	1	...	1	1	...	1	1	...	1	...	1	1	...	1				
	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$:	:	:	1	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \alpha_5$	1	...	1	1	...	1	1	...	1	...	1	1	...	1				
	1	0	1	...	1	0	0	1	1	3	1	...	1	1	...	1	...	1	1	...	1			
	1	0	0	1	...	1	0	1	1	...	1	1	3	1	...	1	1	...	1	...	1	1	...	1		
	:	1	0	1	...	1	0	1	1	...	1	1	1	3	1	...	1	1	...	1	...	1	1	...	1	
	det	0	0	1	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	1	
	1	0	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	1	
	:	1	0	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	1
	1	0	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	1	
	1	1	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	1	
	:	1	1	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	1
	1	1	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	1	
	1	1	1	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	
	0	1	1	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	
	1	1	1	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	1	
	0	1	1	1	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	
	1	0	1	1	0	0	0	1	0	1	1	0	1	1	1	3	1	0	1	1	0	1	1	1	0	

Let $j = (1 \cdots 1)$ be the $1 \times n$ matrix with all one, and J_n be the $n \times n$ matrix with all one. Set $a = 2n + 5$ and $b = mn + 2m + 1$. Then we get:

$$= \frac{1}{b} \times \det \begin{pmatrix} (a+1) & 2 & \dots & \dots & \dots & 2 & -1 & 0 & \dots & \dots & 0 & -1 & -j & 0 & \dots & \dots & 0 & -j \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & -1 & -1 & \ddots & \ddots & \ddots & 0 & -j & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 2 & \vdots & \ddots & \ddots & \ddots & -1 & 0 & \vdots & \ddots & \ddots & \ddots & 0 \\ 2 & \dots & \dots & \dots & \dots & 2 & (a+1) & 0 & \dots & \dots & 0 & -1 & -1 & 0 & \dots & \dots & 0 & -j & -j \\ 1 & 1 & 2 & \dots & \dots & 2 & 2 & 0 & \dots & \dots & 0 & 0 & 0 & \dots & \dots & \dots & 0 & 0 \\ 2 & 1 & \ddots & \ddots & \ddots & 2 & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & 1 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \vdots & \ddots & \ddots & \ddots & \vdots \\ 1 & 2 & \dots & \dots & 2 & 1 & 0 & \dots & \dots & 0 & 2 & 0 & \dots & \dots & \dots & 0 & 0 \\ j^t & j^t & 2j^t & \dots & \dots & 2j^t & 0 & \dots & \dots & \dots & \dots & 0 & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 2j^t & j^t & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 2j^t & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 2j^t & \ddots & \ddots & \ddots & \ddots & \vdots & j^t & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ j^t & 2j^t & \dots & \dots & 2j^t & j^t & 0 & \dots & \dots & \dots & \dots & 0 & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{pmatrix} 2I_{mn}$$

Using Lemma 3, yields

$$\tau(B_n^{(m)}) = \frac{1}{b} \times \det \begin{pmatrix} A & B \\ C & 2I_{mn} \end{pmatrix} = \frac{1}{b} \times \det(A - B \frac{1}{2I_{mn}} C) \times 2^{mn}$$

$$= \frac{1}{b} 2^{mn} \times 2^{-2m} \times \det \begin{pmatrix} (2a+2n+2) & 3n+4 & 4(n+1) & \dots & 4(n+1) & 3n+4 & -2 & 0 & \dots & \dots & 0 & -2 \\ 3n+4 & (2a+2n+2) & 3n+4 & 4(n+1) & \dots & 4(n+1) & -2 & \ddots & \ddots & \ddots & \ddots & 0 \\ 4(n+1) & 3n+4 & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 4(n+1) & \ddots & \ddots & \ddots & \ddots & \vdots \\ 4(n+1) & \ddots & \ddots & \ddots & \ddots & \vdots & 3n+4 & \ddots & \ddots & \ddots & \ddots & 0 \\ 3n+4 & 4(n+1) & \dots & 4(n+1) & 3n+4 & (2a+2n+2) & 0 & \dots & \dots & 0 & -2 & -2 \\ 2 & 2 & 4 & \dots & \dots & 4 & 4 & 0 & \dots & \dots & \dots & 0 \\ 4 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 4 & \ddots & \ddots & \ddots & \ddots & \vdots & 4 & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2 & 4 & \dots & \dots & 4 & 2 & 0 & \dots & \dots & \dots & 0 & 4 \end{pmatrix}$$

Using Lemma 3 again, yields

$$\tau(B_n^{(m)}) = \frac{2^{mn-2m}}{b} \times \det \begin{pmatrix} D & E \\ F & 4I_m \end{pmatrix} = \frac{2^{mn}}{b} \times \det(D - E \frac{1}{4I_m} F)$$

$$\tau(B_n^{(m)}) = \frac{2^{mn}}{b} \times \det \begin{pmatrix} (2a+2n+4) & (3n+7) & 4(n+2) & \dots & 4(n+2) & (3n+7) \\ (3n+7) & (2a+2n+4) & (3n+7) & \ddots & \dots & 4(n+2) \\ 4(n+2) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 4(n+2) & \ddots & \ddots & \ddots & \ddots & (3n+7) \\ (3n+7) & 4(n+2) & \dots & 4(n+2) & (3n+7) & (2a+2n+4) \end{pmatrix}$$

With a straightforward inducement using properties of determinants, we obtain

$$\tau(B_n^{(m)}) = \frac{2^{mn}}{b} \times \frac{4b}{m+n+4} \times \det \begin{pmatrix} (2a-n-3) & 0 & (n+1) & \dots & (n+1) & 0 \\ 0 & (2a-n-3) & 0 & \ddots & \dots & (n+1) \\ (n+1) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & (n+1) \\ (n+1) & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & (n+1) & \dots & (n+1) & 0 & (2a-n-3) \end{pmatrix}$$

$$= \frac{2^{mn+2} \times (n+1)^m}{mn+m+4} \times \det \begin{pmatrix} \frac{(2a-n-3)}{(n+1)} & 0 & 1 & \cdots & 1 & 0 \\ 0 & \frac{(2a-n-3)}{(n+1)} & 0 & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & \frac{(2a-n-3)}{(n+1)} \end{pmatrix}$$

Using Lemma 2, yields

$$\begin{aligned} \tau(B_n^{(m)}) &= 2^{mn+2} \times \frac{(n+1)^m}{mn+m+4} \times \frac{2(\frac{2a-n-3}{n+1} + m - 3)}{\frac{2a-n-3}{n+1} - 3} \times [T_m(\frac{\frac{2a-n-3}{n+1} - 1}{2}) - 1] \\ &= 2^{mn+1} \times (n+1)^m \times [T_m(\frac{n+3}{n+1}) - 1]. \end{aligned}$$

Using Equation (11), yields the result. \square

Definition 3. The pyramid graph $C_n^{(m)}$ is the graph created from the gear graph G_{m+1} with vertices $\{u_0; u_1, u_2, \dots, u_m; w_1, w_2, \dots, w_m\}$ with double external edges and m sets of vertices, say, $\{v_1^1, v_2^1, \dots, v_n^1\}, \{v_1^2, u_2^2, \dots, v_n^2\}, \dots, \{v_1^m, v_2^m, \dots, v_n^m\}$, such that for all $i = 1, 2, \dots, n$ the vertex v_i^j is adjacent to u_j and u_{j+1} , where $j = 1, 2, \dots, m-1$, and v_i^m is adjacent to u_1 and u_m . See Figure 3.

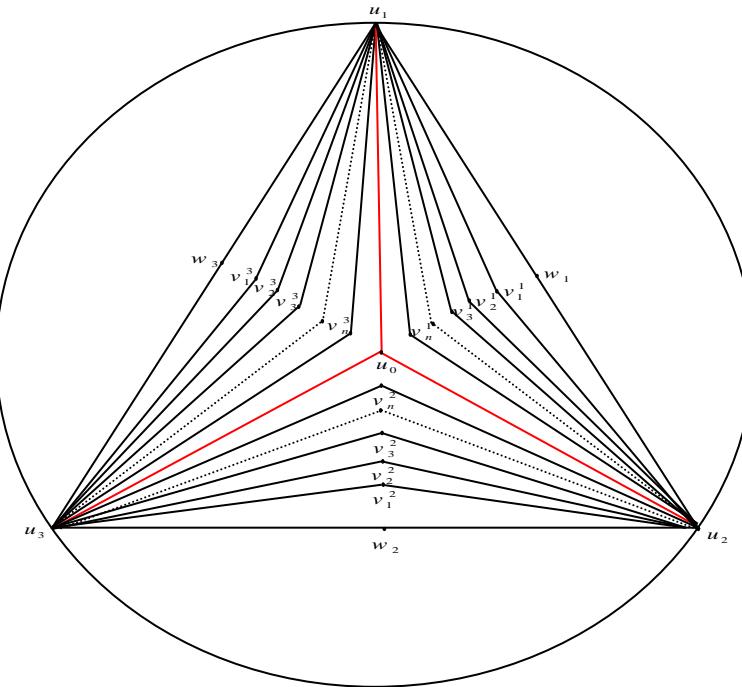


Figure 3. The pyramid graph $C_n^{(3)}$.

Theorem 3. For $n \geq 0, m \geq 3$, $\tau(C_n^{(m)}) = 2^{mn}[(n+4+\sqrt{2n+7})^m + (n+4-\sqrt{2n+7})^m - 2(n+3)^m]$.

Proof. Using Lemma 1, we have:

$$\tau(C_n^{(m)}) = \frac{1}{(mn+2m+1)^2} \times \det((mn+2m+1)I - D^c + A^c) = \frac{1}{(mn+2m+1)^2} \times$$

Let $j = (1 \cdots 1)$ be the $1 \times n$ matrix with all one, and J_n be the $n \times n$ matrix with all one. Set $a = 2n + 6$ and $b = mn + 2m + 1$. Then we have:

$$\tau \left(C_n^{(m)} \right) = \frac{1}{l^2} \times \det \left(\begin{array}{ccccccccc} m+1 & 0 & \cdots & \cdots & \cdots & 0 & 1 & \cdots & \cdots & \cdots & 1 & j & \cdots & \cdots & \cdots & j \\ 0 & a & 0 & 1 & \cdots & 1 & 0 & 0 & 1 & \cdots & \cdots & 1 & 0 & 0 & j & \cdots & \cdots & j & 0 \\ \vdots & 0 & \ddots & \ddots & \ddots & \ddots & 1 & 0 & 0 & \ddots & \ddots & \ddots & 1 & 0 & \ddots & \ddots & \ddots & j \\ \vdots & 1 & \ddots & \ddots & \ddots & \ddots & \vdots & 1 & 0 & \ddots & \ddots & \ddots & \vdots & j & \ddots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & 1 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & 0 & 0 & 1 & \vdots & \ddots & \ddots & \ddots & j \\ 0 & 0 & 1 & \cdots & 1 & 0 & a & 1 & \cdots & \cdots & 1 & 0 & 0 & j & \cdots & \cdots & j & 0 & 0 \\ 1 & 0 & 0 & 1 & \cdots & \cdots & 1 & 3 & 1 & \cdots & \cdots & \cdots & 1 & j & \cdots & \cdots & \cdots & j \\ \vdots & 1 & 0 & \ddots & \ddots & \ddots & 1 & 1 & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & 1 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \vdots & 1 & \vdots & \ddots & \ddots & \ddots & \vdots \\ 1 & 0 & 1 & \cdots & \cdots & 1 & 0 & 1 & \cdots & \cdots & \cdots & 1 & 3 & j & \cdots & \cdots & \cdots & j \\ j^t & 0 & 0 & j^t & \cdots & \cdots & j^t & j^t & \cdots & \cdots & \cdots & \cdots & j^t & \cdots & \cdots & \cdots & \cdots & \vdots \\ \vdots & j^t & 0 & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & j^t & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ j^t & 0 & j^t & \cdots & \cdots & j^t & 0 & j^t & \cdots & \cdots & \cdots & \cdots & j^t & \cdots & \cdots & \cdots & \cdots & \vdots \end{array} \right) - 2I_{mn} + J_{mn}$$

$$= \frac{1}{b} \times \det \left(\begin{array}{ccccccccccccc} a & 0 & 1 & \cdots & 1 & 0 & -1 & 0 & \cdots & \cdots & 0 & -1 & -j & 0 & \cdots & \cdots & 0 & -j \\ 0 & \ddots & \ddots & \ddots & 1 & -1 & -1 & \ddots & \ddots & \ddots & 0 & -j & \ddots & \ddots & \ddots & \ddots & 0 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & -1 & 0 & \vdots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & a & 0 & \cdots & \cdots & 0 & -1 & -1 & 0 & \cdots & \cdots & 0 & -j & -j \\ 0 & 0 & 1 & \cdots & 1 & 2 & 0 & \cdots & \cdots & 0 & 0 & 0 & \cdots & \cdots & \cdots & \cdots & 0 \\ 1 & 0 & \ddots & \ddots & \ddots & 1 & 0 & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & 0 & \vdots & \ddots & \ddots & \vdots \\ 0 & 1 & \cdots & 1 & 0 & 0 & 0 & \cdots & \cdots & 0 & 2 & 0 & \cdots & \cdots & \cdots & \cdots & 0 \\ 0 & 0 & j^t & \cdots & j^t & 0 & 0 & \cdots & \cdots & \cdots & 0 \\ j^t & 0 & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & j^t & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ j^t & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & j^t & \cdots & \cdots & j^t & 0 & 0 & \cdots & \cdots & \cdots & 0 \end{array} \right) \\ 2I_{mn}$$

Using Lemma 3, yields

$$\tau(C_n^{(m)}) = \frac{1}{b} \times \det \left(\begin{array}{cc} A & B \\ C & 2I_{mn} \end{array} \right) = \frac{1}{b} \times \det(A - B \frac{1}{2I_{mn}} C) \times 2^{mn}$$

$$= \frac{1}{b} \times 2^{mn} \times 2^{-2m} \times \det \left(\begin{array}{ccccccccccccc} 2a & n & 2(n+1) & \cdots & 2(n+1) & n & -2 & 0 & \cdots & \cdots & 0 & -2 \\ n & 2a & n+2 & 2(n+1) & \cdots & 2(n+1) & -2 & \ddots & \ddots & \ddots & \ddots & 0 \\ 2(n+1) & n & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 2(n+1) & \vdots & \ddots & \ddots & \ddots & \vdots \\ 2(n+1) & \ddots & \ddots & \ddots & \ddots & \vdots & n & \vdots & \ddots & \ddots & \ddots & 0 \\ n & 2(n+1) & \cdots & 2(n+1) & n & 2a & 0 & \cdots & \cdots & 0 & -2 & -2 \\ 0 & 0 & 2 & \cdots & \cdots & 2 & 4 & 0 & \cdots & \cdots & \cdots & 0 \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 2 & \vdots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \vdots & \ddots & \ddots & \ddots & 0 \\ 0 & 2 & \cdots & \cdots & 2 & 0 & 0 & \cdots & \cdots & \cdots & 0 & 4 \end{array} \right)$$

Using Lemma 3 again, yields

$$\tau(C_n^{(m)}) = \frac{2^{mn-2m}}{b} \times \det \left(\begin{array}{cc} D & E \\ F & 4I_m \end{array} \right) = \frac{2^{mn}}{b} \times \det(D - E \frac{1}{4I_m} F)$$

$$\tau(C_n^{(m)}) = \frac{2^{mn}}{b} \times \det \left(\begin{array}{ccccccccccccc} 2a & (n+1) & 2(n+2) & \cdots & 2(n+2) & (n+1) \\ (n+1) & 2a & (n+3) & \ddots & \cdots & 2(n+2) \\ 2(n+2) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2(n+2) \\ 2(n+2) & \ddots & \ddots & \ddots & \ddots & (n+1) \\ (n+1) & 2(n+2) & \cdots & 2(n+2) & (n+1) & 2a \end{array} \right)$$

Using properties of determinants, we have:

$$\tau(C_n^{(m)}) = \frac{2^{mn}}{b} \times \frac{2b}{mn+3m+2} \times \det \begin{pmatrix} (2a-n-1) & 0 & (n+3) & \cdots & (n+3) & 0 \\ 0 & (2a-n-1) & 0 & \ddots & \cdots & (n+3) \\ (n+3) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & (n+3) \\ (n+3) & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & (n+3) & \cdots & (n+3) & 0 & (2a-n-1) \end{pmatrix}$$

$$= \frac{2^{mn+1}(n+3)^m}{mn+3m+2} \times \det \begin{pmatrix} \frac{(2a-n-1)}{(n+3)} & 0 & 1 & \cdots & 1 & 0 \\ 0 & \frac{(2a-n-1)}{(n+3)} & 0 & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & \frac{(2a-n-1)}{(n+3)} \end{pmatrix}$$

Using Lemma 2, yields:

$$\begin{aligned} \tau(C_n^{(m)}) &= 2^{mn+1} \times \frac{(n+3)^m}{mn+3m+2} \times \frac{2(\frac{2a-n-1}{n+3} + m - 3)}{\frac{2a-n-1}{n+3} - 3} \times [T_m(\frac{2a-n-1}{2} - 1)] \\ &= 2^{mn+1} \times (n+3)^m \times [T_m(\frac{n+4}{n+3}) - 1]. \end{aligned}$$

Using Equation (11), yields the result. \square

Definition 4. The pyramid graph $D_n^{(m)}$ is the graph created from the gear graph G_{m+1} with vertices $\{u_0; u_1, u_2, \dots, u_m; w_1, w_2, \dots, w_m\}$ with double internal and external edges and m sets of vertices, say, $\{v_1^1, v_2^1, \dots, v_n^1\}, \{v_1^2, v_2^2, \dots, v_n^2\}, \dots, \{v_1^m, v_2^m, \dots, v_n^m\}$, such that for all $i = 1, 2, \dots, n$ the vertex v_i^j is adjacent to u_j and u_{j+1} , where $j = 1, 2, \dots, m-1$, and v_i^m is adjacent to u_1 and u_m . See Figure 4.

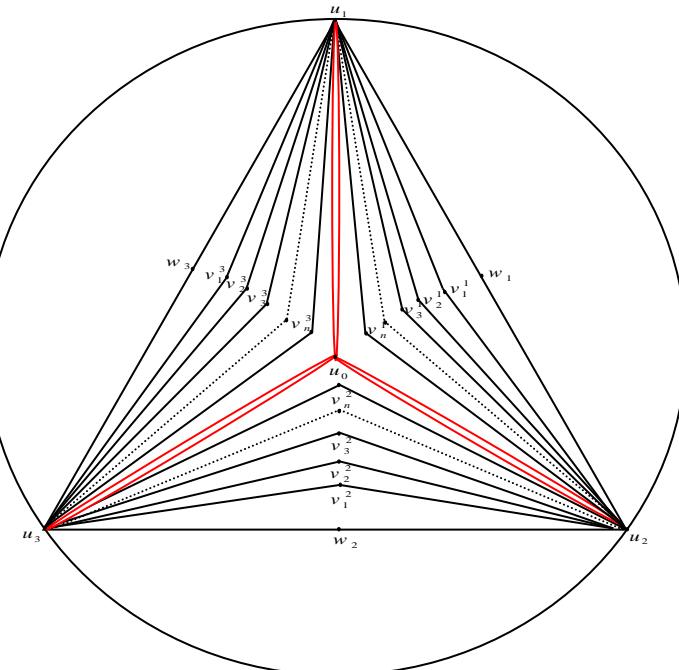


Figure 4. The pyramid graph $D_n^{(3)}$.

Theorem 4. For $n \geq 0$, $m \geq 3$, $\tau(D_n^{(m)}) = 2^{mn}[(n+5+2\sqrt{n+4})^m + (n+5-2\sqrt{n+4})^m - 2(n+3)^m]$.

Proof. Applying Lemma 1, we have:

$$\tau(D_n^{(m)}) = \frac{1}{(mn+2m+1)^2} \times \det((mn+2m+1)I - D^c + A^c) = \frac{1}{(mn+2m+1)^2} \times$$

Let $j = (1 \cdots 1)$ be the $1 \times n$ matrix with all one, and J_n the $n \times n$ matrix with all one. Set $a = 2n + 7$ and $b = mn + 2m + 1$. Then we have:

$$= \frac{1}{b} \det \left(\begin{array}{ccccccccc} (a+1) & 1 & 2 & \cdots & 2 & 1 & -1 & 0 & \cdots & \cdots & 0 & -1 & -j & 0 & \cdots & \cdots & 0 & -j \\ 1 & \ddots & \ddots & \ddots & \ddots & 2 & -1 & -1 & \ddots & \ddots & \ddots & 0 & -j & \ddots & \ddots & \ddots & 0 \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & -1 & 0 & \vdots & \ddots & \ddots & \ddots & 0 \\ 1 & 2 & \cdots & 2 & 1 & (a+1) & 0 & \cdots & 0 & -1 & -1 & 0 & \cdots & \cdots & 0 & -j & -j \\ 1 & 1 & 2 & \cdots & \cdots & 2 & 2 & 0 & \cdots & \cdots & 0 & 0 & \cdots & \cdots & \cdots & \cdots & 0 \\ 2 & 1 & \ddots & \ddots & \ddots & 2 & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \vdots \\ 1 & 2 & \cdots & 2 & 1 & 0 & \cdots & \cdots & 0 & 2 & 0 & \cdots & \cdots & \cdots & \cdots & 0 \\ j^t & j^t & 2j^t & \cdots & \cdots & 2j^t & 0 & \cdots & \cdots & \cdots & 0 & \vdots & \vdots & \ddots & \ddots & \vdots \\ 2j^t & j^t & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2j^t & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ 2j^t & \ddots & \ddots & \ddots & \ddots & j^t & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ j^t & 2j^t & \cdots & \cdots & 2j^t & j^t & 0 & \cdots & \cdots & \cdots & 0 & \vdots & \vdots & \ddots & \ddots & \vdots \end{array} \right)$$

Using Lemma 3, yields

$$\tau(D_n^{(m)}) = \frac{1}{b} \times \det \begin{pmatrix} A & B \\ C & 2I_{mn} \end{pmatrix} = \frac{1}{b} \times \det(A - B \frac{1}{2I_{mn}} C) \times 2^{mn}$$

$$= \frac{1}{b} 2^{mn} \times 2^{-2m} \times \det \left(\begin{array}{ccccccccc} (2a+2n+2) & 3n+2 & 4(n+1) & \cdots & 4(n+1) & 3n+2 & -2 & 0 & \cdots & \cdots & 0 & -2 \\ 3n+2 & (2a+2n+2) & 3n+2 & 4(n+1) & \cdots & 4(n+1) & -2 & \ddots & \ddots & \ddots & \ddots & 0 \\ 4(n+1) & 3n+4 & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 4(n+1) & \ddots & \ddots & \ddots & \ddots & 3n+2 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 3n+2 & 4(n+1) & \cdots & 4(n+1) & 3n+2 & (2a+2n+2) & 0 & \cdots & \cdots & 0 & -2 & -2 \\ 2 & 2 & 4 & \cdots & \cdots & 4 & 4 & 0 & \cdots & \cdots & \cdots & 0 \\ 4 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 4 & \ddots & \ddots & \ddots & \ddots & 4 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2 & 4 & \cdots & \cdots & 4 & 2 & 0 & \cdots & \cdots & \cdots & 0 & 4 \end{array} \right)$$

Using Lemma 3, yields

$$\tau(D_n^{(m)}) = \frac{2^{mn-2m}}{b} \times \det \begin{pmatrix} A & B \\ C & 4I_m \end{pmatrix} = \frac{2^{mn}}{b} \times \det(A - B \frac{1}{4I_m} C)$$

$$\tau(D_n^{(m)}) = \frac{2^{mn}}{b} \times \det \left(\begin{array}{cccccc} (2a+2n+4) & (3n+5) & 4(n+2) & \cdots & 4(n+2) & (3n+5) \\ (3n+5) & (2a+2n+4) & (3n+5) & \ddots & \cdots & 4(n+2) \\ 4(n+2) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 4(n+2) \\ 4(n+2) & \ddots & \ddots & \ddots & \ddots & (3n+5) \\ (3n+5) & 4(n+2) & \cdots & 4(n+2) & (3n+5) & (2a+2n+4) \end{array} \right)$$

Straightforward inducement using properties of determinants, we get:

$$\tau(D_n^{(m)}) = \frac{2^{mn}}{b} \times \frac{4b}{mn+3m+4} \times \det \left(\begin{array}{cccccc} (2a-n-1) & 0 & (n+3) & \cdots & (n+3) & 0 \\ 0 & (2a-n-1) & 0 & \ddots & \cdots & (n+3) \\ (n+3) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & (n+3) \\ (n+3) & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & (n+3) & \cdots & (n+3) & 0 & (2a-n-1) \end{array} \right)$$

$$= \frac{2^{mn+2} (n+3)^m}{mn+3m+4} \times \det \begin{pmatrix} \frac{(2a-n-1)}{(n+3)} & 0 & 1 & \cdots & 1 & 0 \\ 0 & \frac{(2a-n-1)}{(n+3)} & 0 & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & \frac{(2a-n-1)}{(n+3)} \end{pmatrix}$$

Using Lemma 2, yields:

$$\tau(D_n^{(m)}) = 2^{mn+2} \times \frac{(n+3)^m}{mn+3m+4} \times \frac{2(\frac{2a-n-1}{n+3} + m - 3)}{\frac{2a-n-1}{n+3} - 3} \times [T_m(\frac{2a-n-1}{2}) - 1] = 2^{mn+1} \times (n+3)^m \times [T_m(\frac{n+5}{n+3}) - 1].$$

Using Equation (11), yields the result. \square

4. Numerical Results

The following Table 1 illustrates some values of the number of spanning trees of studied pyramid graphs.

Table 1. Some values of the number of spanning trees of studied pyramid graphs.

m	n	$\tau(P_n^{(m)})$	$\tau(A_n^{(m)})$	$\tau(B_n^{(m)})$	$\tau(C_n^{(m)})$
3	0	50	196	242	676
3	1	1024	3200	3136	8192
3	2	15,488	43,264	36,992	92,416
3	3	200,704	524,288	409,600	991,232
3	4	2,367,488	5,914,624	4,333,568	10,240,000
3	5	26,214,400	63,438,848	44,302,336	102,760,448
4	0	192	1152	1792	6400
4	1	11,520	49,152	57,600	184,320
4	2	458,752	1,638,400	1,622,016	4,816,896
4	3	14,745,600	47,185,920	41,746,432	117,440,512
4	4	415,236,096	1,233,125,376	1,006,632,960	2,717,908,992
4	5	10,687,086,592	30,064,771,072	23,102,226,432	60,397,977,600
5	0	722	6724	12,482	58,564
5	1	123,904	739,328	984,064	3,964,928
5	2	12,781,568	59,969,536	65,619,968	237,899,776
5	3	1,007,681,536	4,060,086,272	3,901,751,296	13,088,325,632
5	4	67,194,847,232	243,609,370,624	213,408,284,672	674,448,277,504
5	5	3,995,393,327,104	243,609,370,624	10,953,240,346,624	33,019,708,571,648

5. Conclusions

The number of spanning trees $\tau(G)$ in graphs (networks) is an important invariant. The computation of this number is not only beneficial from a mathematical (computational) standpoint, but it is also an important measure of the reliability of a network and electrical circuit layout. Some computationally laborious problems such as the traveling salesman problem can be resolved approximately by using spanning trees. In this paper, we define some classes of pyramid graphs created by a gear graph and we have studied the problem of computing the number of spanning trees of these graphs.

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