



# Article A Graph-Based Power Flow Method for Balanced Distribution Systems

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**Abstract:** A power flow method based on graph theory is presented for three-phase balanced distribution systems. The graph theory is used to describe the power network and facilitate the derivation of the relationship between bus <u>C</u>urrents and the bus <u>V</u>oltage <u>B</u>ias from the feeder bus (the CVB equation). A distinctive feature of the CVB equation is its unified form for both radial and meshed networks. The method requires neither a tricky numbering and layering of nodes nor breaking meshes and loop-analysis, which are both necessary in previous works for meshed networks. The convergence of the proposed method is proven using the Banach fixed-point theorem.

**Keywords:** distribution system; power flow; graph theory; radial network; meshed network; backward/forward sweep

# 1. Introduction

Power flow calculation is the most fundamental numerical problem for power system analysis. A fast and general power flow method will be required by distribution systems as the development of smart grid and must be as efficient as possible in the future [1]. Methods on transmission networks are well developed such as Gauss–Seidel, Newton–Raphson [2] and Fast-Decoupled method [3]. Distribution networks have some special characteristics such as radial/weakly meshed structure, high R/X ratios of impedances, large number of branches and nodes, etc. These features may cause problems when the algorithms for power flow of transmission networks are applied to distribution systems [4]. Power flow calculations may be executed every five minutes on traditional networks, but as for microgrids, it may not meets the requirements. Microgrids have the natures of uncertainty and volatility, so they need real-time monitoring to guarantee their reliability, and require a faster power flow method. Power flow method is also a very important tool for improving the reliability and efficiency of fault analysis [5], and it can also provide evidence for protection for power distribution systems.

Backward/forward sweep (BFS) method ,which is intended to solve unbalanced radial distribution networks, has a very good performance, where all nodes are labeled into different layers according to distances from the feeder node. The branch currents are calculated in the backward sweep, while the bus voltage is then calculated in the forward sweep [6]. However, the BFS method cannot be applied directly to networks even with weakly meshed structure because the distances from the feeder node are not unique in the presence of loops. Shirmohammadi et al. have proposed a compensation-based power flow method for solving weakly meshed networks by using the multi-port compensation technique and basic formulations of Kirchhoff's laws [7]. Teng has proposed a direct method for both radial and meshed networks by developing the bus-injection to branch-current

(BIBC) matrix and the branch-current to bus-voltage (BCBV) matrix. However, when solving meshed networks, the method has to apply some preliminary operations, including Kron's Reduction and modifying the two matrices by loop-analysis [8]. Wu and Zhang developed a power flow method for dealing with meshed network based on compensation and loop-analysis [9]. These methods all need extra processing for meshed networks.

This paper proposes a graph-based power flow method for distribution systems, which has a unified formulation for both radial and meshed networks. Compared with previous works, the method uses graph theory to directly build the CVB equation, a map to the bus currents from the bus voltage bias from the feeder node. It requires neither a tricky numbering and layering of nodes nor breaking meshes and loop-analysis, which are both necessary in previous works for meshed networks. Although graph theory has been used for power systems in many aspects [10–12], there are a few in power flow calculation. The most relevant is the work published more recently [13], where graph theory is used for building the Z-bus matrix, and the results obtained are only for radial distribution systems. The other contribution of this paper is that the convergence of the graph-based method is addressed by using the Banach fixed-point theorem, associated with the convergence rate of a clear physical meaning.

*Notations:* The following notations are used in this paper:

- CVB: The equation between the bus currents and the bus voltage bias from the feeder bus.
- BFS: The backward/forward sweep method.
- BIBC, BCBV: The bus-injection to branch-current matrix and the branch-current to bus-voltage matrix in [8].
- $\mathbb{R}, \mathbb{C}$  denote the rational, complex number sets.
- \* denotes the conjugate operator of complex number.
- *I*, *U*: The inject current and voltage vectors of all nodes including the feeder node.
- *I<sub>s</sub>*, *U<sub>s</sub>*: The inject current and voltage vectors of all nodes except the feeder node.
- *U<sub>L</sub>*, *I<sub>L</sub>*: The voltage drop, impedance, current vectors of all branches.
- $Z_l, S_s$ : Matrices of all impedance of branches/complex power of nodes except feeder node.
- *H*, *H*<sub>s</sub>: The incidence matrix with/without the row of feeder node.
- $U_d$ : The voltage difference vector between the feeder node and other nodes.
- $\Phi$ ,  $Z_d$ : The mapping between  $U_d$  and  $I_s$ .
- **1**<sub>*n*</sub>: A *n*-order vector with all elements being 1.

# 2. Problem Formulation

# 2.1. Topological Description of the Network

A distribution network has a typical tree structure, the root of which is the feeder node with a known voltage. Sometimes there are some extra branches between nodes so as to form a meshed structure.

We use undirected graph  $\mathcal{G} = \{\mathcal{N}, \mathcal{E}\}$  to depict the topology structure of a given distribution network where  $\mathcal{N} = \{1, \dots, n\}$  and  $\mathcal{E} = \{L_1, \dots, L_m\}$  are node set and branch set, respectively. If there is a branch between two nodes in a graph, the two nodes are connected. Although incidence matrix has been used in previous studies [14,15], it has to be numbered from front to back in sequence. By using of graph theory, we can give an arbitrary numbering to nodes and branches. Figure 1 shows two simple typical distribution network containing 4 nodes, where the feeder node is node 3, not the first node. Without loss of generality, the positive direction of branch current is defined to be always flowing out of the node with the lower number. In this setting, the incidence matrix  $H = (h_{ij})$  of the graph  $\mathcal{G}$  is defined as

$$h_{ij} = \begin{cases} +1, \text{ the branch current } L_j \text{ starts at node } i \\ -1, \text{ the branch current } L_j \text{ ends at node } i \\ 0, \text{ otherwise} \end{cases}$$

The incidence matrices of networks in Figure 1, for instance, are given as

$$H = \begin{bmatrix} L_1 & L_2 & L_3 & & L_1 & L_2 & L_3 & L_4 & L_5 \\ 1 & 1 & 1 & 1 \\ 0 & -1 & 0 \\ -1 & 0 & 0 \\ 4 & 0 & 0 & -1 \end{bmatrix}, \qquad \begin{bmatrix} L_1 & L_2 & L_3 & L_4 & L_5 \\ 1 & 1 & 1 & 0 & 0 \\ 0 & -1 & 0 & 1 & 1 \\ -1 & 0 & 0 & -1 & 0 \\ 0 & 0 & -1 & 0 & -1 \end{bmatrix}.$$

Specially, m = n - 1 for radial networks. It is rational to assume that the considered network is connected, which implies that the rank of *H* is n - 1.



**Figure 1.** A typical radial network and meshed network: (a) radial structure (n = 4, m = 3); and (b) meshed structure (n = 4, m = 5).

### 2.2. Basic Formulations of Kirchhoff's Laws

The power flow is related to the steady-state behavior of the power systems, where all the voltages and currents are sinusoidal signals with the same frequency. For a three-phase balanced distribution system, each signal can be represented by a complex value. Without loss of generality, all the electrical variables are complex numbers in this paper if not specifically stated.

In this sense, the current, voltage and complex power of node k are denoted by complex numbers  $i_k$ ,  $u_k$  and  $s_k$ , respectively. The positive direction of  $i_k$  injects to node k, as illustrated in Figure 1. The concatenated current and voltage vectors are denoted by  $I = [i_1, i_2, \dots, i_n]^T \in \mathbb{C}^n$  and  $U = [u_1, u_2, \dots, u_n]^T \in \mathbb{C}^n$ . Let f be the number of feeder node. The feeding power of distribution networks can be given by  $s_f = u_f i_f^*$ , where \* denotes the conjugate operator.

Let  $i_{Lk}$ ,  $u_{Lk}$  and  $z_{Lk}$  be the current, voltage and impedance of branch k, respectively. The positive directions of  $i_{Lk}$  and  $u_{Lk}$  follow that of  $h_{ij}$ , i.e., from the node with lower number to the node with larger number. Similarly  $I_L = [i_{L_1}, i_{L_2}, \cdots, i_{L_m}]^T \in \mathbb{C}^m$  and  $U_L = [u_{L_1}, u_{L_2}, \cdots, u_{L_m}]^T \in \mathbb{C}^m$ . The diagonal impedance matrix is defined by  $Z_L = \text{diag}(z_{L_1}, z_{L_2}, \cdots, z_{L_m}) \in \mathbb{C}^{m \times m}$ .

Clearly one has

$$I = \left[\frac{s_1^*}{u_1^*}, \frac{s_2^*}{u_2^*}, \dots, \frac{s_n^*}{u_n^*}\right]^T,$$
(1)

and

$$U_L = Z_L I_L. \tag{2}$$

The Kirchhoff's Current and Voltage Laws can be conveniently described by use of incidence matrix, respectively,

$$I = HI_L, \quad U_L = H^T U, \tag{3}$$

and they are valid for both radial and meshed networks. Combining with Equation (2), it follows that

$$I = HZ_I^{-1}H^T U. (4)$$

#### 2.3. Reformulation of Power Flow Equations

Equations (1) and (4) form the power flow equations. However, matrix  $(HZ_L^{-1}H^T)$  is singular, which hampers constructing an identity map from them (an identity map is a function that always returns the same value that was used as its argument.).

For a distribution network, the feeder node is a slack node, whose voltage is fixed as the base value  $V_0 \in \mathcal{R}$ . This paper addresses the case that all the other buses except for the feeder bus are modeled as *P*, *Q* bus. Such a constant power case is increasingly common in the modern distribution systems because more and more power electronics devices are used.

Given the feeder node voltage  $V_0$  and the complex power  $s_k$  for other nodes  $k \in \mathcal{N} \setminus f$ , where notation  $\mathcal{N} \setminus f$  denotes the subset of  $\mathcal{N}$  deleting the element f, the goal of power flow is to calculate the currents and voltages of nodes in the set  $\mathcal{N} \setminus f$ , which we use  $I_s, U_s \in \mathbb{C}^{n-1}$  to denote respectively, that is,

$$I_{s} = [i_{1}, \cdots, i_{f-1}, i_{f+1}, \cdots, i_{n}]^{T}, \quad U_{s} = [u_{1}, \cdots, u_{f-1}, u_{f+1}, \cdots, u_{n}]^{T}.$$
(5)

Correspondingly, let  $H_s \in \mathbb{R}^{(n-1) \times m}$  be the matrix removing the *f*-th row of *H*, with which

$$I_s = H_s I_L. \tag{6}$$

As for the examples in Figure 1, node 3 is the feeder node, then

$$H_{s} = \begin{pmatrix} L_{1} & L_{2} & L_{3} & L_{1} & L_{2} & L_{3} & L_{4} & L_{5} \\ 1 & 1 & 1 & 1 \\ 0 & -1 & 0 \\ 4 & 0 & 0 & -1 \end{pmatrix}, \begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 0 & -1 & 0 & 1 & 1 \\ 0 & 0 & -1 & 0 & -1 \end{bmatrix}.$$
(7)

Denote by  $U_d \in \mathbb{C}^{n-1}$  the voltage differences between the feeder node and other nodes, namely,

$$U_d = U_s - V_0 \mathbf{1}_{n-1}.$$
 (8)

Throughout of this paper, notation  $\mathbf{1}_n$  denotes a *n*-order vector with all elements being 1. Due to  $H^T \mathbf{1}_n = 0$ , the following can be obtained,

$$U_L = H^T U = H^T U - V_0 H^T \mathbf{1}_n = H_s^T U_d.$$
<sup>(9)</sup>

# 3. Main Results

#### 3.1. Review BFS Method

Traditional BFS method generally takes advantage of the radial topology. It starts with numbering and layering from the feeder node to terminal nodes. The backward sweep, starting from the terminal layer and ending at the first layer, is to calculate the branch currents  $I_L$  by a current summation with a possible voltage update. The forward sweep operating in an opposite direction is to calculate the voltage drop of nodes  $U_d$  with the branch currents obtained in the backward process.

Note that  $H_s$  is nonsingular in a connected radial network, we can directly obtain branch current  $I_L$  from  $I_s$  by Equation (6), instead of by current summations in the backward sweep. In our study,

the backward and forward processes can be described, based on graph theory, simply without layering as:

(s1) 
$$U_s^{(k)} \to I_s^{(k)} = \left[\frac{s_1^*}{u_1^*}, \dots, \frac{s_{f-1}^*}{u_{f-1}^*}, \frac{s_{f+1}^*}{u_{f+1}^*}, \dots, \frac{s_n^*}{u_n^*}\right]^T$$

(s2) 
$$I_L^{(k)} = H_s^{-1} I_s^{(k)}$$
,  
(s3)  $I_L^{(k)} - Z_L I^{(k)}$ 

$$(53) \quad U_L = Z_L I_L ,$$

$$(s4) \quad U_d^{(\gamma)} = H_s^{-1} U_L^{(\gamma)},$$

(s5)  $U_s^{(k+1)} = V_0 \mathbf{1}_{n-1} + U_d^{(k)},$  (11)

(s6) repeat (s1) to (s5) until 
$$||U_s^{(k+1)} - U_s^{(k)}|| < \epsilon$$

where superscript (*k*) denotes the values at the *k*-th iteration and  $\epsilon$  is the convergence tolerance.

# 3.2. Unified Method for both Radial and Meshed Networks

The above graph-based process is no longer applicable for meshed network because  $H_s$  is not a square matrix for meshed structure. The traditional methods cannot apply either in that the presence of circulating current prohibits layering nodes.

Generally, radial and meshed networks are dealt with separately when we consider the power flow for distribution systems. For dealing with meshed networks, breaking meshes or loop-analysis were needed in previous studies. Below, a uniform method for both radial and meshed networks is presented.

Combining Equations (6), (2), and (9) yields,

$$I_s = H_s Z_L^{-1} H_s^T U_d = \Phi U_d.$$

$$\tag{10}$$

Here,  $\Phi = H_s Z_L^{-1} H_s^T$  is nothing but the Laplacian matrix weighted by branch admittances of  $\mathcal{G}$  removing the row and column corresponding to the feeder node. For a connected network,  $\Phi$  is always nonsingular no matter if  $H_s$  is a square matrix. Define

$$Z_d := \Phi^{-1}. \tag{11}$$

Equation (10) builds a bijective mapping between  $I_s$  and  $U_d$ , by which the function of the steps (s2)–(s4) in BFS method can be compactly rewritten as

$$U_d^{(k)} = Z_d I_s^{(k)}, (12)$$

which is nothing but the CVB equation. Equations (10) and (12) look significant by themselves since they mean that the injected currents of nodes could be directly related not to node powers but to the node voltage bias from the feeder node.

Now, our uniform graph-based method now can be delivered as follows:

(g1) 
$$U_s^{(k)} \to I_s^{(k)} = \left[\frac{s_1^*}{u_1^*}, \dots, \frac{s_{f-1}^*}{u_{f-1}^*}, \frac{s_{f+1}^*}{u_{f+1}^*}, \dots, \frac{s_n^*}{u_n^*}\right]^T$$

(g2)  $U_d^{(k)} = Z_d I_s^{(k)}$ ,

(g3) 
$$U_s^{(k+1)} = V_0 \mathbf{1}_{n-1} + U_d^{(k)}$$
,

(g4) Repeat (g1) to (g3) until 
$$||U_s^{(k+1)} - U_s^{(k)}|| < \epsilon$$
.

The initial value  $U_s^{(0)}$  is generally set as  $V_0 \mathbf{1}_{n-1}$ . Note that the inversion of  $\Phi$ , i.e.,  $Z_d$  does not need to calculate during the iteration. The flowchart of graph-based method is shown in Figure 2.





Start

Load  $S_s, H_s, Z_L, \epsilon$ 

Figure 2. The flowchart of graph-based method.

# 3.3. Convergence of Method

Since the above algorithm is explicit about the involved electrical variables, its convergence can be analyzed by using Banach fixed-point theorem.

Let  $S_s = \text{diag}(s_1, \dots s_{f-1}, s_{f+1}, \dots, s_n)$ , the diagonal matrix of injected power for node set  $\mathcal{N} \setminus f$ . Denote by  $[x_i]_{\mathcal{N} \setminus f}$  a vector consisting of all  $x_i$  indexed by  $\mathcal{N}$  except for the *f*th one. Let  $V_s$  be the solution of power flow, i.e., the steady state of the algorithm, and  $\underline{v}$  be the element of  $V_s$  with the minimal magnitude, i.e.,  $\underline{v} = \min_{i \in \mathcal{N} \setminus f} |v_{si}|$  where  $v_{si}$  is the *i*th element of  $V_s$ .

**Theorem 1.** Algorithm (g1)–(g4) is stable for all initial value  $U_s^{(0)}$  satisfying

$$\|U_s^{(0)} - V_s\| < R, \tag{13}$$

where

$$R = \underline{v} - \frac{1}{\underline{v}} \| Z_d S_s^* \|.$$
(14)

**Proof.** The proposed algorithm is a mapping from  $U_s$  to itself which essentially is a kind of fixed-point iteration and can be rewritten as

$$U_{s}^{(k+1)} = g\left(U_{s}^{(k)}\right) = V_{0}\mathbf{1}_{n-1} - Z_{d}\left[\frac{s_{i}^{*}}{u_{i}^{(k)*}}\right]_{\mathcal{N}\setminus f}^{I}.$$
(15)

Based onBanach fixed-point theorem, the fixed-point  $V_s$  satisfying  $V_s = g(V_s)$  exists and is unique if  $g(U_s)$  is a contraction mapping on  $U_s$ . It can be seen that

$$\|U_{s}^{(k+1)} - V_{s}\| = \|g\left(U_{s}^{(k)}\right) - g(V_{s})\| = \left\|Z_{d}S_{s}^{*}\left[\frac{1}{v_{si}^{*}} - \frac{1}{u_{i}^{(k)*}}\right]_{\mathcal{N}\setminus f}^{T}\right\|$$

$$\leq \|Z_{d}S_{s}^{*}\|\left\|\left[\frac{u_{i}^{(k)*} - v_{si}^{*}}{u_{i}^{(k)*}v_{si}^{*}}\right]_{\mathcal{N}\setminus f}^{T}\right\| \leq L^{(k)}\|U_{s}^{(k)} - V_{s}\|,$$
(16)

where  $L^{(k)} = \frac{\|Z_d S_s^*\|}{|\underline{v}_s u^{(k)}|}$  with  $|\underline{v}_s u^{(k)}| = \min_{i \in \mathcal{N} \setminus f} |\underline{v}_{si} u^{(k)}_i|$ .

Due to  $||U_s^{(k)} - V_s|| < R$ , one has  $|u_i^{(k)} - v_{si}| < R$  and  $|u_i^{(k)}| > |v_{si}| - R$  for all  $i \in \mathcal{N} \setminus f$ . Subsequently  $|\underline{v}_s u^{(k)}| > (|v_{si}| - R)|v_{si}|$  for one  $i \in \mathcal{N} \setminus f$ . Furthermore,  $(|v_{si}| - R)|v_{si}| \ge \underline{v}(\underline{v} - R)$  for all  $i \in \mathcal{N} \setminus f$ . Therefore, due to Equation (13),

$$L^{(k)} < \frac{\|Z_d S_s^*\|}{\underline{v}(\underline{v} - R)} \le 1,$$
(17)

which implies that  $\|U_s^{(k+1)} - V_s\| < R$ . It together with the initial condition  $\|U_s^{(0)} - V_s\| < R$  shows that  $f(U_s)$  is a contraction mapping on  $U_s$ .  $\Box$ 

**Remark 1.** The condition in Equation (14) implies that  $\underline{v} > R > 0$  and  $\underline{v}^2 > \|Z_d S_s^*\|$ .

**Remark 2.** A smaller  $||Z_d S_s^*||$  will lead to a larger R. This implies roughly that a strong network (a small ratio between transmitted power and the branch admittance,  $s_i/(z_{Li})^{-1}$ ) allows a large permissible region for direct approaches of power flow. Meanwhile a small  $||Z_d S_s^*||$  means a small Lipschitz constant  $L^{(k)}$  and subsequently a fast convergence.

# 3.4. Comparison to the Direct Approach

It can be seen that our Algorithm (g1)–(g3) is similar to that in [8]. This is not surprising, in that both are based on Equation (12), the mapping from  $I_s$  to  $U_d$ . The difference is how to obtain Equation (12). The distinctive feature of our graph-based method is to present a much simpler way than the direct approach in [8].

Recall the CVB equation obtained by the direct approach [8] as: for radial networks,

$$U_d = (BCBV)(BIBC)I_s,$$

and for meshed networks,

$$\begin{bmatrix} U_d \\ 0 \end{bmatrix} = (\text{BCBV})(\text{BIBC}) \begin{bmatrix} I_s \\ B_{new} \end{bmatrix},$$

followed by a Kron's Reduction. The following comparison is stated.

In the radial network, the direct approach needs: (1) sequentially numbering nodes and edges from layer to layer beginning at the feeder node; (2) performing a six-step algorithm build the matrices (BCBV) and (BIBC) ; and (3) obtaining  $Z_d$  by multipling (BCBV) by (BIBC). Our method

needs: (1) an arbitrary numbering nodes and edges; (2) directly writing matrices H,  $H_s$ , and  $Z_L$ ; and (3) calculating  $Z_d = H_s^{-T} Z_L H_s^{-1}$ .

In the meshed network, the direct approach needs an extra drawing the corresponding radial version of the meshed network, adding two steps for every extra branch that makes the network meshed to build matrices (BCBV) and (BIBC), and a Kron's reduction.

Thus, the complexity of the direct approach would increase largely as the degree of mesh increases, while our method has the same procedure for the meshed network as that for the radial network. In fact, our method can apply to any meshed network rather than to the weakly-meshed network. Moreover, our method has a clearer physical meaning because no network reduction has to be made. The above contents are summarized in Table 1 for a clear insight.

	Radial Networks	
	Direct Approach	Proposed Method
Numbering	Sequential	Arbitrary
Matrices	BIBC, BCBV	$H_s, Z_L$
Operation	$Z_d = INV(BIBC \times BCBV)$	$Z_d = H_s^{-T} Z_L H_s^{-1}$
	Meshed Networks	
	Direct Approach	Proposed Method
Meshes	Need Recognition	No Need
Numbering	Sequential for radial structure and place meshes to the end	Arbitrary
Matrices	Modified <i>BIBC</i> , <i>BCBV</i>	$H_s, Z_L$
Operation	Modifying BIBC, BCBV and Applying Kron's Reduction	$Z_{d} = (H_{s}Z_{L}^{-1}H_{s}^{T})^{-1}$

Table 1. Comparison with Direct A	Approach
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# 4. Test Results

The proposed method is tested and compared on both radial and meshed networks on MATLAB. Table 2 shows the distribution systems of 14-, 33-, 69-, 84-, 119-, 135-, and 874-node radial networks and their meshed editions, which are from papers [16–20].

No. of Nodes	No. of Branch (Radial)	No. of Branch (Meshed)
14	13	16
33	32	37
69	68	73
84	83	96
119	118	132
135	134	156
874	873	900

Table 2. Network Configuration.

Four methods are tested. Method I is the Gauss–Seidel Method, Method II is the Newton–Raphson method, Method III is the direct approach proposed in [8] and Method IV is our Graph-based method. The convergence tolerance is set at 0.001 p.u.

Tables 3 and 4 show the performance of these four methods for radial and meshed networks, where "Time" and "ITs" denote the iteration time and iteration numbers, respectively, and "L" denotes the approximate Lipschitz constant, which describes the convergence rate of Method IV. According to the tables, Method I, the Gauss–Seidel method, costs much more time and iteration steps than the three other methods, since it has a very low convergence rate so that even if it does not need much time at each iteration, it still costs much time. Method II, the Newton–Raphson method, needs fewer iteration steps than other methods since it follows the direction of gradient descent at every step.

However, the Newton–Raphson method still costs more time than Method III and IV, because it requires calculating the Jacobian matrix and its inversion matrix at each iteration, which costs majority of time. Therefore, the time consumption of Newton–Raphson is more related to the number of nodes compared with Method III and IV.

No. of Nodes	Method	Ι	Method	II	Method 1	III	Meth	od IV	
ivo. of ivodes	Time	ITs	Time	ITs	Time	ITs	Time	ITs	L
14	$7.7  imes 10^{-3}$	25	$5.2  imes 10^{-4}$	2	$1.2  imes 10^{-5}$	2	$1.1  imes 10^{-5}$	2	0.012
33	0.31	433	$1.2 imes10^{-3}$	3	$2.5 imes10^{-5}$	4	$2.5 imes10^{-5}$	4	0.091
69	0.75	476	$2.4 imes10^{-3}$	3	$3.7 imes10^{-5}$	4	$3.6 imes10^{-5}$	4	0.17
84	0.7	353	$3.8 imes10^{-3}$	3	$4.4 imes10^{-5}$	4	$4.4 imes10^{-5}$	4	0.097
119	1.4	493	$5.7  imes 10^{-3}$	3	$7.6 imes10^{-5}$	4	$7.5  imes 10^{-5}$	4	0.14
135	1.7	519	$6.0 imes10^{-3}$	3	$6.5 imes10^{-5}$	3	$6.6 imes10^{-5}$	3	0.091
874	0.78	27	0.67	3	$4.0  imes 10^{-3}$	4	$4.0  imes 10^{-3}$	4	0.056

Table 3. Radial Network Test.

Table 4. 1	Meshed	Network	Test.
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No. of Nodes	Method	Ι	Method	II	Method 1	II	Μ	ethod	IV
No. of Nodes	Time	ITs	Time	ITs	Time	ITs	Time	ITs	L
14	$8.6 imes10^{-3}$	30	$5.3 imes10^{-4}$	2	$1.2  imes 10^{-5}$	2	$1.1  imes 10^{-5}$	2	$7.5  imes 10^{-3}$
33	0.16	220	$1.2  imes 10^{-3}$	3	$1.9  imes 10^{-5}$	3	$1.9  imes 10^{-5}$	3	0.05
69	0.28	178	$3.0 imes10^{-3}$	3	$2.7 imes10^{-5}$	3	$2.7 imes10^{-5}$	3	0.066
84	0.65	336	$3.0 imes10^{-3}$	3	$3.4 imes10^{-5}$	3	$3.3 imes10^{-5}$	3	0.059
119	0.78	283	$4.8 imes10^{-3}$	3	$5.7 imes10^{-5}$	3	$5.6 imes10^{-5}$	3	0.061
135	1.6	504	$6.1  imes 10^{-3}$	3	$6.5 imes10^{-5}$	3	$6.6 imes10^{-5}$	3	0.028
874	0.5	17	0.45	2	$2.0  imes 10^{-3}$	2	$2.1  imes 10^{-3}$	2	$7.9 imes10^{-3}$

As for Method III and Method IV, the results show the direct approach is approximately equivalent to the proposed method, which is consistent with the theoretical analysis. However, the advantage of our method is the process to get the CVB equation. As mentioned above, the direct approach requires loop-analysis and Kron's Reduction for meshes networks, while our method does not need any extra processing. Table 5 provides a comparison of time spent on getting the CVB equation, and the result shows that the proposed method takes less time than the direct approach to get the CVB equation. Moreover, for the example of the 14-node network, we only increase the number of mesh, and the results show that the time consumptions of proposed method are almost the same when the number of meshes increased, unlike the increasing time consumptions of direct approach, mainly due to the Kron's Reduction. Note that, in the case that the node numbering, the radial structure drawing and the meshed branch identifying have been made in advance, it can be seen that the graph-based method is much better if the time spent on these pretreatments are contained.

# Limitations

The proposed method has shown obvious advantage compared to previous works. However, it has some limitations. First, as for the impact of distributed generators, DGs can be considered as PQ nodes with constant active/reactive powers as well as PV nodes with constant active power and voltage magnitude. If DGs are considered as PQ nodes, our method can deal with it. Alternatively, if DGs are considered as PV nodes, our method is not applicable. Besides, as for the applicability to unbalanced networks, we have a preliminary outline to extend our method for unbalanced networks, but the work still needs deliberate discussion and proof. Finally, the impact of FACTS devices is not discussed in this paper since FACTS devices cannot be simply modeled as PQ nodes and our method aims at distribution networks with PQ nodes.

No. of Node	No. of Branch	Method IV Time	Method III Time
14	16	$5.5  imes 10^{-5}$	$1.2  imes 10^{-4}$
33	37	$1.3 imes10^{-4}$	$2.6 imes10^{-4}$
69	73	$3.9 imes10^{-4}$	$6.4 imes10^{-4}$
84	96	$5.2 imes10^{-4}$	$9.2 imes10^{-4}$
119	132	$8.2 imes10^{-4}$	$1.4 imes10^{-3}$
135	156	$1.1 imes10^{-3}$	$1.8 imes10^{-3}$
874	900	$9.5 imes10^{-2}$	$1.1 imes10^{-1}$
No. of Node	No. of Branch	Method IV Time	Method III Time
No. of Node	No. of Branch	Method IV Time $5.5 \times 10^{-5}$	Method III Time $1.2 \times 10^{-4}$
<b>No. of Node</b> 14 14	No. of Branch 3 4	$\begin{array}{c} \textbf{Method IV Time} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \end{array}$	$\begin{array}{c} \textbf{Method III Time} \\ 1.2 \times 10^{-4} \\ 1.5 \times 10^{-4} \end{array}$
No. of Node 14 14 14	<b>No. of Branch</b> 3 4 5	$\begin{array}{c} \mbox{Method IV Time} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \end{array}$	$\begin{array}{c} \mbox{Method III Time} \\ 1.2 \times 10^{-4} \\ 1.5 \times 10^{-4} \\ 1.7 \times 10^{-4} \end{array}$
No. of Node 14 14 14 14 14	<b>No. of Branch</b> 3 4 5 6	$\begin{array}{c} \text{Method IV Time} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \end{array}$	$\begin{array}{c} \textbf{Method III Time} \\ 1.2 \times 10^{-4} \\ 1.5 \times 10^{-4} \\ 1.7 \times 10^{-4} \\ 2.5 \times 10^{-4} \end{array}$
No. of Node 14 14 14 14 14 14	No. of Branch 3 4 5 6 7	$\begin{array}{c} \text{Method IV Time} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.6 \times 10^{-5} \end{array}$	$\begin{array}{c} \textbf{Method III Time} \\ 1.2 \times 10^{-4} \\ 1.5 \times 10^{-4} \\ 1.7 \times 10^{-4} \\ 2.5 \times 10^{-4} \\ 2.6 \times 10^{-4} \end{array}$
No. of Node 14 14 14 14 14 14 14 14	No. of Branch 3 4 5 6 7 8	$\begin{array}{c} \textbf{Method IV Time} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.5 \times 10^{-5} \\ 5.6 \times 10^{-5} \\ 5.6 \times 10^{-5} \end{array}$	$\begin{array}{c} \textbf{Method III Time} \\ 1.2 \times 10^{-4} \\ 1.5 \times 10^{-4} \\ 1.7 \times 10^{-4} \\ 2.5 \times 10^{-4} \\ 2.6 \times 10^{-4} \\ 2.6 \times 10^{-4} \end{array}$

Table 5. Time to get the CVB equation.

# 5. Conclusions

This paper has proposed a graph-based power flow method for three-phase balanced distribution systems with *PQ* nodes. For the nature of distribution systems such as radial/weakly meshed structure, and large number of branches and nodes, traditional power flow method may fail or cannot meet the requirement. With regards to this, we have made some progress. The proposed method provides a uniform formulation for both radial and meshed networks. The uniform formulation is much simpler than before, requiring neither a tricky numbering and layering of nodes nor breaking meshes and loop-analysis, which are both necessary in previous works for meshed networks. The convergence of the proposed method has been shown by using the Banach fixed-point theorem. The comparison test results show the efficiency of our competitive method.

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